

**Assembly Language:
Part 2**

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Goals of this Lecture

Help you learn:

- Intermediate aspects of x86-64 assembly language...
- Control flow with signed integers
- Control flow with unsigned integers
- Arrays
- Structures

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Agenda

Flattened C code

Control flow with signed integers

Control flow with unsigned integers

Arrays

Structures

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Flattened C Code

Problem

- Translating from C to assembly language is difficult when the C code contains **nested** statements

Solution

- **Flatten** the C code to eliminate all nesting

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Flattened C Code

<p>C</p> <pre>if (expr) { statement1; ... statementN; }</pre>	→	<p>Flattened C</p> <pre>if (! expr) goto endif1; statement1; ... statementN; endif1:</pre>
<pre>if (expr) { statementT1; ... statementTN; } else { statementF1; ... statementFN; }</pre>	→	<pre>if (! expr) goto else1; statement1; ... statementN; goto endif1; else: statementF1; ... statementFN; endif1:</pre>

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Flattened C Code

<p>C</p> <pre>while (expr) { statement1; ... statementN; }</pre>	→	<p>Flattened C</p> <pre>loop1: if (! expr) goto endloop1; statement1; ... statementN; goto loop1; endloop1:</pre>
<pre>for (expr1; expr2; expr3) { statement1; ... statementN; }</pre>	→	<pre> expr1; loop1: if (! expr2) goto endloop1; statement1; ... statementN; expr3; goto loop1; endloop1:</pre>

See Bryant & O' Hallaron book for faster patterns

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Agenda

- Flattened C code
- Control flow with signed integers**
- Control flow with unsigned integers
- Arrays
- Structures

if Example

C

```
int i;
...
if (i < 0)
    i = -i;
```

Flattened C

```
int i;
...
    if (i >= 0) goto endif1;
    i = -i;
endif1:
```

if Example

Flattened C

```
int i;
...
    if (i >= 0) goto endif1;
    i = -i;
endif1:
```

Assem Lang

```
.section ".bss"
i: .skip 4
...
.section ".text"
...
    cml $0, i
    jge endif1
    negl i
endif1:
```

Note:
cmp instruction (counterintuitive operand order)
 Sets CC bits in EFLAGS register
jge instruction (conditional jump)
 Examines CC bits in EFLAGS register

if...else Example

C

```
int i;
int j;
int smaller;
...
if (i < j)
    smaller = i;
else
    smaller = j;
```

Flattened C

```
int i;
int j;
int smaller;
...
    if (i >= j) goto else1;
    smaller = i;
    goto endif1;
else1:
    smaller = j;
endif1:
```

if...else Example

Flattened C

```
int i;
int j;
int smaller;
...
    if (i >= j) goto else1;
    smaller = i;
    goto endif1;
else1:
    smaller = j;
endif1:
```

Assem Lang

```
.section ".bss"
i: .skip 4
j: .skip 4
smaller: .skip 4
...
.section ".text"
...
    movl i, %eax
    cml j, %eax
    jge else1
    movl i, %eax
    movl %eax, smaller
    jmp endif1
else1:
    movl j, %eax
    movl %eax, smaller
endif1:
```

Note:
jmp instruction
 (unconditional jump)

while Example

C

```
int fact;
int n;
...
fact = 1;
while (n > 1)
{ fact *= n;
  n--;
}
```

Flattened C

```
int fact;
int n;
...
    fact = 1;
loop1:
    if (n <= 1) goto endloop1;
    fact *= n;
    n--;
    goto loop1;
endloop1:
```

while Example

Flattened C

```
int fact;
int n;
...
fact = 1;
loop1:
if (n <= 1) goto endloop1;
fact *= n;
n--;
goto loop1;
endloop1:
```

Note:
jle instruction (conditional jump)
imul instruction

Assem Lang

```
.section ".bss"
fact: .skip 4
n: .skip 4
...
.section ".text"
...
movl $1, fact
loop1:
cmpl $1, n
jle endloop1
movl fact, %eax
imull n
movl %eax, fact
decl n
jmp loop1
endloop1:
```

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for Example

C

```
int power = 1;
int base;
int exp;
int i;
...
for (i = 0; i < exp; i++)
power *= base;
```

Flattened C

```
int power = 1;
int base;
int exp;
int i;
...
i = 0;
loop1:
if (i >= exp) goto endloop1;
power *= base;
i++;
goto loop1;
endloop1:
```

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for Example

Flattened C

```
int power = 1;
int base;
int exp;
int i;
...
i = 0;
loop1:
if (i >= exp) goto endloop1;
power *= base;
i++;
goto loop1;
endloop1:
```

Assem Lang

```
.section ".data"
power: .long 1
...
.section ".bss"
base: .skip 4
exp: .skip 4
i: .skip 4
...
.section ".text"
...
movl $0, i
loop1:
movl i, %eax
cmpl exp, %eax
jge endloop1
movl power, %eax
imull base
movl %eax, power
incl i
jmp loop1
endloop1:
```

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Control Flow with Signed Integers

Comparing signed integers

```
cmp(q,l,w,b) srcIRM, destRM Compare dest with src
```

- Sets CC bits in the EFLAGS register
- Beware: operands are in counterintuitive order
- Beware: many other instructions set CC bits
 - Conditional jump should **immediately** follow **cmp**

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Control Flow with Signed Integers

Unconditional jump

```
jmp label Jump to label
```

Conditional jumps after comparing signed integers

```
je label Jump to label if equal
jne label Jump to label if not equal
jl label Jump to label if less
jle label Jump to label if less or equal
jg label Jump to label if greater
jge label Jump to label if greater or equal
```

- Examine CC bits in EFLAGS register

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Agenda

- Flattened C
- Control flow with signed integers
- Control flow with unsigned integers**
- Arrays
- Structures

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Signed vs. Unsigned Integers

In C

- Integers are signed or unsigned
- Compiler generates assem lang instructions accordingly

In assembly language

- Integers are neither signed nor unsigned
- Distinction is in the instructions used to manipulate them

Distinction matters for

- Multiplication and division
- Control flow

Handling Unsigned Integers

Multiplication and division

- Signed integers: `imul, idiv`
- Unsigned integers: `mul, div`

Control flow

- Signed integers: `cmp + {je, jne, jl, jle, jg, jge}`
- Unsigned integers: "unsigned cmp" + `{je, jne, jl, jle, jg, jge}` No!!!
- Unsigned integers: `cmp + {je, jne, jb, jbe, ja, jae}`

while Example

C

```
unsigned int fact;
unsigned int n;
...
fact = 1;
while (n > 1)
{ fact *= n;
  n--;
}
```

Flattened C

```
unsigned int fact;
unsigned int n;
...
fact = 1;
loop1:
  if (n <= 1) goto endloop1;
  fact *= n;
  n--;
  goto loop1;
endloop1:
```

while Example

Flattened C

```
unsigned int fact;
unsigned int n;
...
fact = 1;
loop1:
  if (n <= 1) goto endloop1;
  fact *= n;
  n--;
  goto loop1;
endloop1:
```

Assem Lang

```
.section ".bss"
fact: .skip 4
n: .skip 4
...
.section ".text"
...
movl $1, fact
loop1:
  cmpl $1, n
  jbe endloop1
  movl fact, %eax
  mull n
  movl %eax, fact
  decl n
  jmp loop1
endloop1:
```

Note:
jbe instruction (instead of `jle`)
mull instruction (instead of `imull`)

for Example

C

```
unsigned int power = 1;
unsigned int base;
unsigned int exp;
unsigned int i;
...
for (i = 0; i < exp; i++)
  power *= base;
```

Flattened C

```
unsigned int power = 1;
unsigned int base;
unsigned int exp;
unsigned int i;
...
i = 0;
loop1:
  if (i >= exp) goto endloop1;
  power *= base;
  i++;
  goto loop1;
endloop1:
```

for Example

Flattened C

```
unsigned int power = 1;
unsigned int base;
unsigned int exp;
unsigned int i;
...
i = 0;
loop1:
  if (i >= exp) goto endloop1;
  power *= base;
  i++;
  goto loop1;
endloop1:
```

Assem Lang

```
.section ".data"
power: .long 1
.section ".bss"
base: .skip 4
exp: .skip 4
i: .skip 4
...
.section ".text"
...
movl $0, i
loop1:
  movl i, %eax
  cmpl exp, %eax
  jae endloop1
  movl power, %eax
  mull base
  movl %eax, power
  incl i
  jmp loop1
endloop1:
```

Note:
jae instruction (instead of `jge`)
mull instruction (instead of `imull`)

Control Flow with Unsigned Integers

Comparing unsigned integers

- Same as comparing signed integers

Conditional jumps after comparing unsigned integers

```

je label  Jump to label if equal
jne label Jump to label if not equal
jb label  Jump to label if below
jbe label Jump to label if below or equal
ja label  Jump to label if above
jae label Jump to label if above or equal
    
```

- Examine CC bits in EFLAGS register

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Agenda

- Flattened C
- Control flow with signed integers
- Control flow with unsigned integers
- Arrays**
- Structures**

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Arrays: Indirect Addressing

C

```

int a[100];
int i;
int n;
...
i = 3;
n = a[i]
    
```

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
    
```

One step at a time...

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Arrays: Indirect Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
    
```

Registers

RAX

R10

...

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Indirect Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
    
```

Registers

RAX

R10

...

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Indirect Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
    
```

Registers

RAX

R10

...

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Indirect Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 1012

R10:

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Indirect Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 1012

R10: 123

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Note: Indirect addressing

Arrays: Indirect Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movslq i, %rax
salq $2, %rax
addq $a, %rax
movl (%rax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 1012

R10: 123

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n	123	1404

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Arrays: Base+Disp Addressing

C

```

int a[100];
int i;
int n;
...
i = 3;
n = a[i]
                    
```

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
call $2, %eax
movl a(%eax), %r10d
movl %r10d, n
...
                    
```

One step at a time...

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Arrays: Base+Disp Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
call $2, %eax
movl a(%eax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX:

R10:

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Base+Disp Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
call $2, %eax
movl a(%eax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 3

R10:

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Base+Disp Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
sall $2, %eax
movl a(%eax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 12

R10:

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Base+Disp Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
sall $2, %eax
movl a(%eax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 12

R10: 123

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Note:
Base+displacement addressing

Arrays: Base+Disp Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
sall $2, %eax
movl a(%eax), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 12

R10: 123

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n	123	1404

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Arrays: Scaled Indexed Addressing

C

```

int a[100];
int i;
int n;
...
i = 3;
n = a[i]
                    
```

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
movl a(,%eax,4), %r10d
movl %r10d, n
...
                    
```

One step at a time...

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Arrays: Scaled Indexed Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
movl a(,%eax,4), %r10d
movl %r10d, n
...
                    
```

Registers

RAX:

R10:

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Scaled Indexed Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
movl a(,%eax,4), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 3

R10:

Memory

0		1000
1		1004
2		1008
3	123	1012
...		
100		1396
i	3	1400
n		1404

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Arrays: Scaled Indexed Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
movl a(,%eax,4), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 3
R10: 123

Memory

0	1000
1	1004
2	1008
3	1012
...	...
100	1396
i	3
...	...
n	1404

Note: **Scaled indexed** addressing

Arrays: Scaled Indexed Addressing

Assem Lang

```

.section ".bss"
a: .skip 400
i: .skip 4
n: .skip 4
...
.section ".text"
...
movl $3, i
...
movl i, %eax
movl a(,%eax,4), %r10d
movl %r10d, n
...
                    
```

Registers

RAX: 12
R10: 123

Memory

0	1000
1	1004
2	1008
3	1012
...	...
100	1396
i	3
...	...
n	1404

Generalization: Memory Operands

Full form of memory operands:

displacement (base, index, scale)

- **displacement** is an integer or a label (default = 0)
- **base** is a 4-byte or 8-byte register
- **index** is a 4-byte or 8-byte register
- **scale** is 1, 2, 4, or 8 (default = 1)

Meaning

- Compute the sum (displacement) + (contents of base) + ((contents of index) * (scale))
- Consider the sum to be an address
- Load from (or store to) that address

Note:

- All other forms are subsets of the full form...

Generalization: Memory Operands

Valid subsets:

- **Direct addressing**
 - displacement
- **Indirect addressing**
 - (base)
- **Base+displacement addressing**
 - displacement (base)
- **Indexed addressing**
 - (base, index)
 - displacement (base, index)
- **Scaled indexed addressing**
 - (, index, scale)
 - displacement (, index, scale)
 - (base, index, scale)
 - displacement (base, index, scale)

Operand Examples

Immediate operands

- \$5 => use the number 5 (i.e. the number that is available immediately within the instruction)
- \$i => use the address denoted by i (i.e. the address that is available immediately within the instruction)

Register operands

- %rax => read from (or write to) register RAX

Memory operands: direct addressing

- 5 => load from (or store to) memory at address 5 (silly; seg fault)
- i => load from (or store to) memory at the address denoted by i

Memory operands: indirect addressing

- (%rax) => consider the contents of RAX to be an address; load from (or store to) that address

Operand Examples

Memory operands: base+displacement addressing

- 5(%rax) => compute the sum (5) + (contents of RAX); consider the sum to be an address; load from (or store to) that address
- i(%rax) => compute the sum (address denoted by i) + (contents of RAX); consider the sum to be an address; load from (or store to) that address

Memory operands: indexed addressing

- 5(%rax,%r10) => compute the sum (5) + (contents of RAX) + (contents of R10); consider the sum to be an address; load from (or store to) that address
- i(%rax,%r10) => compute the sum (address denoted by i) + (contents of RAX) + (contents of R10); consider the sum to be an address; load from (or store to) that address

Operand Examples

Memory operands: scaled indexed addressing

- `5(%rax,%r10,4)` => compute the sum (5) + (contents of RAX) + ((contents of R10) * 4); consider the sum to be an address; load from (or store to) that address
- `i(%rax,%r10,4)` => compute the sum (address denoted by i) + (contents of RAX) + ((contents of R10) * 4); consider the sum to be an address; load from (or store to) that address

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Aside: The lea Instruction

lea: load effective address

- Unique instruction: suppresses memory load/store

Example

- `movq 5(%rax), %r10`
 - Compute the sum (5) + (contents of RAX); consider the sum to be an address; load 8 bytes from that address into R10
- `leaq 5(%rax), %r10`
 - Compute the sum (5) + (contents of RAX); move that sum to R10

Useful for

- Computing an address, e.g. as a function argument
- See precept code that calls `scanf()`
- Some quick-and-dirty arithmetic

What is the effect of this?
`leaq (%rax,%rax,4), %rax`

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Agenda

- Flattened C
- Control flow with signed integers
- Control flow with unsigned integers
- Arrays
- Structures**

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Structures: Indirect Addressing

<p>C</p> <pre>struct S { int i; int j; }; ... struct S myStruct; ... myStruct.i = 18; ... myStruct.j = 19;</pre>	<p>Assem Lang</p> <pre>.section ".bss" myStruct: .skip 8section ".text" ... movq %myStruct, %rax movl \$18, (%rax) ... movq %myStruct, %rax addq \$4, %rax movl \$19, (%rax)</pre>
---	---

Note:
Indirect addressing

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Structures: Base+Disp Addressing

<p>C</p> <pre>struct S { int i; int j; }; ... struct S myStruct; ... myStruct.i = 18; ... myStruct.j = 19;</pre>	<p>Assem Lang</p> <pre>.section ".bss" myStruct: .skip 8section ".text" ... movl \$0, %eax movl \$18, myStruct(%eax) ... movl \$4, %eax movl \$19, myStruct(%eax)</pre>
---	--

Note:
Base+displacement addressing

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Structures: Padding

<p>C</p> <pre>struct S { char c; int i; }; ... struct S myStruct; ... myStruct.c = 'A'; ... myStruct.i = 18;</pre>	<p>Assem Lang</p> <pre>.section ".bss" myStruct: .skip 8section ".text" ... movl \$0, %eax movb \$'A', myStruct(%eax) ... movl \$4, %eax movl \$18, myStruct(%eax)</pre>
---	---

Three-byte pad here

Beware:
Compiler sometimes inserts padding after fields

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Structures: Padding

x86-64/Linux rules

Data type	Within a struct, must begin at address that is evenly divisible by:
(unsigned) char	1
(unsigned) short	2
(unsigned) int	4
(unsigned) long	8
float	4
double	8
long double	16
any pointer	8

- Compiler may add padding after last field if struct is within an array

Summary

Intermediate aspects of x86-64 assembly language...

Flattened C code

Control transfer with signed integers

Control transfer with unsigned integers

Arrays

- Full form of instruction operands

Structures

- Padding

Appendix

Setting and using CC bits in EFLAGS register

Setting Condition Code Bits

Question

- How does `cmp1` set condition code bits in EFLAGS register?

Answer

- (See following slides)

Condition Code Bits

Condition code bits

- **ZF: zero** flag: set to 1 iff result is **zero**
- **SF: sign** flag: set to 1 iff result is **negative**
- **CF: carry** flag: set to 1 iff **unsigned overflow** occurred
- **OF: overflow** flag: set to 1 iff **signed overflow** occurred

Condition Code Bits

Example: `addl src, dest`

- Compute sum (`dest+src`)
- Assign sum to `dest`
- ZF: set to 1 iff `sum == 0`
- SF: set to 1 iff `sum < 0`
- CF: set to 1 iff unsigned overflow
 - Set to 1 iff `sum < src`
- OF: set if signed overflow
 - Set to 1 iff `(src > 0 && dest > 0 && sum < 0) || (src < 0 && dest < 0 && sum >= 0)`

Condition Code Bits

Example: `subl src, dest`

- Compute sum (`dest+(-src)`)
- Assign sum to `dest`
- ZF: set to 1 iff sum == 0
- SF: set to 1 iff sum < 0
- CF: set to 1 iff unsigned overflow
 - Set to 1 iff `dest < src`
- OF: set to 1 iff signed overflow
 - Set to 1 iff `(dest > 0 && src < 0 && sum < 0) || (dest < 0 && src > 0 && sum >= 0)`

Example: `cml src, dest`

- Same as `subl`
- But does not affect `dest`

Using Condition Code Bits

Question

- How do conditional jump instructions use condition code bits in EFLAGS register?

Answer

- (See following slides)

Conditional Jumps: Unsigned

After comparing unsigned data

Jump Instruction	Use of CC Bits
<code>je label</code>	ZF
<code>jne label</code>	~ZF
<code>jb label</code>	CF
<code>jae label</code>	~CF
<code>jbe label</code>	CF ZF
<code>ja label</code>	~(CF ZF)

Note:

- If you can understand why `jb` jumps iff CF
- ... then the others follow

Conditional Jumps: Unsigned

Why does `jb` jump iff CF? Informal explanation:

(1) `largenum – smallnum (not below)`

- Correct result
- => CF=0 => don't jump

(2) `smallnum – largenum (below)`

- Incorrect result
- => CF=1 => jump

Conditional Jumps: Signed

After comparing signed data

Jump Instruction	Use of CC Bits
<code>je label</code>	ZF
<code>jne label</code>	~ZF
<code>jl label</code>	OF ^ SF
<code>jge label</code>	~(OF ^ SF)
<code>jle label</code>	(OF ^ SF) ZF
<code>jg label</code>	~((OF ^ SF) ZF)

Note:

- If you can understand why `j1` jumps iff `OF^SF`
- ... then the others follow

Conditional Jumps: Signed

Why does `jl` jump iff `OF^SF`? Informal explanation:

(1) `largeposnum – smallposnum (not less than)`

- Certainly correct result
- => OF=0, SF=0, `OF^SF`=0 => don't jump

(2) `smallposnum – largeposnum (less than)`

- Certainly correct result
- => OF=0, SF=1, `OF^SF`=1 => jump

(3) `largenegnum – smallnegnum (less than)`

- Certainly correct result
- => OF=0, SF=1 => `(OF^SF)`=1 => jump

(4) `smallnegnum – largenegnum (not less than)`

- Certainly correct result
- => OF=0, SF=0 => `(OF^SF)`=0 => don't jump

Conditional Jumps: Signed



- (5) **posnum – negnum (not less than)**
 - Suppose correct result
 - => $OF=0, SF=0 \Rightarrow (OF \wedge SF)=0 \Rightarrow$ don't jump
- (6) **posnum – negnum (not less than)**
 - Suppose incorrect result
 - => $OF=1, SF=1 \Rightarrow (OF \wedge SF)=0 \Rightarrow$ don't jump
- (7) **negnum – posnum (less than)**
 - Suppose correct result
 - => $OF=0, SF=1 \Rightarrow (OF \wedge SF)=1 \Rightarrow$ jump
- (8) **negnum – posnum (less than)**
 - Suppose incorrect result
 - => $OF=1, SF=0 \Rightarrow (OF \wedge SF)=1 \Rightarrow$ jump