



Comparison of Four Parallel Algorithms for Domain Decomposed Implicit Monte Carlo

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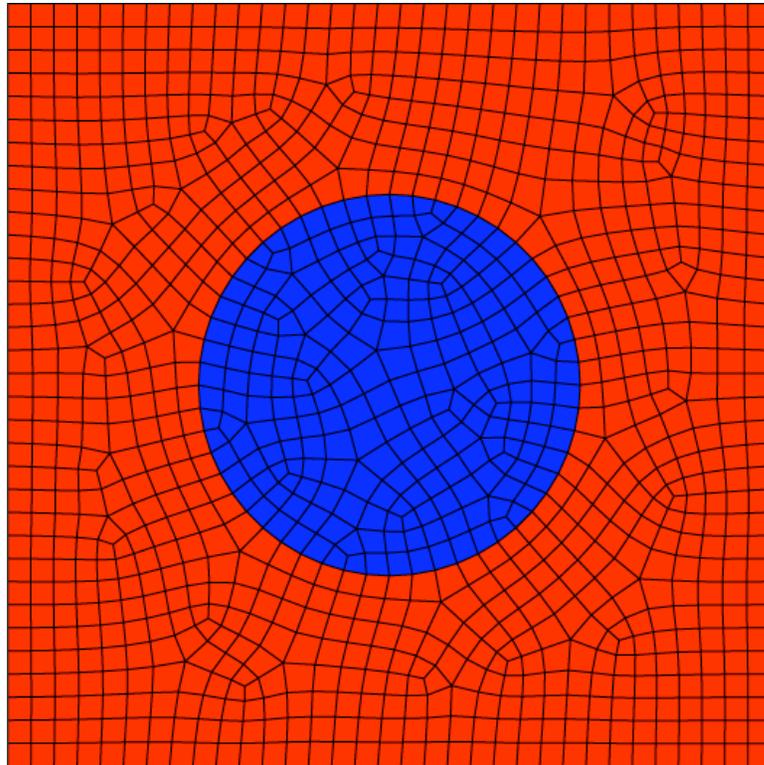
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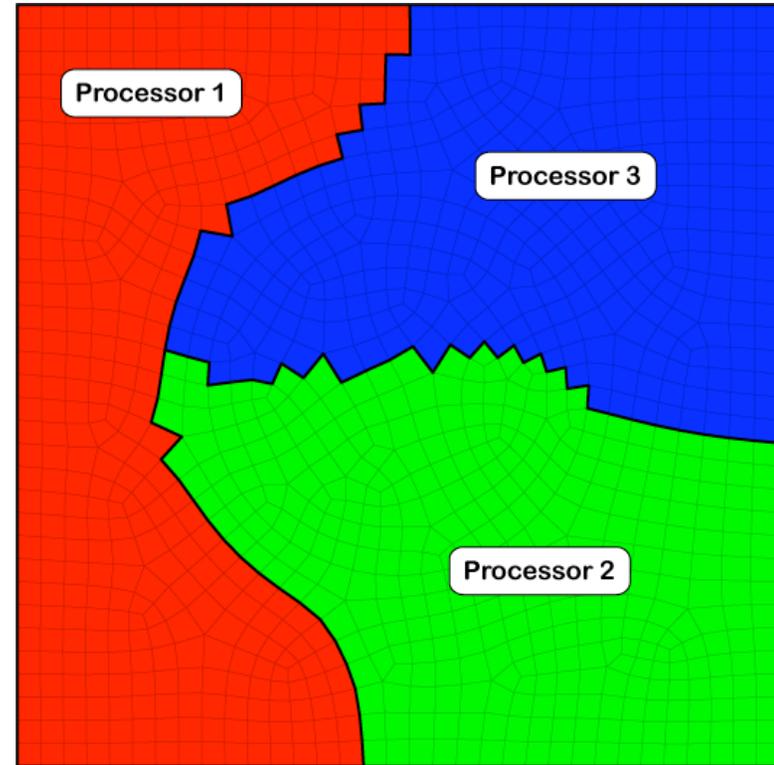




Particle Simulation On A Decomposed Mesh



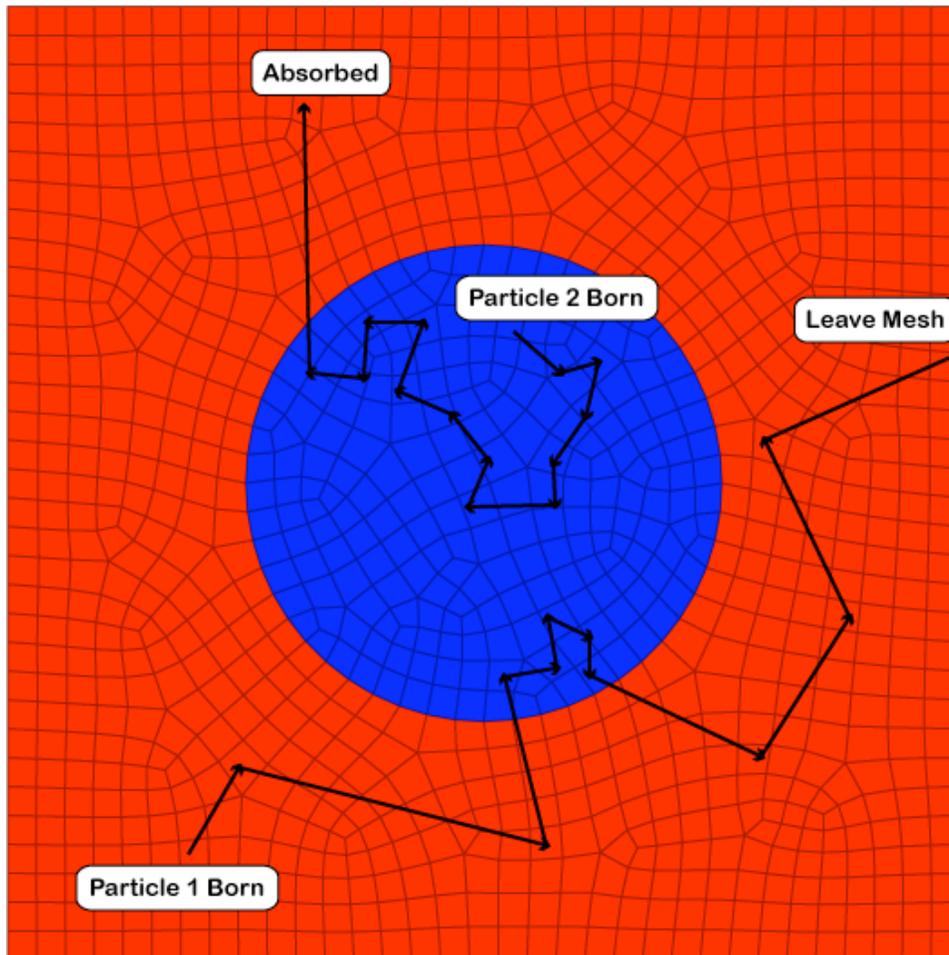
- Typical user mesh
- Cylinder of one material
- Another material outside



- Typical decomposition
- Load balances zones
- Irregular boundaries



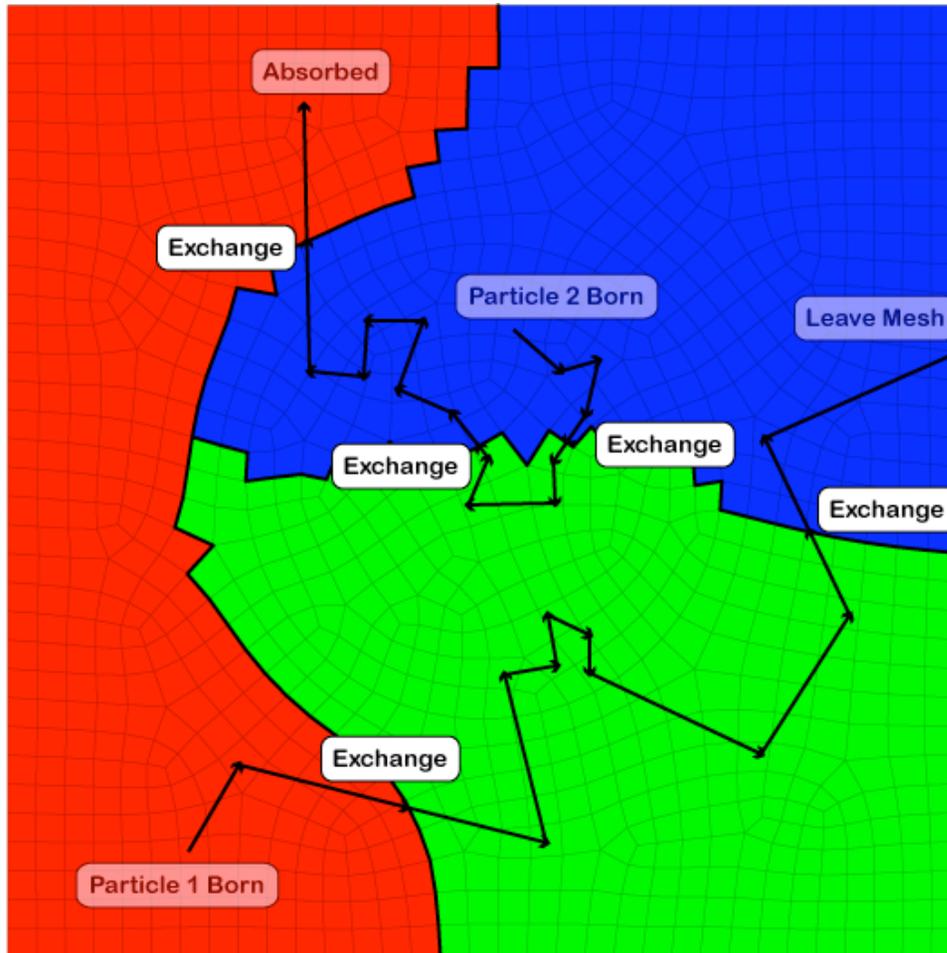
Physical Simulation



- Particles born in mesh
- Scatter in materials
- Deposit energy in zones
- Photons are terminated
 - by leaving mesh
 - being absorbed
 - or reaching census (the end of a time step).



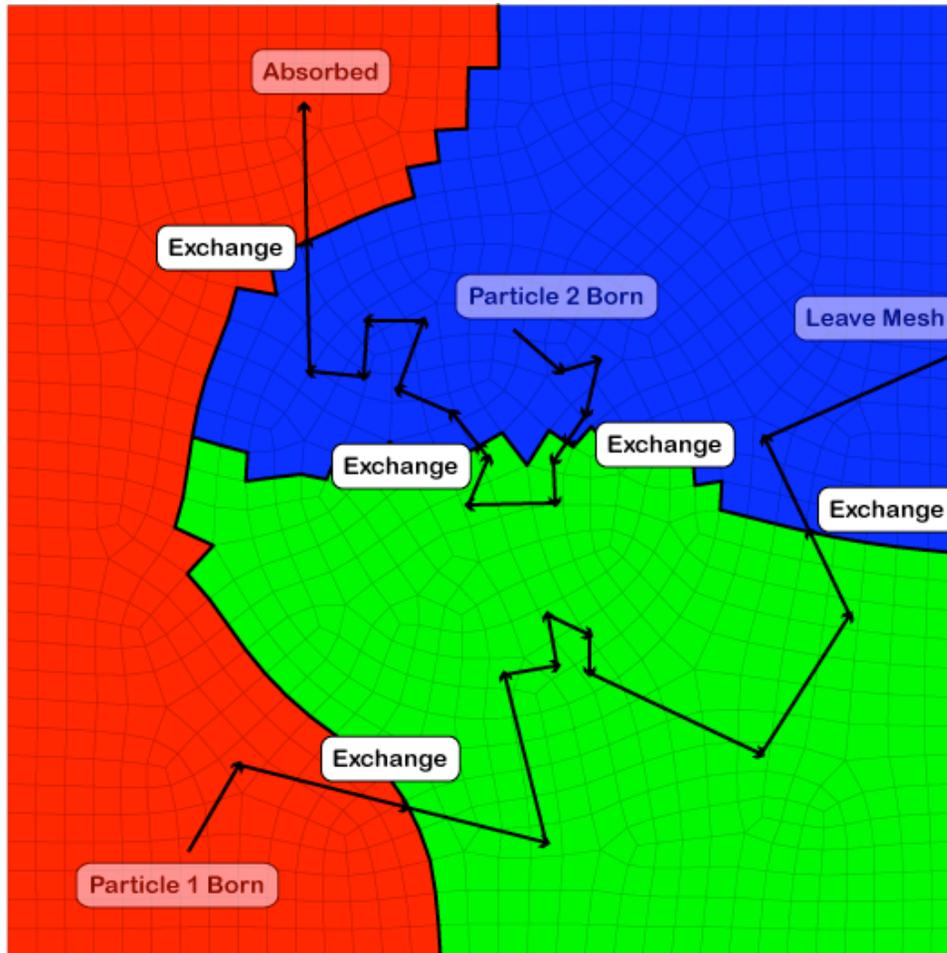
Decomposed Simulation



- **New events for decomposed mesh:**
 - Particles must be exchanged between processors.
 - Need to determine globally when all particles are finished.
- **These four algorithms do these two steps in different ways.**



Algorithm 1: KULL

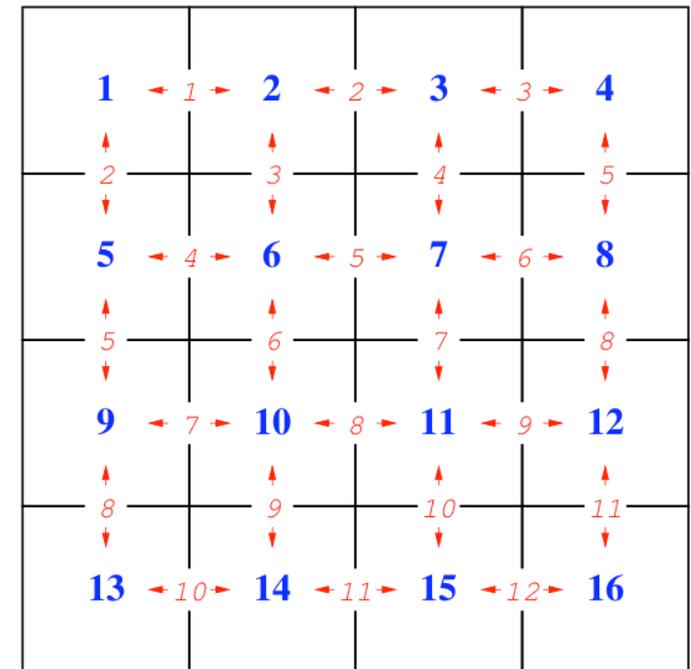


1. Each processor transports local particles until all have:
 - hit processor boundaries,
 - been absorbed,
 - leaked off mesh, or
 - reached census.
2. Blocking exchange particles buffered on processor boundaries with neighbors.
3. Global sum tallying number of remaining particles to simulate on all processors.
4. If particle sum is nonzero, go to step 1.



Problems With Algorithm 1

- **Blocking sends for particle exchanges cause serialization of communication.**
- **Multiple synchronization points during a time step for global sum that is used to determine if all particles have finished.**
- **Algorithm 2 modifies Algorithm 1 to use asynchronous particle exchanges, fixing the first problem.**
- **Algorithm 3 eliminates the synchronization points, but it still does not scale well.**

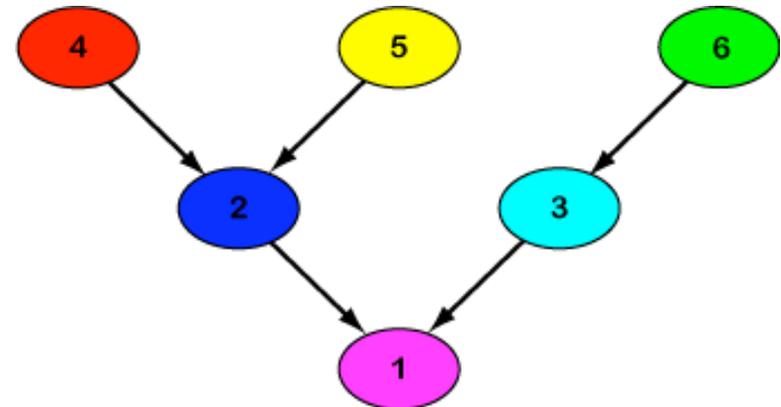
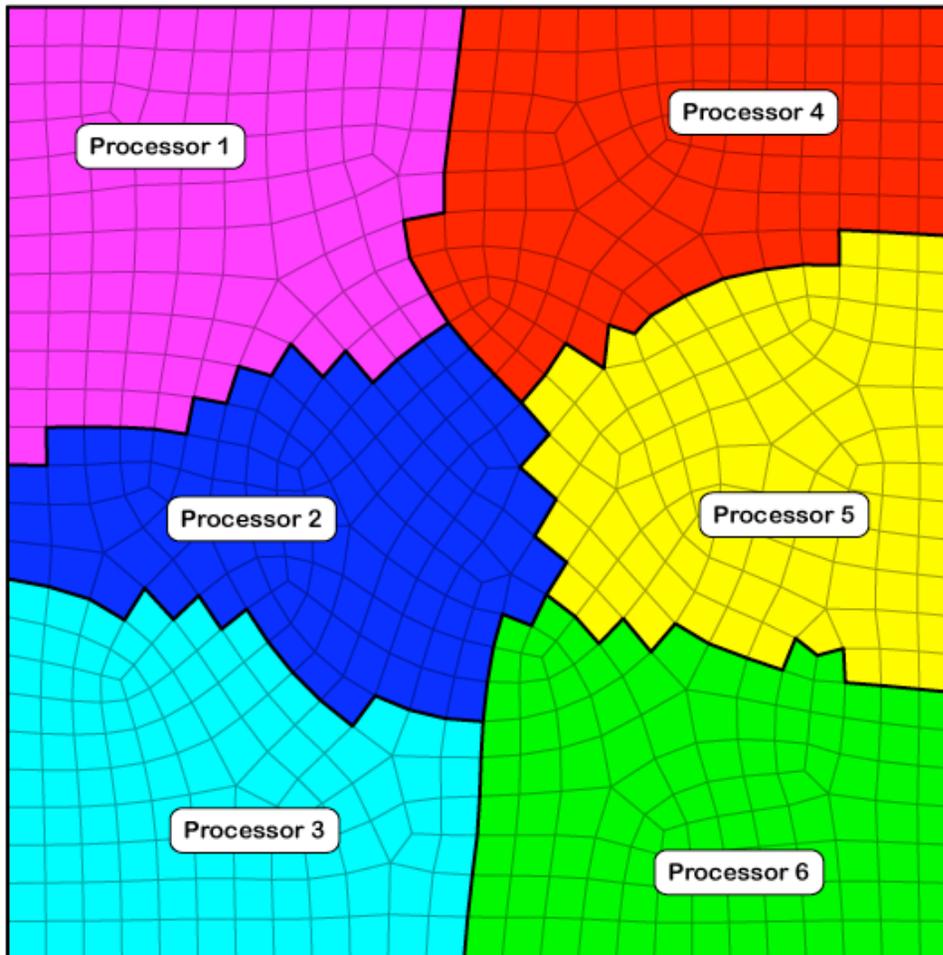


Key: **Processor/Domain Number**
← *Communication Step* →

Another decomposition to illustrate the serialization of Algorithm 1.



Algorithm 4: Improved Milagro



- More frequent asynchronous particle exchanges of smaller, fixed-size buffers.
- Master (Processor 1) keeps track of total number of particles completed
 - Collected asynchronously through a binary tree.

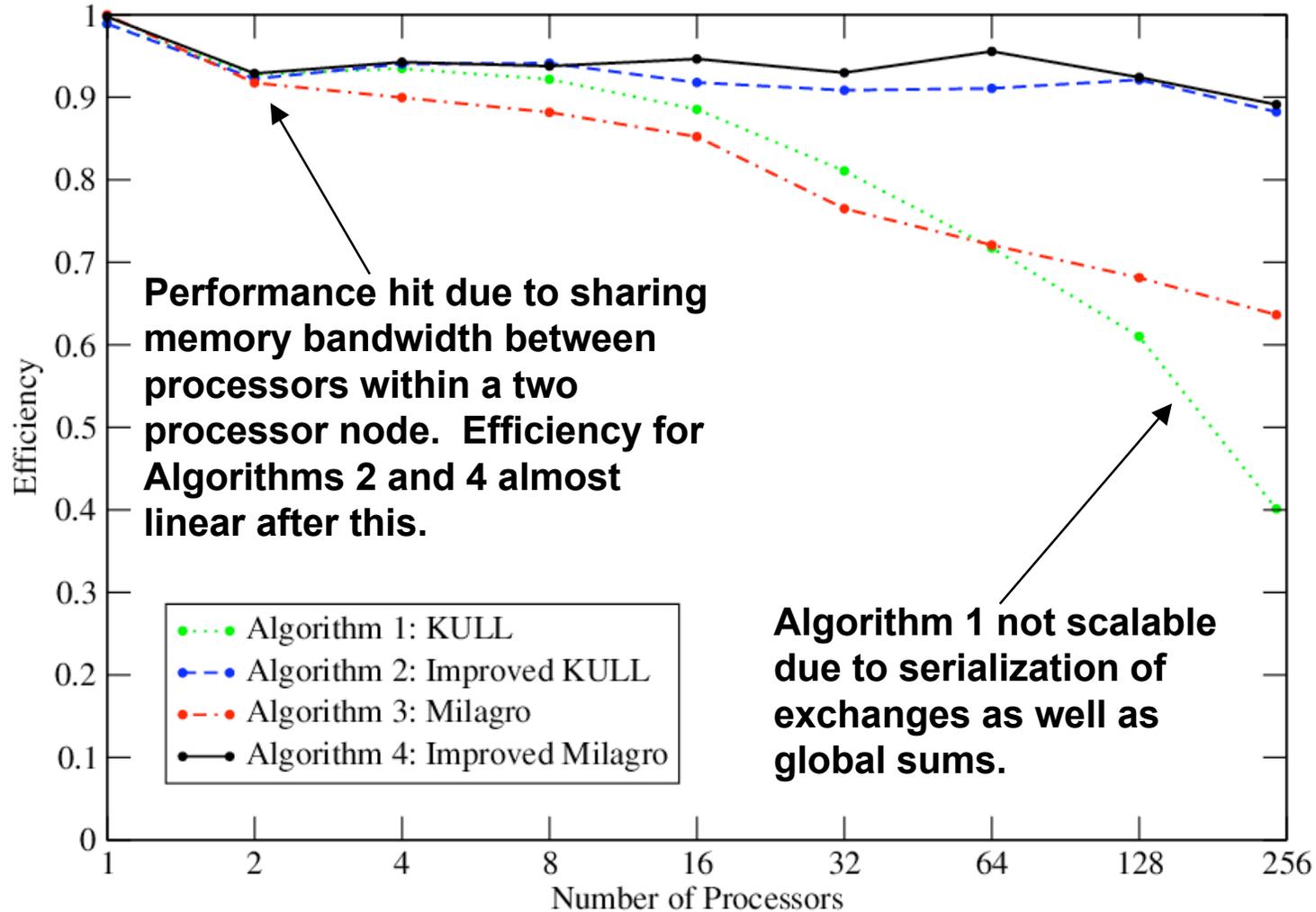


Test Problem 1: A Hot Box

- **Uniform temperature box**
 - 60x60x60 zones
 - Each zone about one mean free path
- **Reflecting boundary conditions**
- **Perfectly load balanced**
- **Constant work (strong) scaling study**
 - Subdivided same exact problem on multiple decompositions
 - Same exact result from all four algorithms



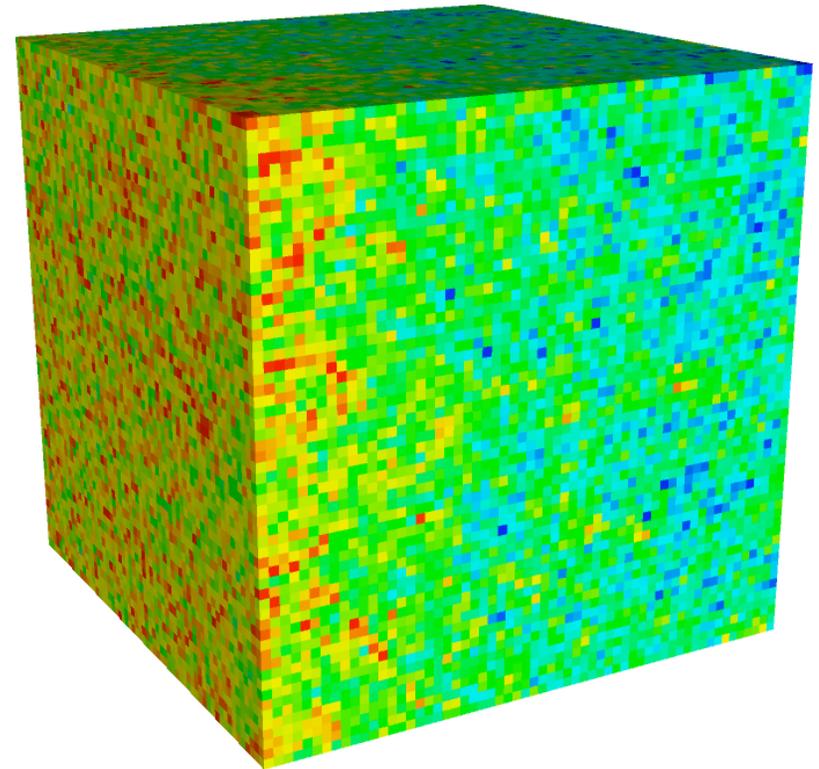
Hot Box Constant Work Study





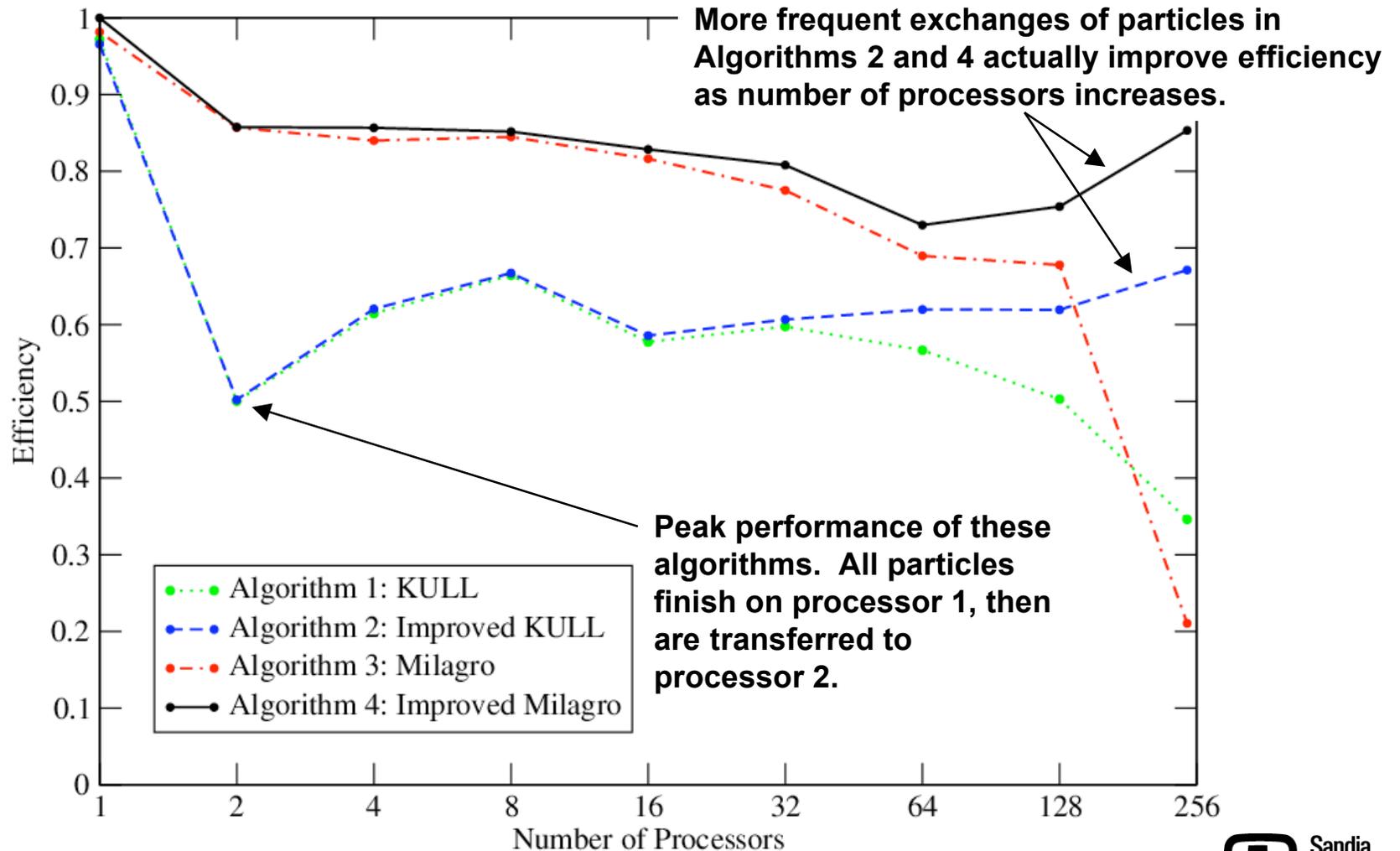
Test Problem 2: Mild Load Imbalance

- Vacuum Box
- Reflecting boundaries
- Initially Cold
- Temperature source on one side
- Initially load imbalanced
- Becomes balanced by end





Vacuum Box Constant Work Scaling





Conclusions

- **The improved Milagro algorithm, Algorithm 4, scales almost linearly for load balanced problems.**
- **Critical to avoid all synchronization points in algorithms, either through blocking sends or global sums.**
- **Each processor must have the same parallel overhead.**
- **More frequent exchanges of particles can help in load imbalanced problems.**