

A Novel Control Strategy for Fail-Safe Cyclic Data Exchange in Wireless Sensor Networks

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Wireless Sensor Networks Fundamentals

- ✓ Sensors are cheap, small and with **little processing capacity**
- ✓ Sensors are usually **battery-powered** or **self-powered**
- ✓ A Wireless Sensor Network typically consists of **several hundreds** of sensors
- ✓ In a multi-hop network, **key nodes (gateways)** are used to route data from one cluster to another

Wireless Sensor Networks Fundamentals

- ✓ Administering sensors' energy enhances their **lifetime** and thus **network availability**
- ✓ The underlying problem is **mathematically complex** and requires proper modeling of **physical and application layers**
- ✓ Wireless Sensor Networks are a “**hot topic**” in both academia and industry
- ✓ “**Fail-safe**” is a hot topic in Wireless Sensor Networks

In this scenario, our goal is ...

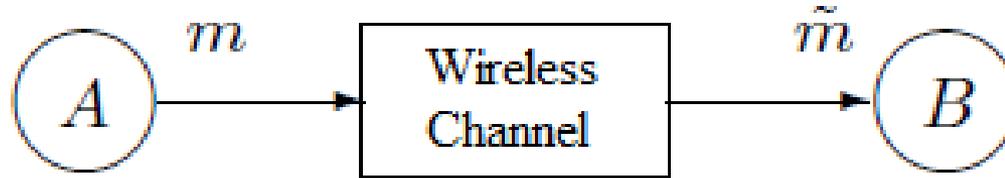
✓ to propose a **mathematical model** that enables the designer to ensure a **target network availability** by trading off between **message transmission rate** and **outage probability**

- **Outage Probability (P_{out})**: probability that the **signal quality** received by a given node **falls under a minimum acceptable threshold value (i.e., bad signal quality)**

Definitions

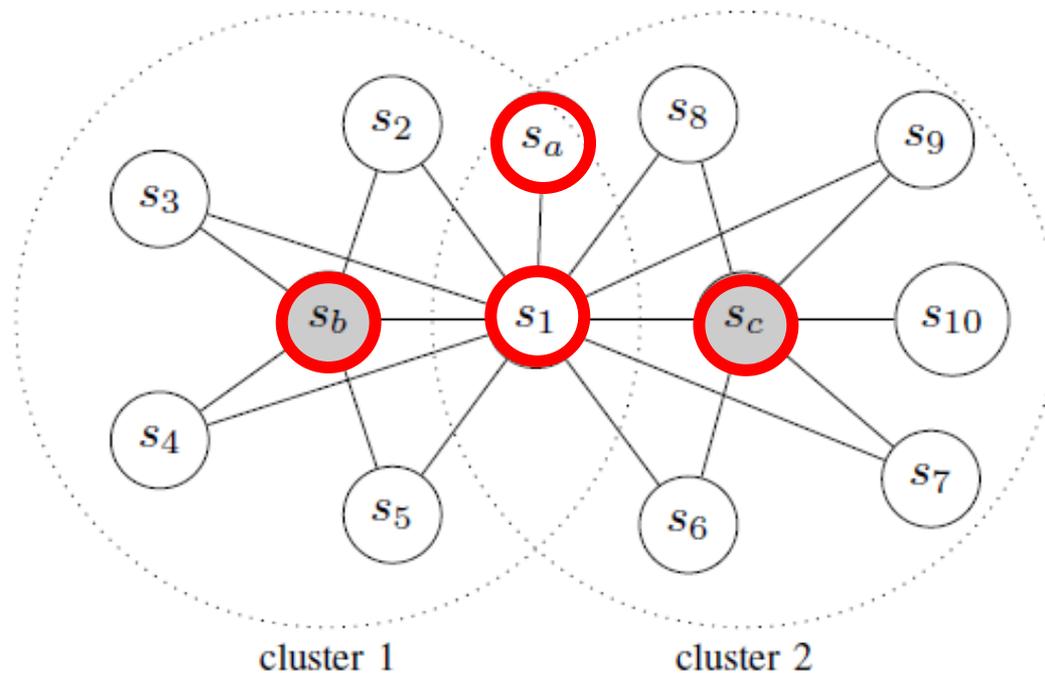
One-Way Message Exchange Mechanism:

Node A transmits m messages to node B through a wireless fading channel. Node B successfully receives \tilde{m} messages.



- Under small-scale fading effects, the system's performance is governed by the signal-to-noise ratio.

System Model



- s_1 and s_a are **Cluster Head Nodes** between **cluster 1** and **cluster 2**
- s_b and s_c are **Coordinator Nodes** control the quality of the communication links between nodes within a single broadcast domain, by overhearing messages during an observation **Window T_w**
- Remaining nodes are **Application Nodes** (e.g., sensing nodes)
- Nodes transmit data at **constant transmission rate R**
- Nodes are synchronized and use **TDMA** to avoid medium contention

Energy Function and Metrics

- Outage Probability (P_{out}):

of successfully received messages by the Coordinator Node from a Sensing Node over a T_w

$$\hat{P}_{out} = 1 - \frac{\tilde{M}}{\hat{m}}$$

of estimated transmitted messages from the sensing node



Receiver Sensitivity

$$\approx \frac{\gamma_0 \sigma_N^2}{S}$$

Channel Noise

Transmit Power

- Where:

Transmission Window

Transmission Rate

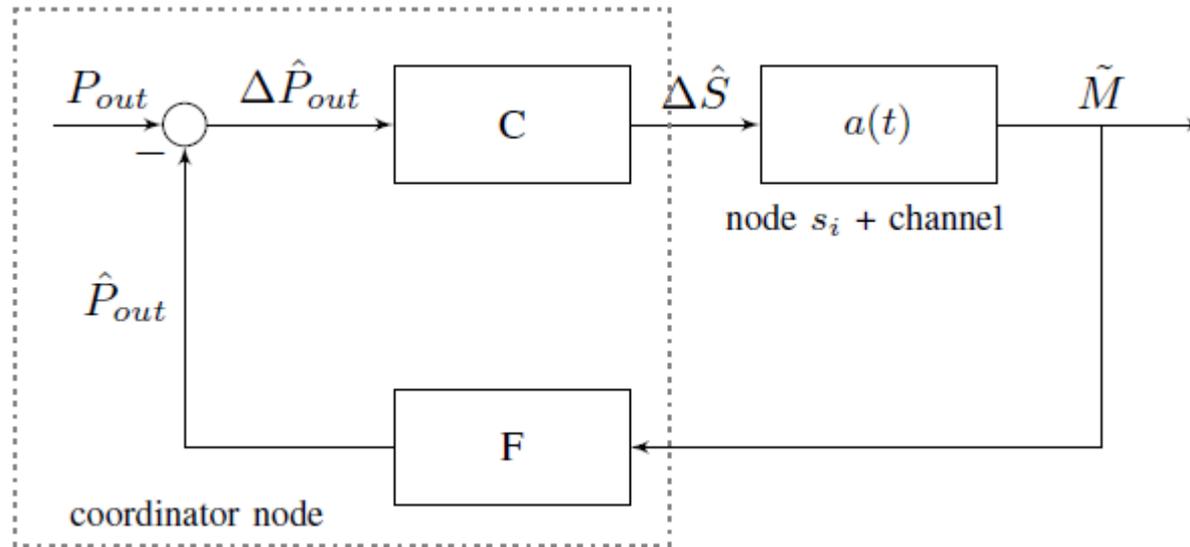
$$\hat{m} = T_w R$$

- Expected number of successfully received messages:

Outage Probability

$$\tilde{m} \longrightarrow = m (1 - P_{out})$$

Control Loop Strategy



- **Control loop** runs on the **Coordinator Node** to regulate **Application Nodes** energy S
- The possible control actions are
 - regulate the **transmit power** of Application Nodes
 - reconfigure network to achieve a **target availability**

Node Energy and Lifetime

- The instantaneous energy of a node is defined by:

$$E_T(t) = E_0 + E_H(t) - E_C(t)$$

- The consumed energy of a node can be defined as:

$$E_C(t) = S_{avg}t + V_{dd}I_qt$$

$$S_{avg} = T_m R \hat{S}$$

- Rewriting $E_T(t)$: $E_T(t) = E_0 + E_H(t) - (T_m R \hat{S} + V_{dd}I_q)t$

- Therefore, the lifetime of a node is defined by:

$$t_{lt} = \frac{E_0 + E_H - E_{min}}{T_m R \hat{S} + V_{dd}I_q}$$

Network Availability

- The network availability α is defined as the probability that the lifetime of a key node falls below a minimum value t_{min}

$$P_r (t_{lt} < t_{min}) = 1 - \alpha$$

- The closed form of the network availability is given by the estimated transmitted messages and the outage probability in the form of a binomial probability as defined below:

$$P_r (t_{lt} < t_{min}) = 1 - \sum_{k=0}^{\lfloor B(R) \rfloor} \binom{\lfloor \hat{m} \rfloor}{k} (1 - P_{out})^k (P_{out})^{\lfloor \hat{m} \rfloor - k}$$

with

$$A(R) \triangleq \frac{1}{RT_m} \left(\frac{E_0 + E_H - E_{min}}{t_{min}} - V_{dd}I_q \right)$$

$$B(R) \triangleq RT_w \left(1 - \frac{\gamma_0 \sigma_N^2}{A(R)} \right)$$

Network Availability vs. Estimation Accuracy

- The estimation accuracy ϵ of an application parameter must be guaranteed in order to meet the application's demands

$$\sigma_{\theta}^2(\tilde{m}) \triangleq \frac{\sigma_V^2}{\tilde{m}} \leq \epsilon$$

- Thus, for a target network availability α , we can state:

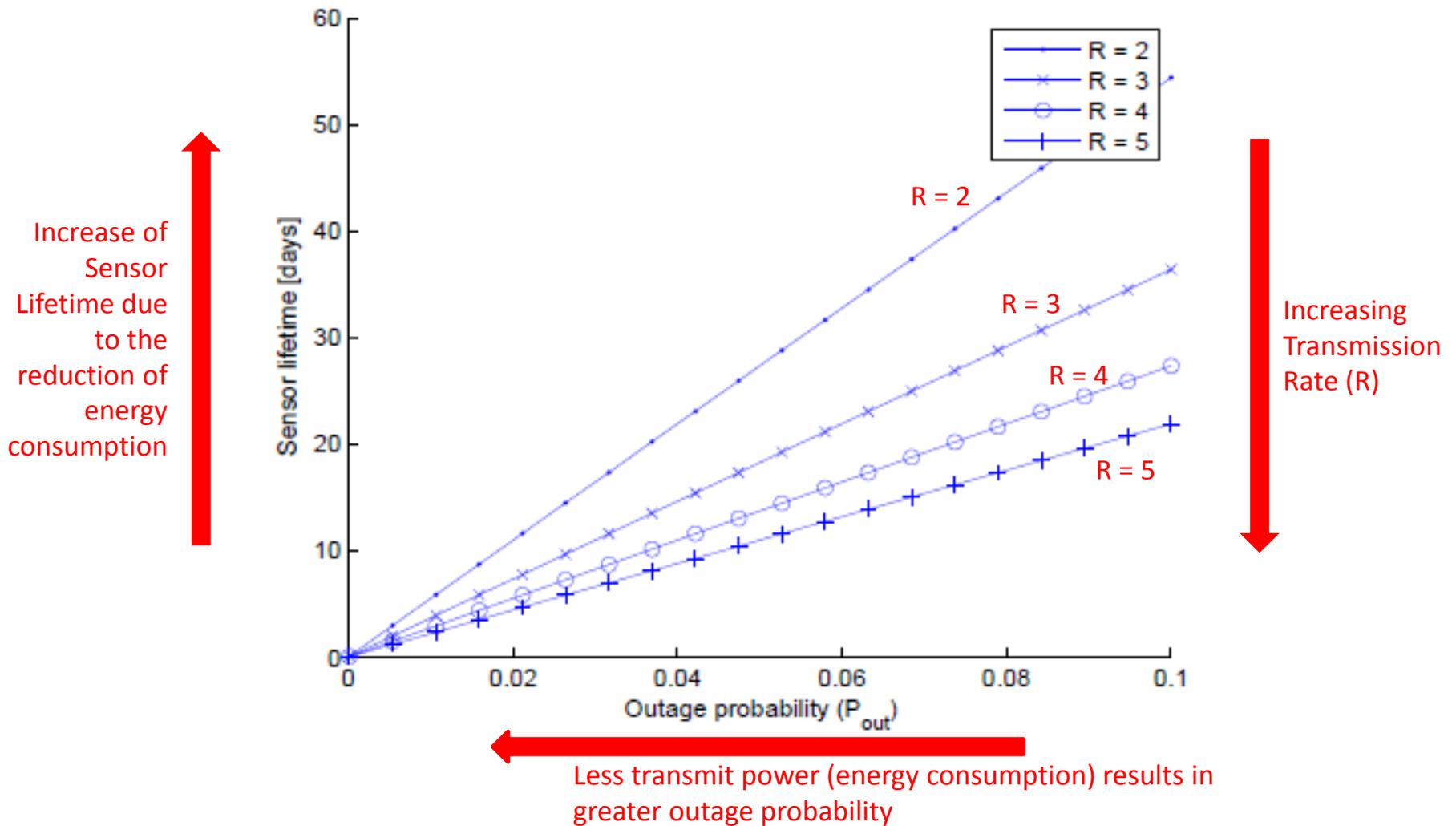
$$\sum_{k=0}^{\lfloor B(R) \rfloor} \binom{\lfloor (RT_w) \rfloor}{k} \left(\frac{\sigma_V^2}{\epsilon T_w R} \right)^k \left(1 - \frac{\sigma_V^2}{\epsilon T_w R} \right)^{\lfloor (RT_w) \rfloor - k} = \alpha$$

- The message rate R_s that satisfied the previous equation defines the setpoint of P_{out} as given by:

$$P_{out} \leq \left(1 - \frac{\sigma_V^2}{\epsilon R_s T_w} \right) H \left(1 - \frac{\sigma_V^2}{\epsilon R_s T_w} \right)$$

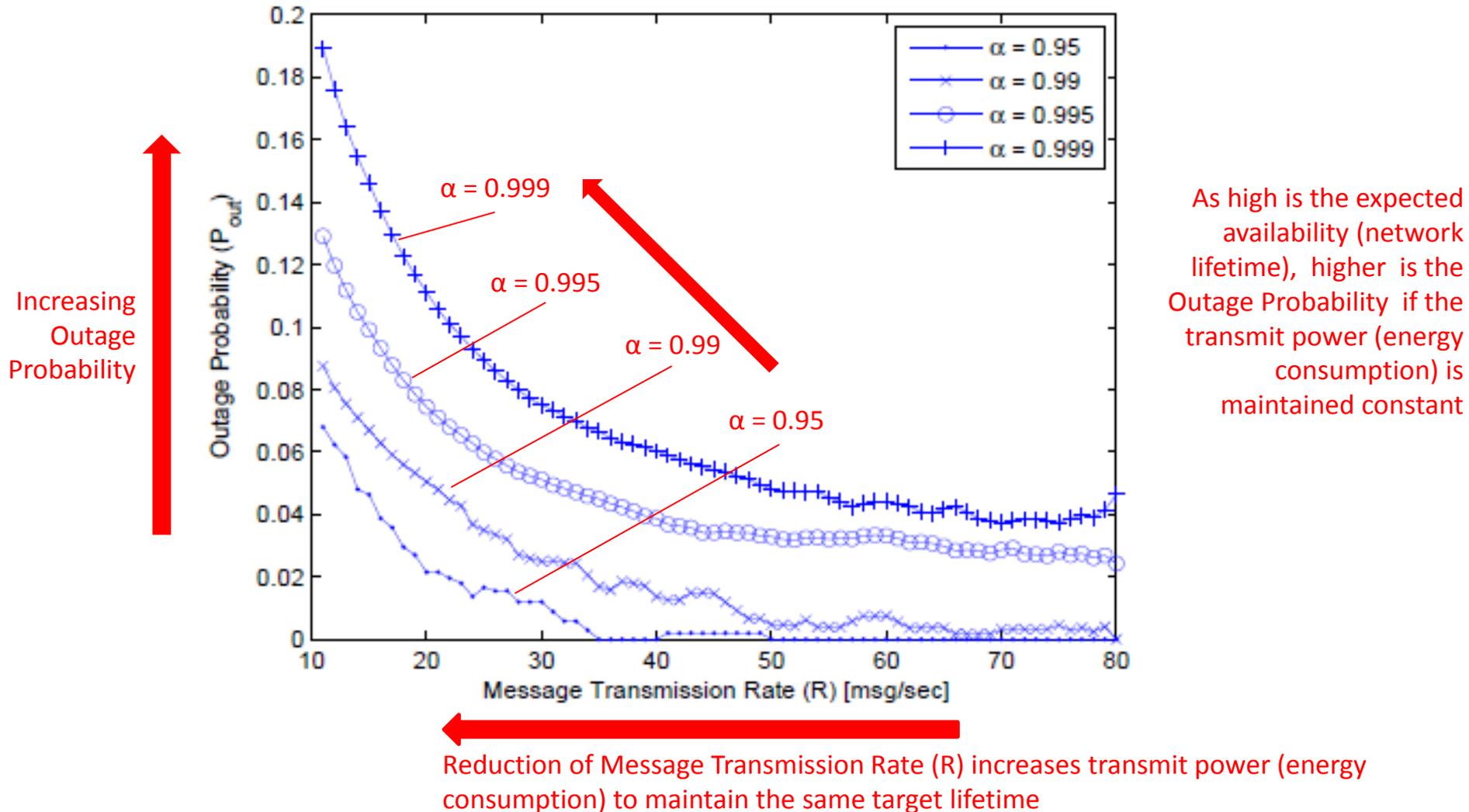
A target α can be achieved by tuning R , P_{out} and trading ϵ

Simulation Results*



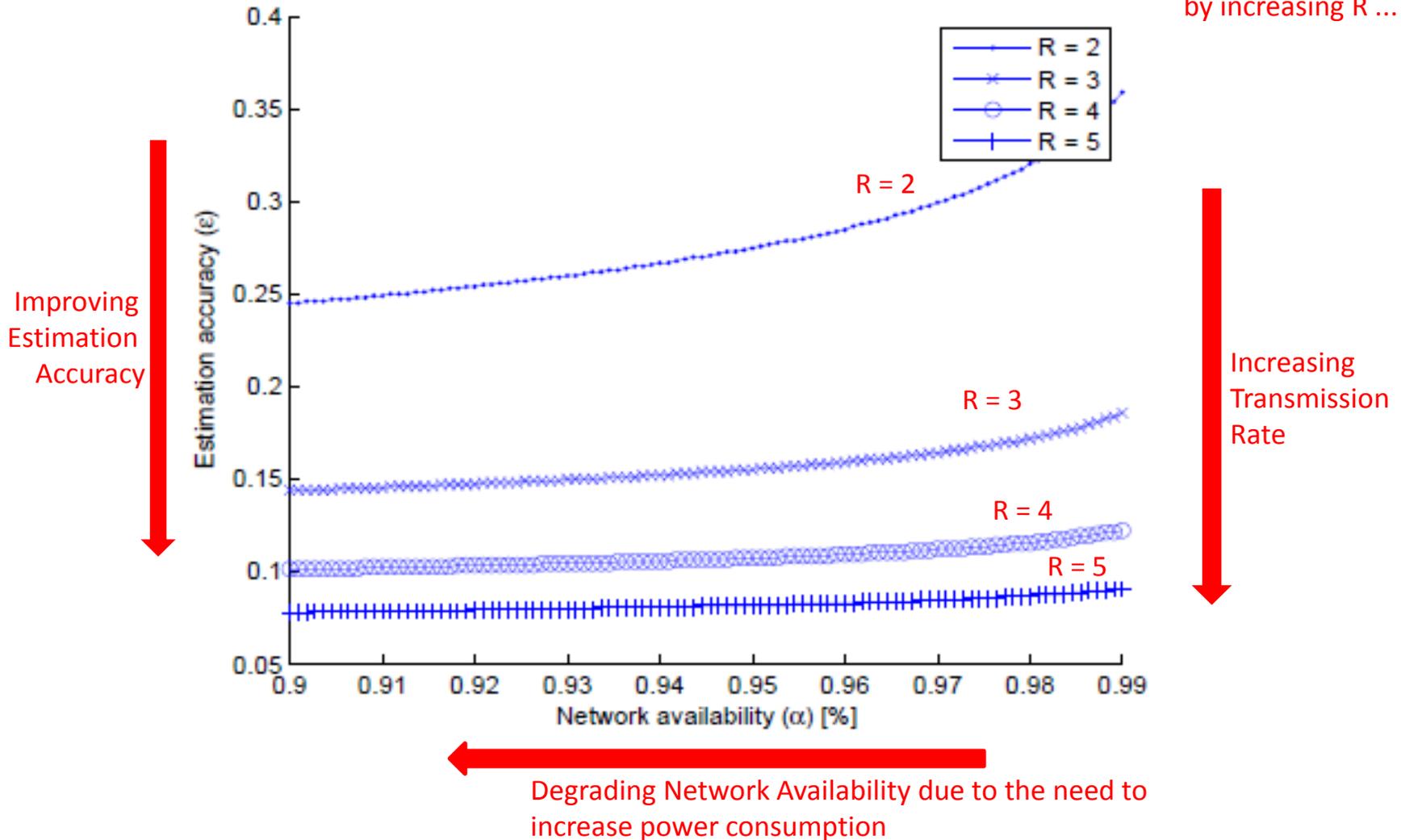
* Proposed system modeled and simulated in MATLAB®

Simulation Results (Cont'd)



The message transmission rate that fulfills the target estimation accuracy requirement defines the outage probability setpoint. **The greater R and P_{out} , the greater the energy consumption.**

Simulation Results (Cont'd)



Better (i.e., smaller) **estimation accuracy** implies worse (i.e., smaller) **network availability** due to **increased power consumption**, which reduces node lifetime.

Questions?

Thanks for your kind attention!