

# Analyzing fuzzy and contextual approaches to vagueness by semantic games

PhD Thesis

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# Motivation

## Vagueness

- ▶ ubiquitous in **all** natural languages
- ▶ still neglected by most natural language processing systems

Vagueness is modeled by **logicians**, **philosophers**, **linguists**, etc., but

- ▶ ... different aims and methodologies
- ▶ ... no systematic comparison of different approaches
- ▶ ... lack of interpretations

# Outline

## What is Vagueness

### Fuzzy approaches

- Giles's Game

- Generalizing Giles's Game

- Fuzzy Quantification

### Barker's 'Dynamics of Vagueness'

- Bridges to fuzzy logics by measuring contexts

- Saturated Contexts

### Shapiro's 'Vagueness in Context'

- Semantic games for Shapiro's logic

### Relating Barker's and Shapiro's approaches

# What is Vagueness?

## Related phenomena

'Vagueness' in colloquial language may be used for:

### Underspecificity / Unwanted generality

*Someone said something.*

### Ambiguity

*I'm at the bank.*

### 'Genuine' vagueness

*He is rather tall.*

### Vagueness due to imprecision

*We'll meet at 12 o'clock.*

# What is Vagueness?

## Characteristics of Vagueness

### Borderline cases



*red*



???



*not red*

Is the patch in the middle considered *red* or *not red*?

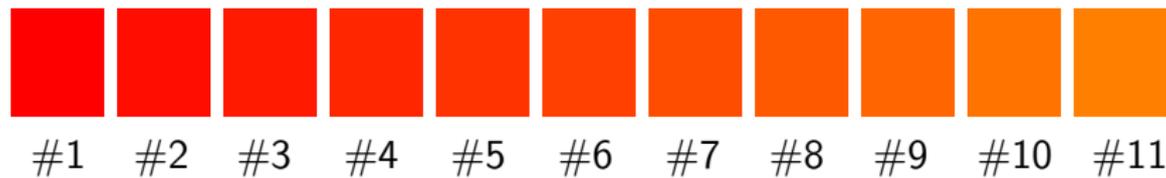
### Blurry boundaries



Where do the red patches and the not-red patches start?

# What is Vagueness?

## Characteristics of Vagueness (contd.)



### Sorites paradox

(Premise 1) Patch #1 is red.

(Premise 2) If a patch is red and you change it just a bit, it remains red.

(Conclusion) Patch #11 is red.

- ▶ Sorites series can be formed for all vague predicates

# Fuzzy approaches to vagueness

## Fuzzy approaches

- ▶ “*Truth comes in degrees.*”: unit interval  $[0, 1]$
- ▶ truth-functional
- ▶ based on continuous *t*-norms

## Interpretation of fuzzy models

- ▶ dialogue semantics (Giles 1970)
- ▶ characterize truth by evaluation games

# Giles's game

## Overview

### Giles's game for Łukasiewicz logic $\mathcal{L}_\infty$

- ▶ evaluation game: determine **truth** in a model
  - ▶ each player asserts a **multiset** of formulas
  - ▶ initially, **I** assert the formula in question
  - ▶ **dialogue part** and **betting part**
- 
- ▶ can be motivated as a generalization of Hintikka's evaluation game for classical logic

# Giles's game

## Dialogue part

- ▶ Decompose formulas according to dialogue rules

$(R_{\rightarrow})$ : If I assert  $F \rightarrow G$  then you may attack by asserting  $F$ , which obliges me to defend by asserting  $G$ .

$(R_{\forall})$ : If I assert  $\forall x.F(x)$ , then you may attack by choosing  $c$ , which obliges me to assert  $F(c)$ .

$(LLA)$ : I can declare not to attack an assertion.

$(LLD)$ : I can just assert  $\perp$  in reply to your attack.

- ▶ rules for your assertions are dual
- ▶ no further regulations

# Giles's game

Betting part and correspondence to  $\mathbb{L}_\infty$

- ▶ dispersive binary experiments
- ▶ success probabilities correspond to  $v_M$
- ▶ experiment fails: the asserting player pays 1€
- ▶ **Risk**: the **expected** amount of money I have to pay

## Theorem (Giles)

The least final risk  $r$  I can enforce directly corresponds to  $v_M(F)$ :

- ▶ I have a strategy, that **my** risk is not greater than  $r$
- ▶ **you** have a strategy, that **my** risk is not smaller than  $r$

# Generalizing Giles's game

## Motivation

### Generalize Giles's game

Payoff function: take **some** payoff function

Truth values: transformations between **payoff** and **truth** values

Dialogue rules: define a **general format** of dialogue rules

- ▶ conditions on payoff functions and dialogue rules
- ▶ derive truth functions from dialogue rules

# Generalizing Giles's game

## Results

### Theorem

Let  $\mathcal{D}$  be a game with discriminating payoff function and decomposing dual rules. Then for each connective  $\diamond$  there is a function  $f_\diamond$  such that  $\langle \mid \diamond(F_1, \dots, F_n) \rangle = f_\diamond(\langle \mid F_1 \rangle, \dots, \langle \mid F_n \rangle)$  for all formulas  $F_1, \dots, F_n$ .

Logics which can be characterized in this framework:

- ▶ Łukasiewicz logic  $\mathbf{L}_\infty$ ,  $\mathbf{L}_n$ , Abelian logic  $\mathbf{A}$ ,  $\mathbf{CHL}$
- ▶  $f_\diamond$  composed by min, max, and  $\oplus$

# Fuzzy Quantification

## Motivation

- ▶ model vague **natural language quantifiers** like *about half, nearly a third*
- ▶ focus on **semi-fuzzy** quantifiers

### Random witness selection

If I assert  $\Pi x. \hat{F}(x)$  then I have to assert  $\hat{F}(c)$  for a **randomly** picked  $c$ .

- ▶ **uniform** distribution over a **finite** domain  $\mathcal{D}$

# Modeling fuzzy quantification

Giles's game with random witness selection

- ▶ random witness selection
- ▶ blind choice and deliberate choice

## Blind choice quantifiers

- ▶ Place bets **for** and **against** a fixed number of instances.
- ▶ Identities of random instances are revealed only afterwards.

# Modeling fuzzy quantification

## Blind choice quantifiers - Results

### Theorem

All blind choice quantifiers can be expressed using quantifiers of the form  $L_m^k$ ,  $G_m^k$ , disjunction  $\vee$ , conjunction  $\wedge$ , and  $\perp$ .

### Theorem

The blind choice quantifiers  $G_m^k$  and  $L_m^k$  can be expressed in  $\mathbf{L}_\infty(\Pi)$  via the following reductions:

$$G_m^k x. \hat{F}(x) \equiv \neg((\neg \Pi x. \hat{F}(x))^{m+1}) \ \& \ (\Pi x. \hat{F}(x))^{k-1}$$

$$L_m^k x. \hat{F}(x) \equiv \neg((\Pi x. \hat{F}(x))^{k+1}) \ \& \ (\Pi x. \neg \hat{F}(x))^{m-1}$$

### Corollary

All blind choice quantifiers can be expressed in  $\mathbf{L}_\infty(\Pi)$ .

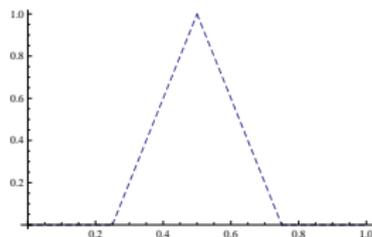
# Modeling fuzzy quantification

## Blind choice quantifiers - Example

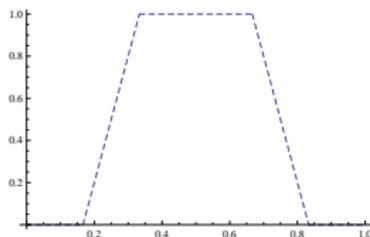
### Example - Around half

Consider the quantifier  $H_t^s x. \hat{F}(x) =_{df} G_{s-t}^{s+t} x. \hat{F}(x) \wedge L_{s-t}^{s+t} x. \hat{F}(x)$

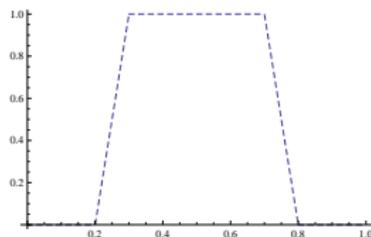
- ▶  $s$ : determines sample size
- ▶  $t$ : determines tolerance



$$H_0^2 x. \hat{F}(x)$$



$$H_1^3 x. \hat{F}(x)$$



$$H_2^5 x. \hat{F}(x)$$

# Modeling fuzzy quantification

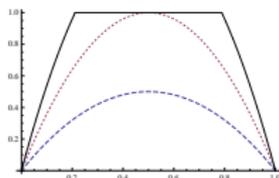
## Deliberate choice quantifiers

- $(R_{\Pi_m^k})$ :
- ▶ choose  $k + m$  constants randomly
  - ▶ defender bets for  $k$  and against  $m$  constants
  - ▶ may also invoke LLD instead

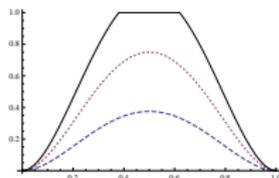
## Proposition

The dialogue rule  $R_{\Pi_m^k}$  matches the specification

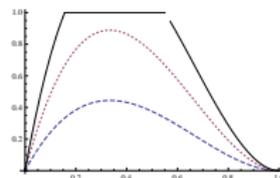
$$v_M(\Pi_m^k \hat{F}(x)) = \binom{k+m}{k} (Prop_x \hat{F}(x))^k (1 - Prop_x \hat{F}(x))^m$$



$$W_i(\Pi_1^1)_x \hat{F}(x)$$



$$W_i(\Pi_2^2)_x \hat{F}(x)$$



$$W_i(\Pi_2^1)_x \hat{F}(x)$$

- ▶ operator  $W_i$  to project towards 1

# Barker's 'Dynamics of Vagueness'

## Overview

- ▶ linguistic **dynamic**, **scale-based** approach to vagueness

### The 'Dynamics of vagueness'

- ▶ objects possess vague / gradable properties **to some degree**
- ▶ context is defined as a set of **classical worlds**
- ▶ **meaning** defined as **context change potential**

# Barker's 'Dynamics of Vagueness'

## Example Context

### Example context $C$

$w$	$\delta(w)(\uparrow tall)$	$tall(w)(b)$	$\delta(w)(\uparrow clever)$	$clever(w)(b)$
$w_1$	170	175	100	105
$w_2$	160	170	120	125
$w_3$	170	180	100	95
$w_4$	180	175	105	100
$w_5$	170	165	110	115

- ▶ values for *tallness*, *cleverness*, Bill's height and cleverness

# Barker's 'Dynamics of Vagueness'

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context update: filter out worlds

- ▶  $\llbracket Bill\ is\ tall \rrbracket(C) = \{w_1, w_2, w_3\}$

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context update: filter out worlds

- ▶  $\llbracket \text{Bill is tall} \rrbracket(C) = \{w_1, w_2, w_3\}$
- ▶  $\llbracket \text{Bill is clever} \rrbracket(C) = \{w_1, w_2, w_5\}$

# Barker's 'Dynamics of Vagueness'

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context update: filter out worlds

- ▶  $\llbracket Bill\ is\ tall \rrbracket(C) = \{w_1, w_2, w_3\}$
- ▶  $\llbracket Bill\ is\ clever \rrbracket(C) = \{w_1, w_2, w_5\}$
- ▶  $\llbracket Bill\ is\ tall\ and\ Bill\ is\ clever \rrbracket(C) = \{w_1, w_2\}$

# Barker's 'Dynamics of Vagueness'

## Changes to Barker's formalism

Barker: several issues and shortcomings

- ▶ simplify notation (**linear scales**)
- ▶ **compose** predicate modifiers:  
*definitely very tall, very very tall*
- ▶ introduce logical operators
  - ▶ **disjunction, conjunction, negation** via **union, intersection, and complement** of contexts.
  - ▶ **conditional** via **material implication**

# Barker's 'Dynamics of Vagueness'

## Measuring contexts

Find connections to fuzzy logics

- ▶ determine proportion of worlds filtered out

## Example context $C$

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- ▶  $\|Bill\ is\ tall\ \|_C = 3/5$

# Barker's 'Dynamics of Vagueness'

## Measuring contexts

Find connections to fuzzy logics

- ▶ determine proportion of worlds filtered out

## Example context $C$

$w$	$\delta(w)(\uparrow tall)$	$tall(w)(b)$	$\delta(w)(\uparrow clever)$	$clever(w)(b)$
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$w_5$	170	165	110	115

- ▶  $\|Bill\ is\ tall\|_C = 3/5$
- ▶  $\|Bill\ is\ tall\ and\ Bill\ is\ clever\|_C = 2/5$

# Barker's 'Dynamics of Vagueness'

## Measuring contexts

- ▶ **not** truth-functional

### Theorem

The following boundaries are tight:

- ▶  $*_{\mathbf{L}}(\|\phi\|_c, \|\psi\|_c) \leq \|\phi \wedge \psi\|_c \leq *_{\mathbf{G}}(\|\phi\|_c, \|\psi\|_c)$ ,
- ▶  $\bar{*}_{\mathbf{G}}(\|\phi\|_c, \|\psi\|_c) \leq \|\phi \vee \psi\|_c \leq \bar{*}_{\mathbf{L}}(\|\phi\|_c, \|\psi\|_c)$ , and
- ▶  $I_{\mathbf{G}}(\|\phi\|_c, \|\psi\|_c) \leq \|\phi \rightarrow \psi\|_c \leq \|\phi\|_c \Rightarrow_{\mathbf{L}} \|\psi\|_c$ .

- ▶ *t*-norms, co-*t*-norms, etc., occur as bounds

# Barker's 'Dynamics of Vagueness'

## Saturated contexts

### Saturated context

- ▶ predicates and individuals are **independent** of each other
- ▶ contexts are **'dense'**
- ▶ contexts determined by lower and upper bounds for values
- ▶ fuzzy truth values for atomic propositions directly computable

For  $\phi$  and  $\psi$  with disjoint predicate symbols:

$$\|\phi \wedge \psi\|_c = \|\phi\|_c *_{\cap} \|\psi\|_c$$

$$\|\phi \vee \psi\|_c = \|\phi\|_c \bar{*}_{\cap} \|\psi\|_c$$

# Shapiro's 'Vagueness in Context'

## Motivation

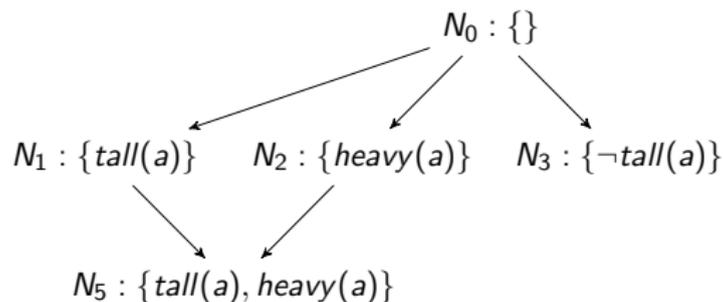
- ▶ philosophical **contextual**, **delineation-based** approach

## 'Vagueness in Context'

- ▶ contexts are **partial valuations**
  - ▶ fixed context space: all **admissible** contexts
  - ▶ define **local** and **global** connectives and quantifiers
- 
- ▶ local connectives: **strong Kleene** tables

# Shapiro's 'Vagueness in Context'

## Contexts



- ▶ directed acyclic structures with root  $N_0$
- ▶ ordered by  $\succeq$  'sharpens'
- ▶ truth defined in terms of forcing

# Shapiro's 'Vagueness in Context'

## Evaluation games

a Hintikka-style evaluation game for Shapiro's logic:

- ▶ players refer to truth values

a Giles-style evaluation game:

- ▶ three-player game between **you**, **me**, and **nature**
- ▶ assertions relative to sharpenings
- ▶ modifiers '**maybe**' and '**surely**'

# Shapiro's 'Vagueness in Context'

## Evaluation games

a Hintikka-style evaluation game for Shapiro's logic:

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a Giles-style evaluation game:

- ▶ three-player game between **you**, **me**, and **nature**
- ▶ assertions relative to sharpenings
- ▶ modifiers '**maybe**' and '**surely**'

## Theorem

For the game  $\mathcal{G} = [ | F_P ]$  starting with my assertion of  $F$  at  $P$

- (i) I have a winning strategy for  $\mathcal{G}$  iff  $F$  is **true**,
- (ii) you have a winning strategy for  $\mathcal{G}$  iff  $F$  is **false**, and
- (iii) neither of us has a winning strategy for  $\mathcal{G}$  iff  $F$  is **indefinite** at the world  $P$ .

# Relating Barker's and Shapiro's approaches

## Motivation

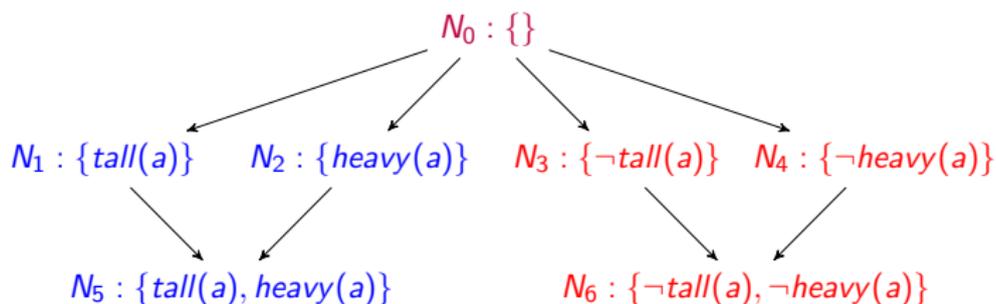
- ▶ different motivations and formalisms

## Questions

- ▶ compare **possible inferences** using Barker / Shapiro
- ▶ compare **expressiveness** of approaches
  
- ▶ Shapiro: handle **conflicting** information
- ▶ Barker: all worlds are **classical**
  - ▶ require **completeness** for Shapiro's contexts

# Relating Barker's and Shapiro's approaches

Partitioning the context space - Shapiro



- ▶ partition possible worlds by their **completions**
- ▶ same propositions are **forced** within a partition block

$${}^W T_{sh}(N) =_{df} \{N' \in W : N' \succeq N, N' \text{ is complete}\}$$

# Relating Barker's and Shapiro's approaches

Partitioning the context space - Barker

$w$	$\delta(w)(\uparrow tall)$	$tall(w)(a)$	$\delta(w)(\uparrow heavy)$	$heavy(w)(a)$
$w_1$	170	175	100	105
$w_2$	165	170	90	125
$w_3$	170	177	95	100
$w_4$	180	175	90	80
$w_5$	170	165	105	100

- ▶ partition **sets** of possible worlds
- ▶ three **partitions**:
  - ▶  $C \subseteq \{w_1, w_2, w_3\}$
  - ▶  $C \subseteq \{w_4, w_5\}$
  - ▶  $C \cap \{w_1, w_2, w_3\} \neq \emptyset$  and  $C \cap \{w_4, w_5\} \neq \emptyset$

# Relating Barker's and Shapiro's approaches

## Partitioning the context space - Barker

$w$	$\delta(w)(\uparrow tall)$	$tall(w)(a)$	$\delta(w)(\uparrow heavy)$	$heavy(w)(a)$
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▶ partition **sets** of possible worlds

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$T_b(C) =_{df} \{s(w) : w \in C\}$  with  $s(w)$  classical s.t.

$R(u)$  is true at  $s(w)$  iff  $w \in \llbracket R \rrbracket(u)(C)$  and

$R(u)$  is false at  $s(w)$  iff  $w \notin \llbracket R \rrbracket(u)(C)$

# Relating Barker's and Shapiro's approaches

## Corresponding models

### Definition

$\langle W, N_0 \rangle_{Sh}$  ... Shapiro model

$\langle \mathcal{P}(C_0), C_0 \rangle_B$  ... Barker model

Corresponding models iff:

- ▶ agree on relevant predicates and individuals
- ▶  $N \in W$ : there exists  $C \in \mathcal{P}(C_0)$  such that  ${}^W T_{Sh}(N) = T_b(C)$
- ▶  $C \in \mathcal{P}(C_0)$ : there exists  $N \in W$  such that  $T_b(C) = {}^W T_{Sh}(N)$

# Relating Barker's and Shapiro's approaches

## Corresponding models

### Definition

$\langle W, N_0 \rangle_{Sh}$  ... Shapiro model

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Corresponding models iff:

- ▶ agree on relevant predicates and individuals
- ▶  $N \in W$ : there exists  $C \in \mathcal{P}(C_0)$  such that  ${}^W T_{Sh}(N) = T_b(C)$
- ▶  $C \in \mathcal{P}(C_0)$ : there exists  $N \in W$  such that  $T_b(C) = {}^W T_{Sh}(N)$

### Proposition

$\langle W, N_0 \rangle_{Sh}$  and  $\langle \mathcal{P}(C_0), C_0 \rangle_B$  corresponding models:

There exists a **bijection** between  $[\cdot]_{Sh}^W$  and  $[\cdot]_B^{C_0}$  induced by

$${}_{C_0} T_b^{-1} \circ {}^W T_{Sh} \quad \text{and} \quad {}^W T_{Sh}^{-1} \circ T_b$$

## Theorem

$\langle W, N_0 \rangle_{Sh}$  and  $\langle \mathcal{P}(C_0), C_0 \rangle_B$  corresponding models,  
 $\phi$  unsettled in  $N_0$  and in  $C_0$

- ▶ perform **context update** with  $\phi$
  - ▶ resulting contexts:  $N_1$  and  $C_1$
  - ▶  $[N_1]_{Sh}^W \simeq [C_1]_B^{C_0}$  **holds**
- 
- ▶  $N_1$  and  $C_1$ : **satisfy same first-order formulas**
  - ▶ can be iterated

# Conclusion

## Summary

- ▶ **generalization** of Giles's game to a more abstract framework
- ▶ random witness selection for **semi-fuzzy quantification**
- ▶ recover **(co-)t-norms** within Barker's approach to vagueness
- ▶ present **evaluation games** for Shapiro's logic
- ▶ find **connections** between Barker's and Shapiro's approaches

## Outlook

- ▶ logically model conclusions drawn due to **scale structure**
- ▶ combine Kyburg and Moreau's **contextual** approach with **degree-based** decision making

# Publications



C.G. Fermüller and C. Roschger.

From games to truth functions: A generalization of Giles's game.

*Studia Logica, Special issue on Logic and Games, 2014.*



C.G. Fermüller and C. Roschger.

Randomized game semantics for semi-fuzzy quantifiers.

*Logic Journal of the IGPL, 2014.*



C.G. Fermüller and C. Roschger.

Bridges between contextual linguistic models of vagueness and t-norm based fuzzy logic.

*Petr Hájek on Mathematical Fuzzy Logic, 2015.*



C. Roschger.

Evaluation games for Shapiro's logic of vagueness in context.

*The Logica Yearbook 2009, 2010.*



C. Roschger.

Comparing context updates in delineation and scale based models of vagueness.

*Understanding Vagueness - Logical, philosophical, and linguistic perspectives, 2011.*