

# **An Architecture for Distributed Source Localization in Wireless Sensor Network**

Problem:

Estimate the location of a moving  
source using sensor networks

Solution:

Algorithms:

Collaborative signal processing (CSP)

Recent advancements:

Incremental non-linear Optimization for  
use with distributed/ decentralized CSP  
(1996, 2004)

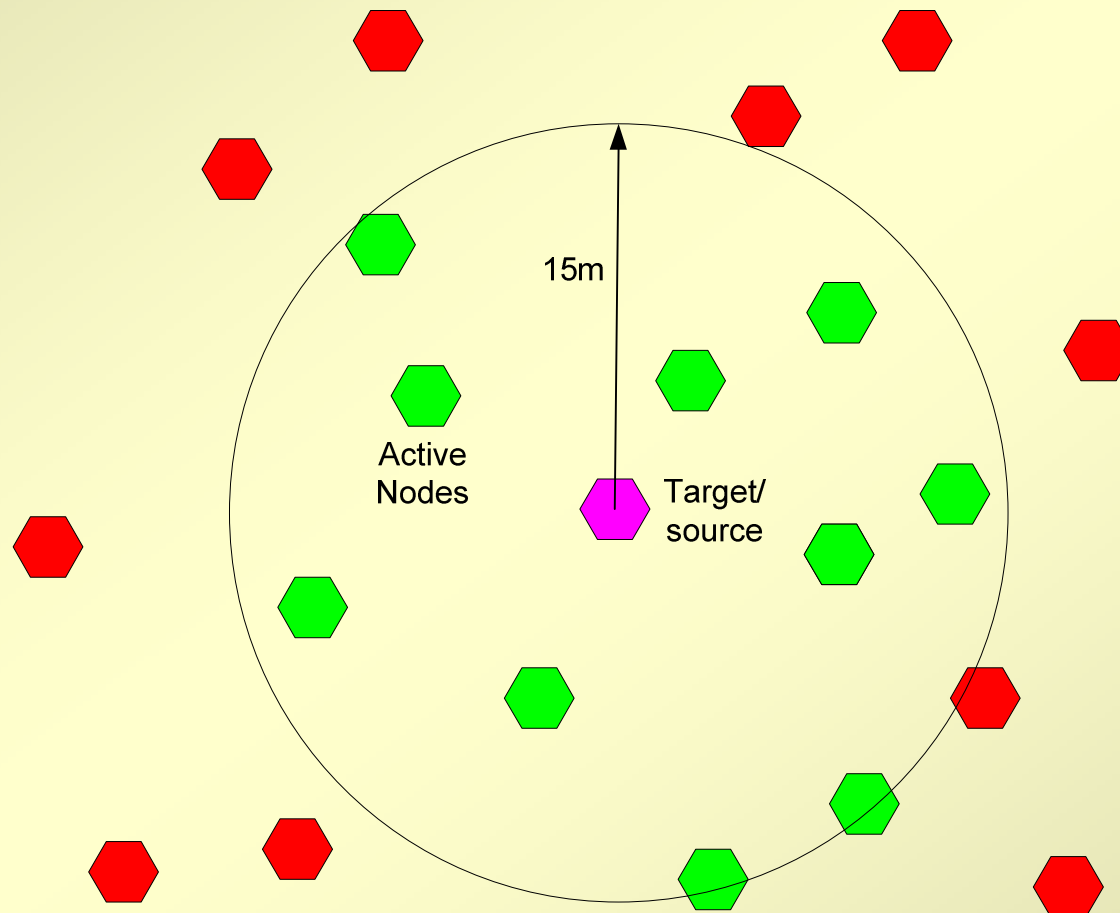
Architecture:

Topic of this presentation

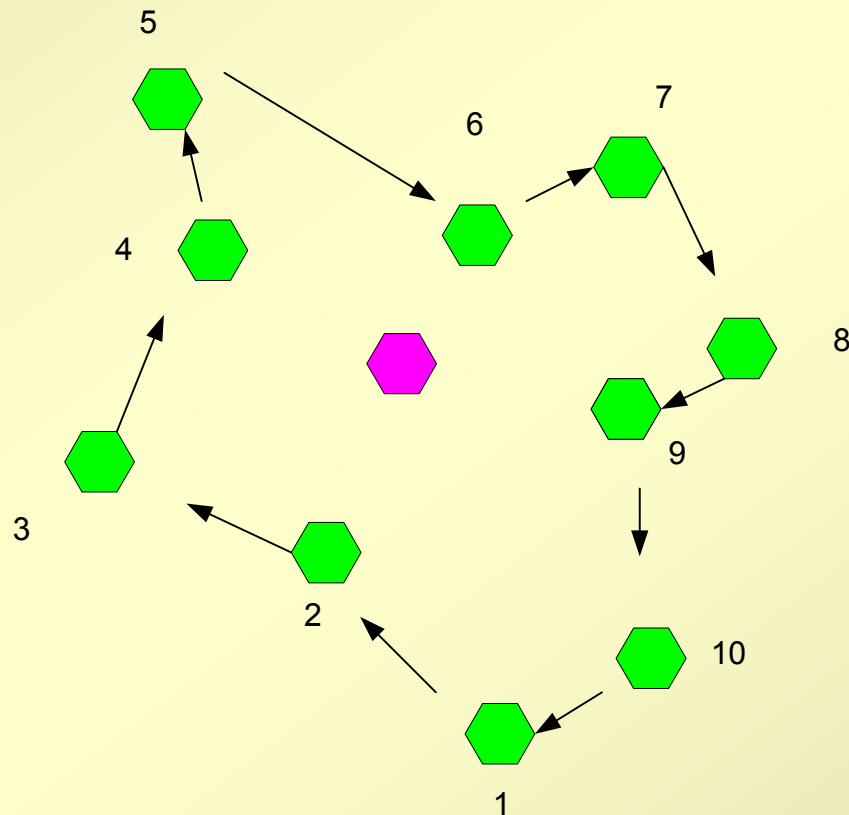
Literature:

Numerous *incoherent* protocols

Objective: Estimate the source/ target location



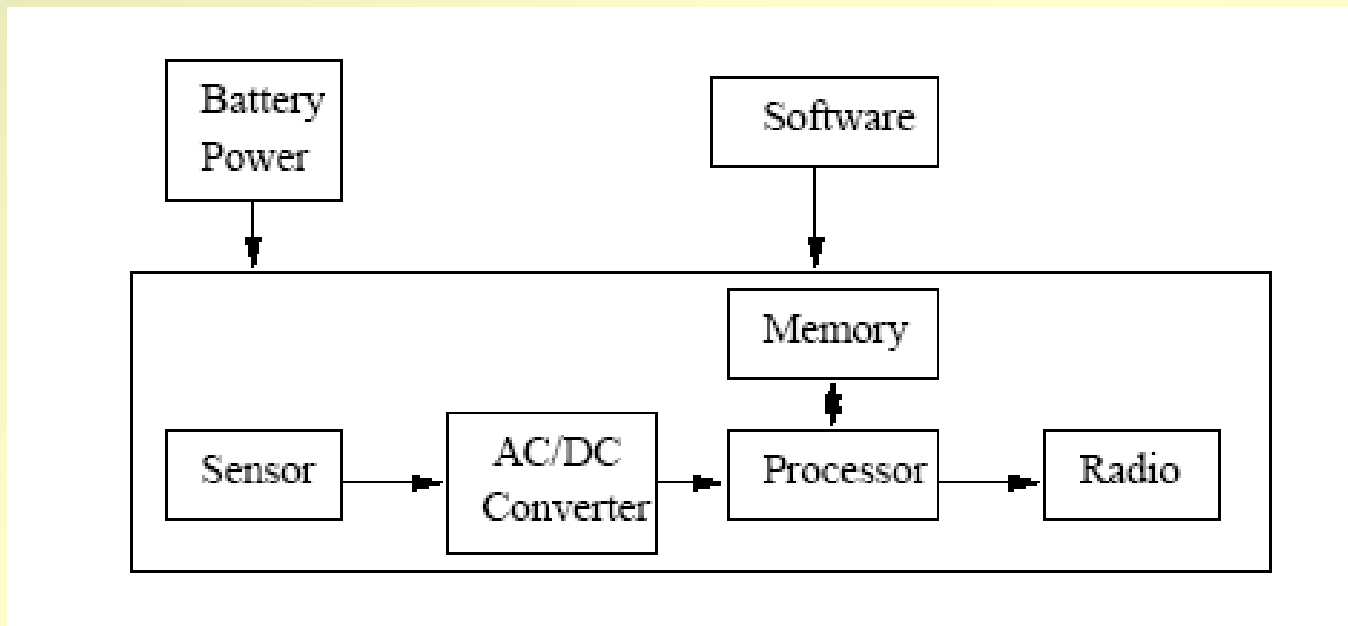
Objective: Estimate the source/ target location using



Node-1 back to Node-1: thatz one cycle



## Components



## Requirements (of the algorithm) to be met (by the architecture):

- Ad-hoc network
- Self-organizing MAC
- Scalable
- Node insertion (new node)
- Node deletion (node out-of-range)
- Directed graph (cyclical communication)

## Pessimist's take!

- **Conserve energy**
  - Bad nodes
  - node mobility
  - refresh a static network

NO SINGLE SOLUTION SATISFIES ALL

## Assumptions:

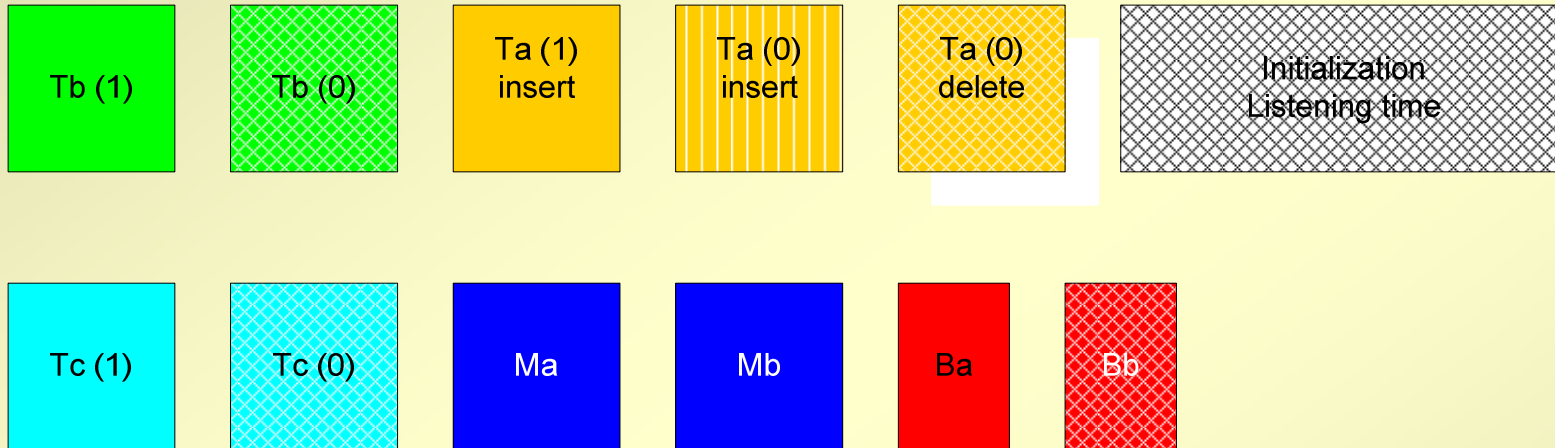
- **Absolute time synchronization**
- Nodes aware of their locations
- Can simultaneously receive two frames operating in two bands

## Risk factors:

- **Absolute time synchronization: difficult in practice.** As an example, USA's patriot missiles could not detect the Scud missiles launched by Iraq during Gulf-war just because the clocks were off by 1/3s which further resulted in estimating the target off by 600m
- The network formed may collapse if the initial link formed is too bad a guess
- If the source is moving too fast, latency in network formation may be counter-productive in the distributed CSP setting



## Types of frames



### Broadcast

- TYPE-a (ta)
- TYPE-b (tb)
- TYPE-c (tc)
- Beacon-a (ba)
- Beacon-b (bb)

### Message

- message-a (ma)
- message-b (mb)





## Parameters/ Variables

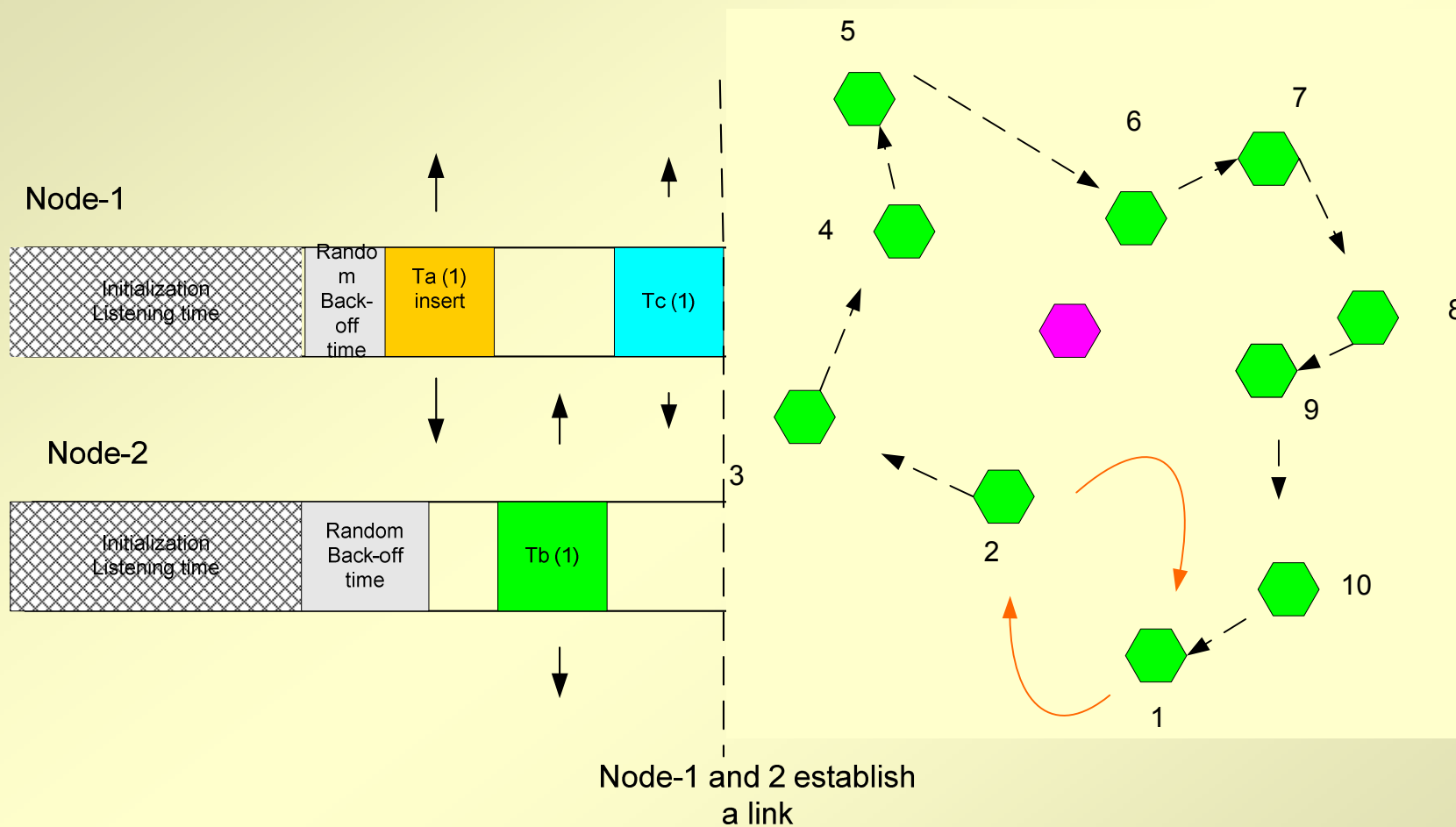
N: current number of sensors in the network  
T<sub>Tx</sub>: Transmission time  
T<sub>Rx</sub>: Reception time  
T<sub>rand\_backoff\_max</sub>: maximum backoff time  
t<sub>bkoff</sub>: random back-ff time  
Counter<sub>cycle</sub>: current cycle number  
Counter<sub>total\_cycles</sub>:  $\log_2(N)$   
T<sub>alg</sub>: algorithm time  
T<sub>initiaze</sub>: Initialization listening time



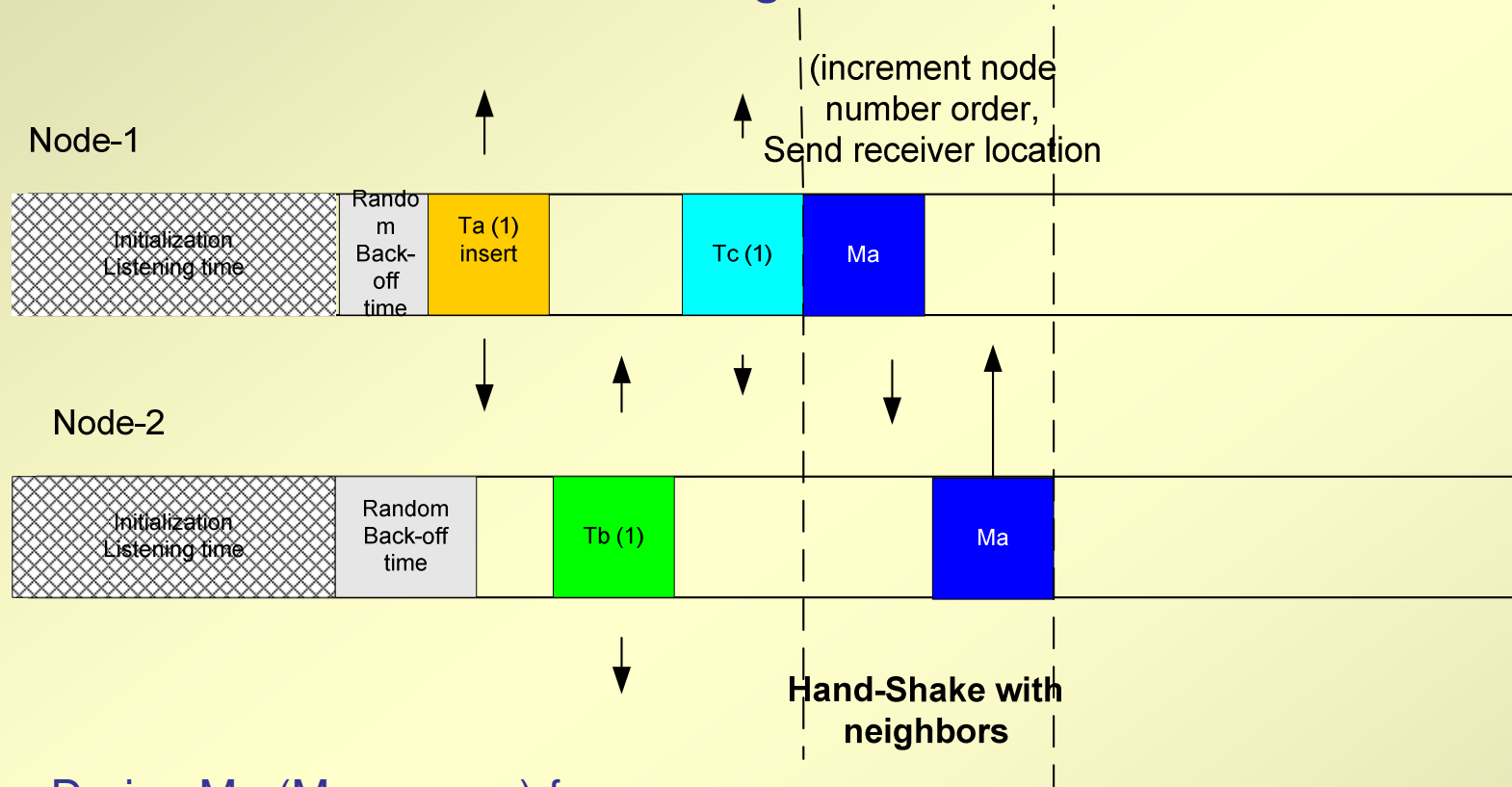
## Frequency bands

f<sub>d</sub>: delete a node (used in ta)  
f<sub>i</sub>: insert a node (used in ta, tb and tc)  
f<sub>b</sub>: beacon frame  
f<sub>m</sub>: messege frame (used in ma,mb)

## Very Initial Link formation



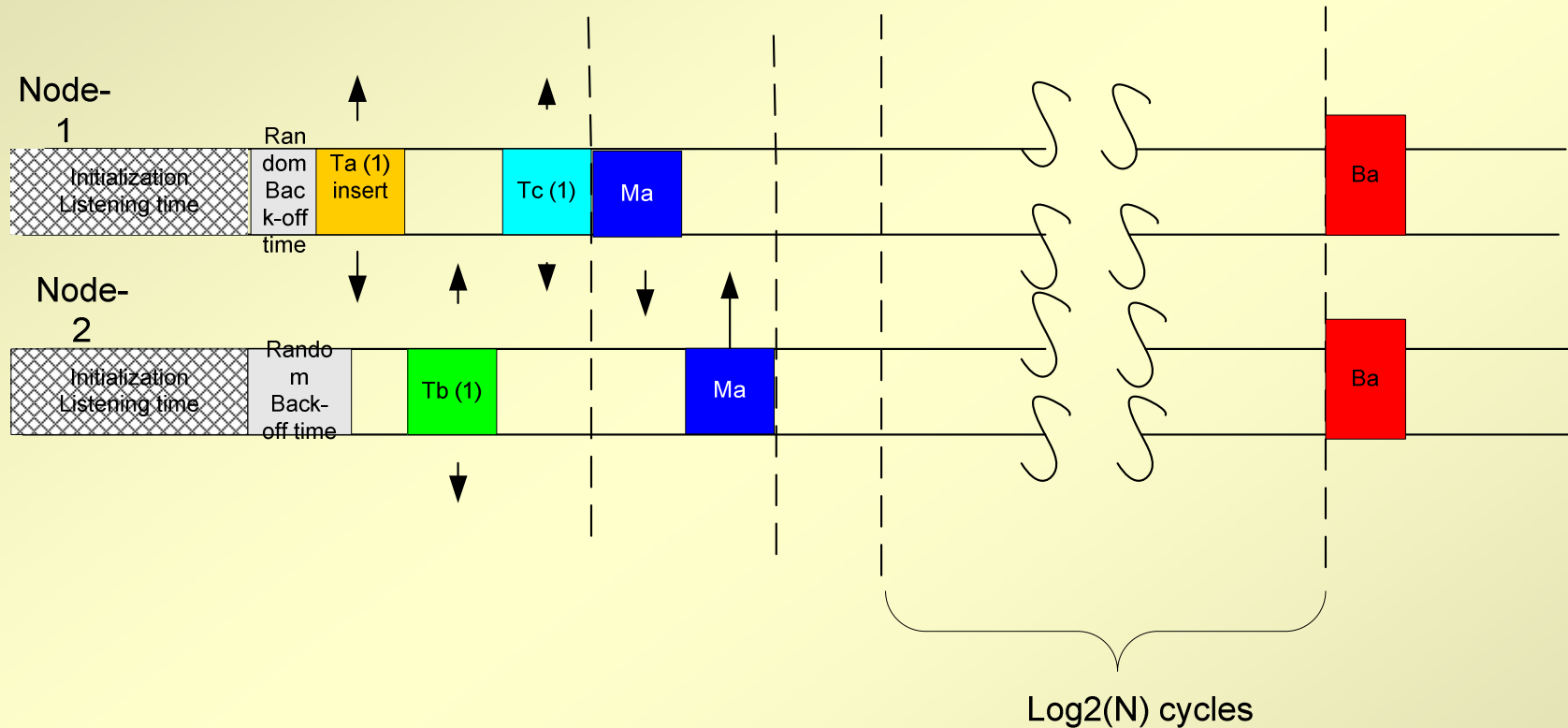
## Information exchange in the new link



### During Ma (Message-a) frame

- node number, intended receiver location and total number of nodes are transmitted
- Logical indices are formed at the end of this cycle which helps in defining the TDMA slots

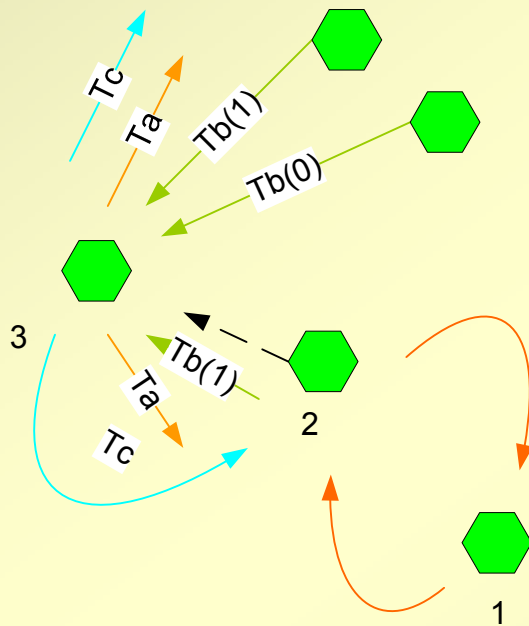
## Network management phase begins to check for inserting new nodes or deleting nodes



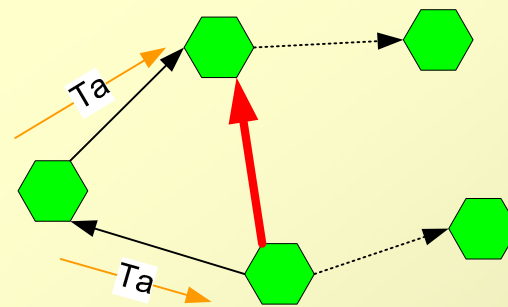
- At the end of beacon interval, nodes can send ta type frames
- The routing table is exactly like a linked-list except that it is maintained in pieces by all the nodes

## Addition and Deletion of a node

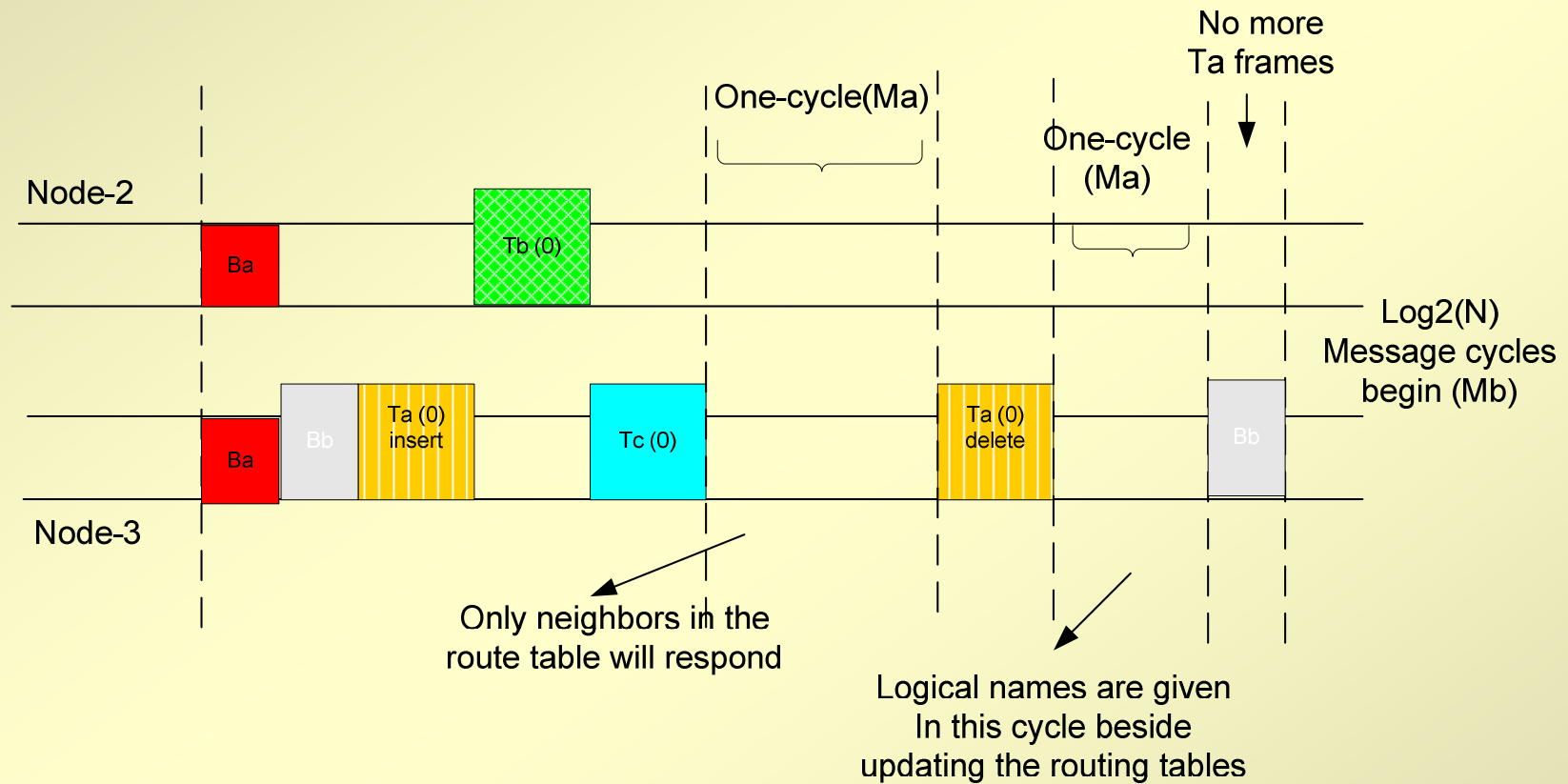
Adding a node



Deleting a node



## Time-Diagram



## Conclusions:

- meets “requirements”
- works under “assumptions”
- has “risks” associated with
- still evolving far away from implementation

## Future course:

- define a software architecture for implementation
- and implement

## Acknowledgements:

- Dr. Chang, Instructor, ee543x, ISU
- Bernard Lwakabamba, Grad. Student, ECpE Dept, ISU
- WWW
- and patient audience

## References:

Soma Sekhar Dhavala, [ee543x project report](#)