

Contextual Abductive Reasoning with Side-Effects

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How to explain the forest fire?



we observe a forest fire

How to explain the forest fire?



we observe a forest fire

forest fire occurred because of lightning



lightning

forest fire occurred because of barbecue



barbecue

How to explain the forest fire?



we observe a storm

forest fire occurred because of lightning

only if a storm was observed



lightning



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We want to express that when observing a storm, the forest fire occurred because of a lightning. Otherwise, it occurred because of a barbecue.

Motivation

Evans, Barston, and Pollard carried out an experiment where participants had to evaluate whether certain syllogisms are valid.

- ▶ Their answers were strongly influenced by the **contextual setting** and whether the syllogisms' conclusions confirmed or contradicted the participants' **beliefs**.

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Classical logic cannot adequately represent this syllogistic reasoning task.

Hölldobler and Kencana Ramli [2009] propose to model human reasoning by

- ▶ logic programs
- ▶ under weak completion semantics
- ▶ based on the three-valued Łukasiewicz (1920) logic.

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While modeling this syllogistic reasoning task under weak completion semantics we needed to define a **contextual abductive framework**.

Logic Programs

We restrict ourselves to datalog programs. A logic program \mathcal{P} is a finite set of clauses

$$A \leftarrow A_1 \wedge \dots \wedge A_n \wedge \neg B_1 \wedge \dots \wedge \neg B_m, \quad (1)$$

$$A \leftarrow \perp, \quad (2)$$

- ▶ where A and A_i , $0 \leq i \leq n$, are **atoms** and $\neg B_j$, $1 \leq j \leq m$, are **negated atoms**.
- ▶ If $i = 0$, then we write $A \leftarrow \top$, which is called a **positive fact**.
- ▶ A clause of the form (2) is called a **negative fact**.
- ▶ A is **undefined** if it is not the head of any clause.
- ▶ $g\mathcal{P}$ denotes **ground** \mathcal{P} , that is, it consists of all the ground instances of its clauses.
- ▶ $\text{undef}(\mathcal{P})$ is the **set of all undefined atoms** in $g\mathcal{P}$.

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The following transformation is the **weak completion** of \mathcal{P}

1. Replace all clauses in $g\mathcal{P}$ with the same head $A \leftarrow body_1, \dots, A \leftarrow body_n$ by the single expression $A \leftarrow body_1 \vee \dots \vee body_n$.
2. Replace all occurrences of \leftarrow by \leftrightarrow .

Three-Valued Łukasiewicz [1920] Logic

		\neg
\top		\perp
\perp		\top
\mathbf{U}		\mathbf{U}

\wedge		\top	\mathbf{U}	\perp
\top		\top	\mathbf{U}	\perp
\mathbf{U}		\mathbf{U}	\mathbf{U}	\perp
\perp		\perp	\perp	\perp

\vee		\top	\mathbf{U}	\perp
\top		\top	\top	\top
\mathbf{U}		\top	\mathbf{U}	\mathbf{U}
\perp		\top	\mathbf{U}	\perp

\leftarrow_L		\top	\mathbf{U}	\perp
\top		\top	\top	\top
\mathbf{U}		\mathbf{U}	\top	\top
\perp		\perp	\mathbf{U}	\top

\leftrightarrow_L		\top	\mathbf{U}	\perp
\top		\top	\mathbf{U}	\perp
\mathbf{U}		\mathbf{U}	\top	\mathbf{U}
\perp		\perp	\mathbf{U}	\top

Table: \top , \perp , and \mathbf{U} denote *true*, *false*, and *unknown*, respectively.

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\perp		\top
\mathbf{U}		\mathbf{U}

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\top		\top	\mathbf{U}	\perp
\mathbf{U}		\mathbf{U}	\mathbf{U}	\perp
\perp		\perp	\perp	\perp

\vee		\top	\mathbf{U}	\perp
\top		\top	\top	\top
\mathbf{U}		\top	\mathbf{U}	\mathbf{U}
\perp		\top	\mathbf{U}	\perp

\leftarrow_L		\top	\mathbf{U}	\perp
\top		\top	\top	\top
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An **interpretation** I of \mathcal{P} is a mapping of the **Herbrand base** $\mathcal{B}_{\mathcal{P}}$ to $\{\top, \perp, \mathbf{U}\}$ and is represented by an unique pair, $\langle I^{\top}, I^{\perp} \rangle$, where

$$I^{\top} = \{A \in \mathcal{B}_{\mathcal{P}} \mid A \text{ is mapped to } \top\} \text{ and } I^{\perp} = \{A \in \mathcal{B}_{\mathcal{P}} \mid A \text{ is mapped to } \perp\}.$$

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\top		\top	\top	\top
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\leftarrow_L		\top	\mathbf{U}	\perp
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- ▶ For every I it holds that $I^{\top} \cap I^{\perp} = \emptyset$.
- ▶ A **model of a formula** F is an interpretation I such that F is true under I .
- ▶ A **model of** $g\mathcal{P}$ is an interpretation that is a model of each clause in $g\mathcal{P}$.

Computing Least Models

Hölldobler and Kencana Ramli [2009] propose to compute the **least model of the weak completion of \mathcal{P}** ($\text{lm}_{\text{wc}}\mathcal{P}$) which is identical to the **least fixed point of $\Phi_{\mathcal{P}}$** , by an operator defined by Stenning and van Lambalgen [2008].

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Let I be an interpretation in $\Phi_{\mathcal{P}}(I) = \langle J^{\top}, J^{\perp} \rangle$, where

$$J^{\top} = \{A \mid \text{there exists } A \leftarrow \text{body} \in \text{g} \mathcal{P} \text{ with } I(\text{body}) = \top\},$$

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In (Dietz, Hölldobler, and Wernhard [2013]) we show that weak completion semantics corresponds to **well-founded semantics** for modified tight logic programs.

- ▶ \mathcal{P} is **tight**, that is, it does not contain positive cycles.
- ▶ **Modified \mathcal{P}** is $\mathcal{P} \cup \{A \leftarrow \neg nA, nA \leftarrow \neg A \mid \text{for all undefined atoms } A \in \mathcal{P}\}$.

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- ▶ set of **abducibles** \mathcal{A} contains all positive and negative facts of each $A \in \text{undef}(\mathcal{P})$,
- ▶ \mathcal{E} is an **explanation** and a consistent subset of \mathcal{A} ,
- ▶ **logical consequence relation** $\models_{\mathcal{L}}^{\text{Imwc}}$, where $\mathcal{P} \models_{\mathcal{L}}^{\text{Imwc}} F$ iff $\text{Im}_{\mathcal{L}}\text{wc } \mathcal{P}(F) = \top$, and
- ▶ \mathcal{O} is an **observation** which is a set of (at least one) literals.

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\mathcal{O} is **explained by \mathcal{E} given \mathcal{P}** iff $\mathcal{P} \cup \mathcal{E} \models_{\mathcal{L}}^{\text{Imwc}} \mathcal{O}$, where $\mathcal{P} \not\models_{\mathcal{L}}^{\text{Imwc}} \mathcal{O}$.

\mathcal{O} is **explained given \mathcal{P}** iff there exists an \mathcal{E} such that \mathcal{O} is explained by \mathcal{E} given \mathcal{P} .

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F follows skeptically from \mathcal{P} , and \mathcal{O} iff \mathcal{O} can be explained given \mathcal{P} , and for all minimal explanations \mathcal{E} we find that $\mathcal{P} \cup \mathcal{E} \models_{\mathcal{L}}^{\text{lmwc}} \mathcal{O}$.

F follows credulously from \mathcal{P} , and \mathcal{O} iff there exists a minimal explanation \mathcal{E} such that $\mathcal{P} \cup \mathcal{E} \models_{\mathcal{L}}^{\text{lmwc}} \mathcal{O}$.

Explaining the Forest Fire

Consider $\mathcal{P}_{fire,dry}$

$storm$	\leftarrow	$lightning \wedge tempest,$	$forest_fire$	\leftarrow	$lightning \wedge smoke \wedge \neg ab,$
ab	\leftarrow	$\neg dry_leaves,$	$forest_fire$	\leftarrow	$barbecue \wedge smoke \wedge \neg ab,$
$smoke$	\leftarrow	$\top,$	dry_leaves	\leftarrow	$\top.$

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Let us assume that we observe a forest fire

$$\mathcal{O}_{forest_fire} = \{forest_fire\}$$

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$smoke$	\leftarrow	\top ,	dry_leaves	\leftarrow	\top .

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The set of abducibles, $\mathcal{A}_{\mathcal{P}_{fire,dry}}$, is

$barbecue$	\leftarrow	\top ,	$barbecue$	\leftarrow	\perp ,
$lightning$	\leftarrow	\top ,	$lightning$	\leftarrow	\perp ,
$tempest$	\leftarrow	\top ,	$tempest$	\leftarrow	\perp .

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There are two minimal explanations for $\mathcal{O}_{forest_fire}$

$$\mathcal{E}_{barbecue} = \{barbecue \leftarrow \top\},$$

$$\mathcal{E}_{lightning} = \{lightning \leftarrow \top\}.$$

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There are two minimal explanations for $\mathcal{O}_{forest_fire}$

$$\mathcal{E}_{barbecue} = \{barbecue \leftarrow \top\},$$

$$\mathcal{E}_{lightning} = \{lightning \leftarrow \top\}.$$

The least model of the weak completion of $\mathcal{P}_{fire,dry}$ with each of them is

$$\text{Im}_{\perp} \text{WC}(\mathcal{P}_{fire,dry} \cup \mathcal{E}_{barbecue}) = \langle \{forest_fire, barbecue, smoke, dry_leaves\}, \{ab\} \rangle,$$

$$\text{Im}_{\perp} \text{WC}(\mathcal{P}_{fire,dry} \cup \mathcal{E}_{lightning}) = \langle \{forest_fire, lightning, smoke, dry_leaves\}, \{ab\} \rangle.$$

Contextual Abductive Reasoning

How to express that *barbecue* describes the *usual* and *lightning* the *exceptional case*, only explaining the *forest fire* in the context of a *storm*?

Contextual Abductive Reasoning

How to express that *barbecue* describes the usual and *lightning* the exceptional case, only explaining the *forest fire* in the context of a *storm*?

Introduce for every A , two reserved (meta-) predicates (Pereira and Pinto [2011]).

$\text{inspect}(A)$ and $\text{inspect}_{\text{not}}(A)$

These are special abducibles only to be abduced if A or $\neg A$ are abduced somewhere else.

Contextual Side-effects

Consider $\mathcal{P}_{fire,dry}$

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forest fire

$$\leftarrow \underline{\underline{\mathcal{E}_{inspect(lightning)} = \{inspect(lightning) \leftarrow \top\}}}$$



lightning

$$\leftarrow \underline{\underline{\mathcal{E}_{barbecue} = \{barbecue \leftarrow \top\}}}$$



barbecue

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$\mathcal{E}_{inspect(lightning)} = \{inspect(lightning) \leftarrow \top\}$



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Contextual Side-effects

Consider $\mathcal{P}_{\text{fire,dry}}$

$\text{storm} \leftarrow \text{lightning} \wedge \text{tempest}, \quad \text{forest_fire} \leftarrow \text{inspect}(\text{lightning}) \wedge \text{smoke} \wedge \neg \text{ab},$
 $\text{ab} \leftarrow \neg \text{dry_leaves}, \quad \text{forest_fire} \leftarrow \text{barbecue} \wedge \text{smoke} \wedge \neg \text{ab},$
 $\text{smoke} \leftarrow \top, \quad \text{dry_leaves} \leftarrow \top.$



storm

$$\begin{aligned} \mathcal{E}_{\text{lightning}, \text{tempest}} &= \{ \text{lightning} \leftarrow \top, \text{tempest} \leftarrow \top \} \\ \mathcal{E}_{\text{insp}(\text{lightning})} &= \{ \text{inspect}(\text{lightning}) \leftarrow \top \} \end{aligned}$$



lightning



forest fire

$$\mathcal{E}_{\text{barbecue}} = \{ \text{barbecue} \leftarrow \top \}$$



barbecue

$\text{inspect}(\text{lightning}) \leftarrow \top$ can only be abduced by $\mathcal{O}_{\text{forest_fire}} = \{ \text{forest_fire} \}$ if $\text{lightning} \leftarrow \top$ is explained somewhere else, e.g. by observing a *storm*.

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storm

$\mathcal{E}_{lightning, tempest} = \{ lightning \leftarrow \top, tempest \leftarrow \top \}$
 $\mathcal{E}_{inspect(lightning)} = \{ inspect(lightning) \leftarrow \top \}$



lightning



forest fire

~~$\mathcal{E}_{barbecue} = \{ barbecue \leftarrow \top \}$~~



barbecue

$\mathcal{O}_{forest_fire}$ is a **contextual side-effect** of \mathcal{O}_{storm} given $\mathcal{P}_{fire,dry}$.

Contestable Contextual Side-effects

Consider $\mathcal{P}_{fire,rained}$

$storm \leftarrow lightning \wedge tempest,$ $forest_fire \leftarrow inspect(lightning) \wedge smoke \wedge \neg ab,$
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- ▶ The *rained*-clause states that if for some other observation, we explained that the leaves are not dry, then it rained.

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- ▶ When we observe $\mathcal{O}_{forest_fire}$, we need to abduce $dry_leaves \leftarrow \top$.

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- ▶ The *rained*-clause states that if for some other observation, we explained that the leaves are not dry, then it rained.
- ▶ When we observe $\mathcal{O}_{forest_fire}$, we need to abduce $dry_leaves \leftarrow \top$.
- ▶ This will make $inspect_{not}(dry_leaves)$ false and we conclude that it did not rain.

Contestable Contextual Side-effects

Consider $\mathcal{P}_{fire,rained}$

$$\begin{array}{ll} storm \leftarrow lightning \wedge tempest, & forest_fire \leftarrow inspect(lightning) \wedge smoke \wedge \neg ab, \\ ab \leftarrow \neg dry_leaves, & forest_fire \leftarrow barbecue \wedge smoke \wedge \neg ab, \\ smoke \leftarrow \top, & rained \leftarrow inspect_{not}(dry_leaves). \end{array}$$

- ▶ The *rained*-clause states that if for some other observation, we explained that the leaves are not dry, then it rained.
- ▶ When we observe $\mathcal{O}_{forest_fire}$, we need to abduce $dry_leaves \leftarrow \top$.
- ▶ This will make $inspect_{not}(dry_leaves)$ false and we conclude that it did not rain.

$\mathcal{O}_{\neg rained}$ is a **contested contextual side-effect** of $\mathcal{O}_{forest_fire}$ given $\mathcal{P}_{fire,rained}$.

Jointly Supported Contextual Relevant Consequences

Consider $\mathcal{P}_{\text{ffight}}$

$\text{smoke} \leftarrow \text{fire} \quad \wedge \quad \text{inspect}(\text{fire_fighters}),$

$\text{sirens} \leftarrow \text{inspect}(\text{fire}) \quad \wedge \quad \text{fire_fighters}.$

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- ▶ $\text{inspect}(\text{fire_fighters}) \leftarrow \top$ is only abducible if $\text{fire_fighters} \leftarrow \top$ is abduced.

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- ▶ Given $\mathcal{O}_{\text{sirens}}$, we can abduce $\text{fire_fighters} \leftarrow \top$ but not $\text{inspect}(\text{fire}) \leftarrow \top$.

Jointly Supported Contextual Relevant Consequences

Consider $\mathcal{P}_{\text{ffight}}$

```
smoke ← fire ∧ inspect(fire_fighters),  
sirens ← inspect(fire) ∧ fire_fighters.
```

- ▶ Let's observe $\mathcal{O}_{\text{smoke}} = \{\text{smoke}\}$. We abduce $\text{fire} \leftarrow \top$.
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If we observe both, $\mathcal{O}_{\text{smoke}}$ and $\mathcal{O}_{\text{sirens}}$ then the minimal explanation is

```
fire ← T, fire_fighters ← T,  
inspect(fire) ← T, inspect(fire_fighters) ← T.
```

Jointly Supported Contextual Relevant Consequences

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$$\begin{array}{ll} \text{fire} & \leftarrow \top, & \text{fire_fighters} & \leftarrow \top, \\ \text{inspect}(\text{fire}) & \leftarrow \top, & \text{inspect}(\text{fire_fighters}) & \leftarrow \top. \end{array}$$

$\mathcal{O}_{\text{smoke}}$ and $\mathcal{O}_{\text{sirens}}$ are jointly supported contextual relevant consequences given $\mathcal{P}_{\text{ffight}}$.

Conclusion

- ▶ Weak completion semantics is based on a previously proposed approach that seems to adequately model
 - ▶ Wason's selection task, and
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Conclusion

- ▶ Weak completion semantics is based on a previously proposed approach that seems to adequately model
 - ▶ Wason's selection task, and
 - ▶ Byrne's suppression task.
- ▶ While modeling other human reasoning tasks, we identified the need to express **contextual abductive reasoning**.
- ▶ By modeling contextual abduction with **inspection points**, more specific relations between observations can be defined as
 - ▶ contextual side-effects,
 - ▶ (jointly supported) contextual relevant consequences, and
 - ▶ contestable contextual side-effects.

Thank you very much for your attention!

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If something is a cigarette, then it is abnormal (wrt (3)). (4)

Modeling S_{add} of the Syllogistic Reasoning Task

\mathcal{P}_{add} represents the first two premises of S_{add}

$$add'(X) \leftarrow inex(X) \wedge \neg ab_{add'}(X), \quad add(X) \leftarrow \neg add'(X), \quad \text{PREMISE 1}$$

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To explain why people validate S_{add} we need to show that they reason **abductively**.

Abducing the CONCLUSION

Given our background knowledge we know, there are additive things, let's say about b

$$\mathcal{O}_{add(b)} = \{add(b)\}$$

We have two minimal explanations for $\mathcal{O}_{add(b)}$

$$\begin{aligned} \text{Im}_{\text{LWC}}(\mathcal{P}_{add} \cup \mathcal{E}_{cig(b)}) &= \langle \{add(b), cig(b), inex(b), \dots\}, \quad \{\dots\} \rangle \\ \text{Im}_{\text{LWC}}(\mathcal{P}_{add} \cup \mathcal{E}_{-cig(b)}) &= \langle \{add(b), \dots\}, \quad \{cig(b), inex(b), \dots\} \rangle \end{aligned}$$

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Recall \mathcal{P}_{add} . Together with $\mathcal{E}_{cig(b)}$ it contains

$$\begin{aligned} add'(X) &\leftarrow inex(X) \wedge \neg ab_{add'}(X), & add(X) &\leftarrow \neg add'(X), \\ inex(X) &\leftarrow cig(X) \wedge \neg ab_{inex}(X), & ab_{add'}(X) &\leftarrow cig(X), \\ ab_{add'}(X) &\leftarrow \perp, & ab_{inex}(X) &\leftarrow \perp, \\ cig(b) &\leftarrow \top. \end{aligned}$$

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Credulously, we validate **some addictive things are not cigarettes**.

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Rich people are millionaires. (1)

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*If someone is a hard worker and **not abnormal**, then this person is not a millionaire.* (3)
Nobody is abnormal (wrt (3)).

The belief in (1) and (2) would generate the exception for rich people

*If someone is rich, then this person is **abnormal** (wrt (3)).* (4)

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Even though not tested yet, our hypothesis is, while checking S_{rich} , participants did not make these assumptions and thus, had not been influenced by the belief-bias effect.

Contextual Abductive Reasoning

How to express that PREMISE 1 describes the **usual** and PREMISE 2 the **exceptional case**? Inexpensive cigarette should be the **exception** in the context of addictive things.

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How to express that PREMISE 1 describes the usual and PREMISE 2 the exceptional case? Inexpensive cigarette should be the exception in the context of addictive things.

Introduce for every A , two reserved (meta-) predicates (Pereira and Pinto [2011]).

$\text{inspect}(A)$ and $\text{inspect}_{\text{not}}(A)$

These are special abducibles only to be abduced if A or $\neg A$ are abduced somewhere else.

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$$\begin{array}{ll} add'(X) \leftarrow inex(X) \wedge \neg ab_{add'}(X), & add(X) \leftarrow \neg add'(X), \\ inex(X) \leftarrow cig(X) \wedge \neg ab_{inex}(X), & ab_{add'}(X) \leftarrow \text{inspect}(cig(X)), \\ ab_{add'}(X) \leftarrow \perp, & ab_{inex}(X) \leftarrow \perp, \end{array}$$

Contextual Abductive Reasoning

How to express that PREMISE 1 describes the usual and PREMISE 2 the exceptional case? Inexpensive cigarette should be the exception in the context of additive things.

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Suppose again b is additive, i.e. $\mathcal{O}_{add(b)} = \{add(b)\}$. $\mathcal{E}_{\text{cig}} = \{\text{cig}(b) \leftarrow \top\}$ cannot be abduced anymore to explain $\mathcal{O}_{add(b)}$. Its only minimal explanation is $\mathcal{E}_{\neg \text{cig}(b)}$.

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$$\begin{array}{ll} \text{add}'(X) \leftarrow \text{inex}(X) \wedge \neg ab_{\text{add}'}(X), & \text{add}(X) \leftarrow \neg \text{add}'(X), \\ \text{inex}(X) \leftarrow \text{cig}(X) \wedge \neg ab_{\text{inex}}(X), & ab_{\text{add}'}(X) \leftarrow \text{inspect}(\text{cig}(X)), \\ ab_{\text{add}'}(X) \leftarrow \perp, & ab_{\text{inex}}(X) \leftarrow \perp, \end{array}$$

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Skeptically, we validate **some additive things are not cigarettes**.

Contextual Relevant Consequences

Consider \mathcal{P}_{fire}

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Given \mathcal{O}_{storm} explained by $\mathcal{E}_{lightning}$, $\mathcal{O}_{forest_fire}$'s minimal explanation is

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$\mathcal{O}_{forest_fire}$ is a **contextual relevant consequence** of \mathcal{O}_{storm} given \mathcal{P}_{fire} .