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# **Exterminator: Automatically Correcting Memory Errors with High Probability**

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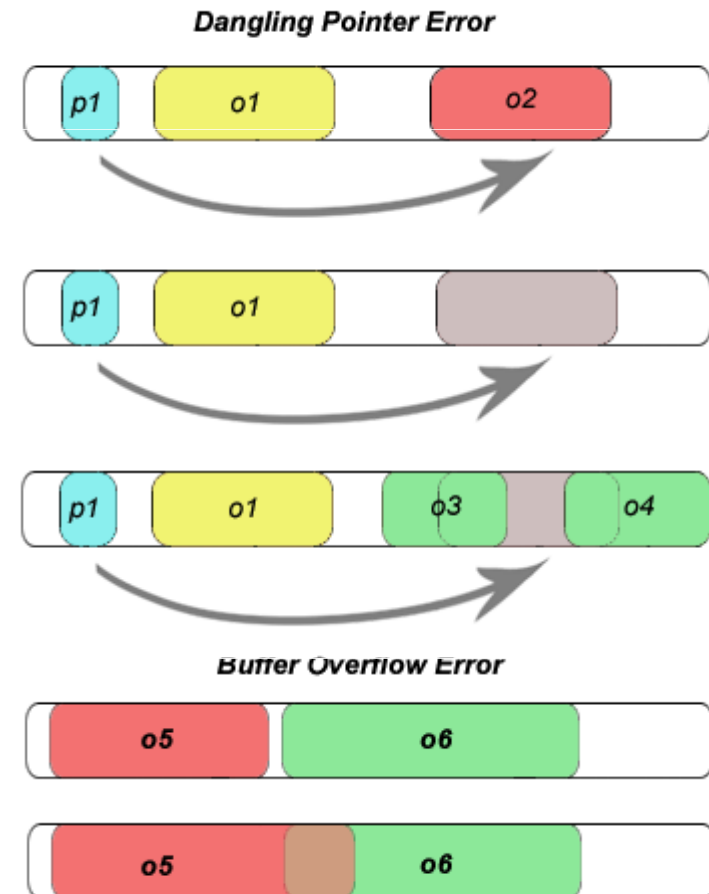
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# Overview

- Manual memory management in C and C++ results in vulnerability to various memory errors.
- In particular **Dangling-Pointers**, **Buffer Overflows**.
- Traditional debugging methods are costly and time consuming.
- Symantec: avg time between discovery and patch deployment for an exploitable bug is 28 days.
- The author presents Exterminator, a runtime system that detects, **isolates** & **corrects** two classes of memory error.

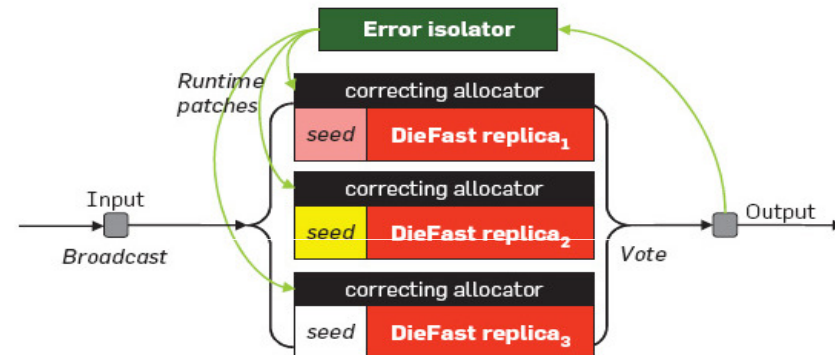
# Target Memory Errors

- Dangling pointers and Buffer-overflow errors are the most commonly **exploited** memory errors.
- Previous works could detect and **tolerate** these. Exterminator aims to **correct** them.



# Modes of Operation

- Iterative mode:
  - **Sequentially** replays the program, each time with the same input but using different random-seeds.
- Replicated mode :
  - Monitors **concurrent** executions with different random-seeds.

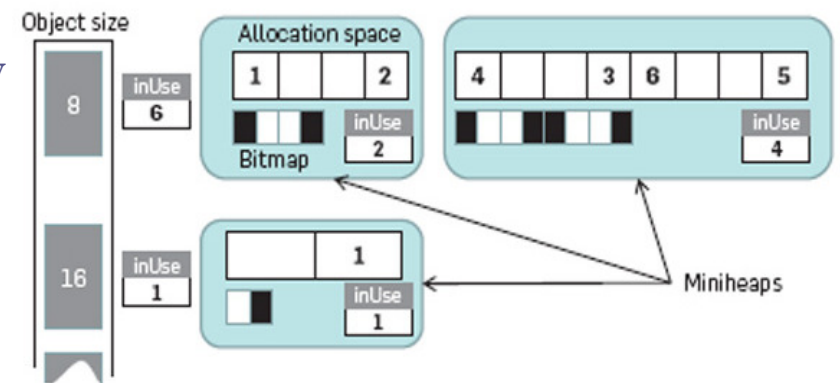


- Cumulative mode:
  - Gradually builds up statistical data **over normal runs** of the program.

# Components of Exterminator

- Exterminator features 3 core components:
  - Probabilistic debugging allocator.
  - Probabilistic error isolation algorithm.
  - Correcting memory allocator.
- The Probabilistic Debugging allocator:
  - Ensures the heap is always  $M$  times larger than max 'needed' ( $1/M$  ratio).
  - Does this using 'mini-heaps' .
  - Randomised Bitmaps tests allow for  $O(1)$  allocation/de-allocation.
  - Ensures:
    - Probable that overflows go into free space.
    - Probable that freed objects aren't immediately reused.

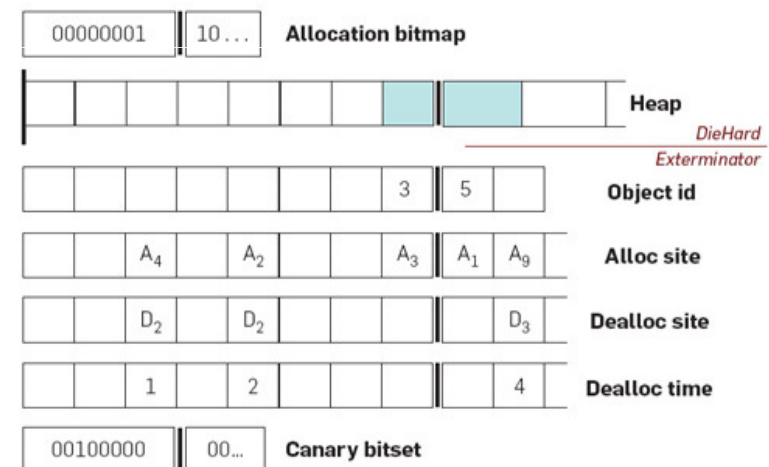
Figure 1: The adaptive (new) DieHard heap layout, used by Exterminator. Objects in the same size class are allocated randomly from separate *miniheaps*, which combined hold  $M$  times more memory than required (here,  $M = 2$ ).



# Heap Architecture

- Object ID's are sequential (an ID of  $n$  indicates the object as the  $n$ th object allocated).
- 'Site' information fields capture the calling contexts.
- Canary bit indicates if the object location was filled with canaries.
- This data persists from allocation, until an object is allocated in its place.
- This adds a memory overhead of 16 bytes and 2 bits per allocated object.

Figure 2: An abstract view of Exterminator's heap layout. Metadata below the horizontal line contains information used for error isolation and correction (see Section 3.2).



# Debugging Allocator

- 32bit randomised canary padding:
  - **Buffer overflows** detected when canary ‘fence-posts’ are overrun.
  - **Dangling pointers** detected when canaries are read, and the program throws an error.
    - May match true data values; in '*cumulative mode*' Exterminator only sets canaries with probability  $P$ .
  - When Exterminator allocates memory it checks canary integrity (based on bit-set).
    - If not intact then it signals an error and isolates this memory.
  - When Exterminator de-allocates memory, it also checks neighbours. If free it performs the same canary-check.



# Error Isolation

## In Iterative and Replicated modes

- **Buffer overflows:**
  - Compare objects across heaps (by Object ID). They should be *identical*.
  - If they are not identical, they are considered a candidate '**victim**' of an overflow.
    - Some exceptions (pointers, handles).
  - For each candidate-victim found Exterminator then examines heap images to find a matching '**culprit**' object.
    - Must be the same distance  $\delta$ , from the victim within each heap-image.
  - Once a culprit has been deduced, overflow size is recorded as max-victim- $\delta$ .
- **Dangling pointers:**
  - Under these modes, only **read & write** dangling pointers are found.
    - Would require >20 replicas for read only dangling pointers.
  - When a **freed** object is overwritten with the same values across all heap-images, classified as a dangling pointer error.

# Error Isolation

## In Cumulative mode

- Approach is to consider exceptional program events (crashes) and generate **statistics**.
- **Dangling pointers:**
  - To force different runs to respond differently, de-allocated objects are only filled with canaries with a **probability  $P$** .
  - The choice of  $P$  reflects a trade-off between dangling pointer isolation, and buffer overflow isolation (converse effects).
- **Buffer overflows:**
  - Identifies corruption by looking for overwritten canary values.
  - For each allocation site, computes probability that the 'culprit' is within.
  - Combines these independent statistics from each run to deduce sites that consistently appear as candidates for housing the culprit.
  - Deduces the culprit from these sites.

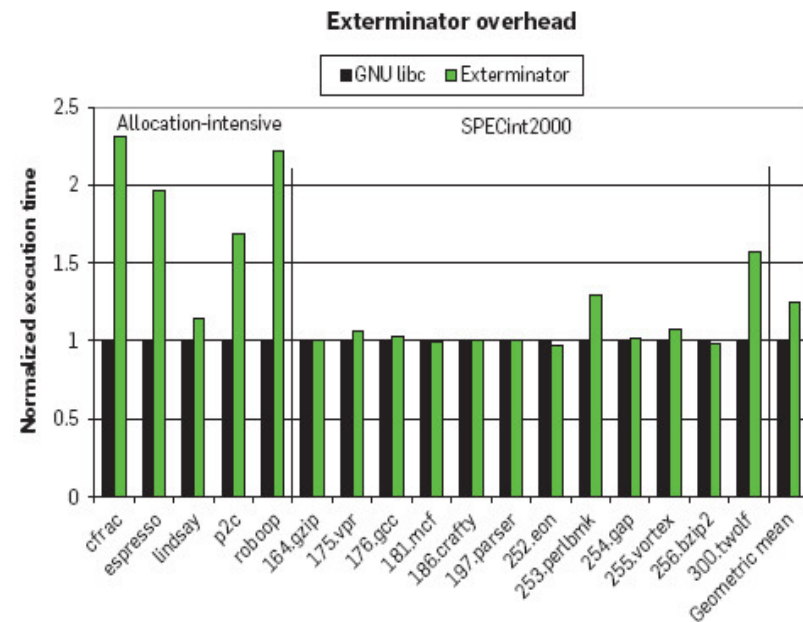
# Error Correction

- Buffer overflow correction
  - Runtime patch contains allocation-site hash and amount of **padding** needed to contain the overflow ( $\delta$  + size of the overflow). If a patch for the allocation site already exists, the Max is used.
- Dangling Pointer Correction
  - Runtime patch contains allocation/de-allocation-site hash and amount of time by which to **delay** de-allocation:
    - Let  $\tau$  be the recorded de-allocation time of this object.  $T$  be the time at which the program crashed or corruption was detected.
    - Delay will be:  $2 \times (T - \tau) + 1$ .

# Results Found

- Exterminator degrades performance by from 0% to 132%.
- Overall geometric mean of 25.1%.
- Most extensive across allocation-intensive suite (81.2% geometric mean).
- Less pronounced across SPECint2000 suite (7.2% g.m.).

Figure 4: Runtime overhead for Exterminator across a suite of benchmarks, normalized to the performance of GNU libc (Linux) allocator.



# Results Found

## Injected Errors

Test Type	Iterations Taken	False-negative rate
Buffer-overflows	3	0%

Test Type	Successful Runs	Iterations Taken
Dangling Ptr. (iterative-mode)	4/10	1
Dangling Ptr. (cumulative-mode)	N/A	22-30

## Real Errors

Software	Problem	Result
Squid Web Cache	Input-dependant overflow Leads to crash	6-byte pad generated Error corrected
Mozilla Firefox	Overflow Cumulative mode only	23 runs with quick-trigger 34 runs with prolonged trigger Error corrected

# Conclusion

- Exterminator automatically detects and **corrects** runtime heap-based memory errors in C/C++ programs, to a high probability.
- Consists of:
  - Probabilistic debugging allocator.
  - Probabilistic error isolation algorithm.
  - Correcting memory allocator.
- Provably low false-positive and false-negative rates.
- Runtime patches are generated that the correcting-memory-allocator corrects memory errors with.
- Exterminator is useful during testing **and** in deployed applications.