

Weak Cardinality Theorems for First-Order Logic

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Fundamentals of Computation Theory 2003

Outline

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Motivation of Enumerability

Problem

Many functions are not computable or not efficiently computable.

Example

- #SAT:
How many satisfying assignments does a formula have?

Motivation of Enumerability

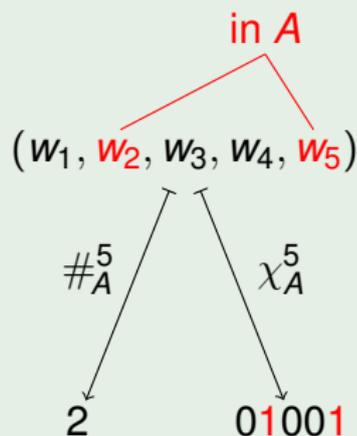
Problem

Many functions are not computable or not efficiently computable.

Example

For difficult languages A :

- Cardinality function $\#_A^n$:
How many input words are in A ?
- Characteristic function χ_A^n :
Which input words are in A ?



Motivation of Enumerability

Problem

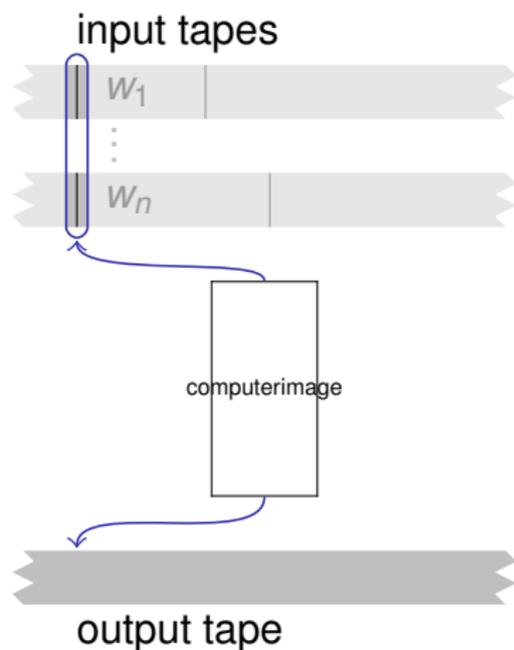
Many functions are not computable or not efficiently computable.

Solutions

Difficult functions can be

- computed using probabilistic algorithms,
- computed efficiently on average,
- approximated, or
- **enumerated.**

Enumerators Output Sets of Possible Function Values

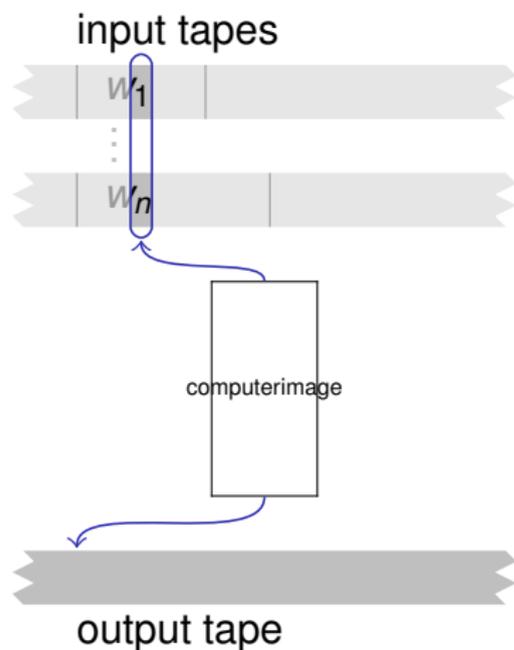


Definition (1987, 1989, 1994, 2001)

An *m*-enumerator for a function f

- 1 reads n input words w_1, \dots, w_n ,
- 2 does a computation,
- 3 outputs at most m values,
- 4 one of which is $f(w_1, \dots, w_n)$.

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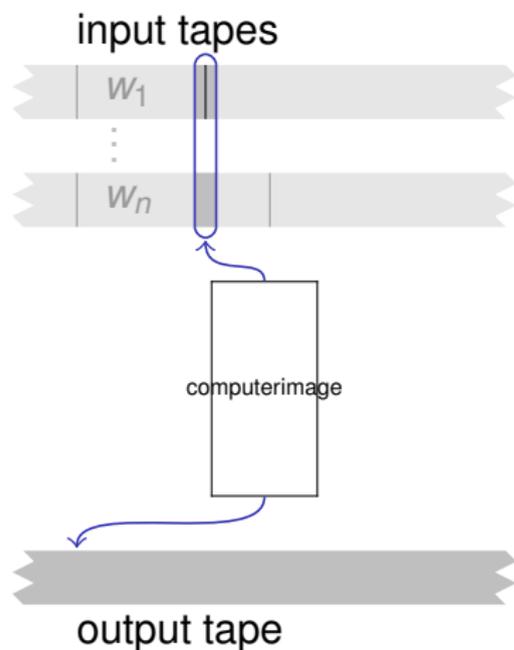


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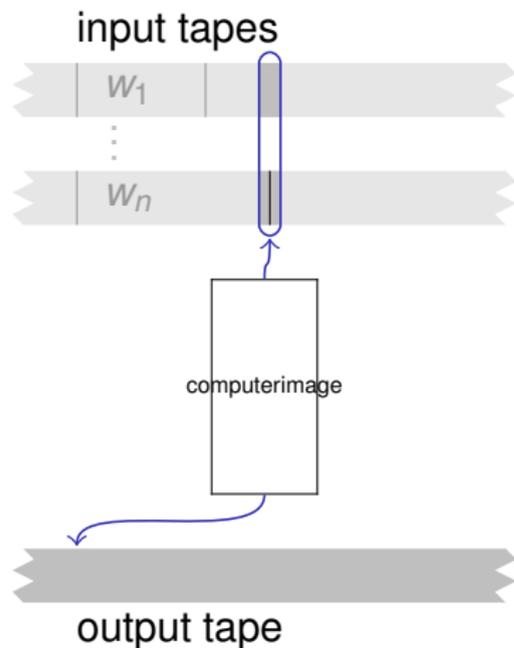


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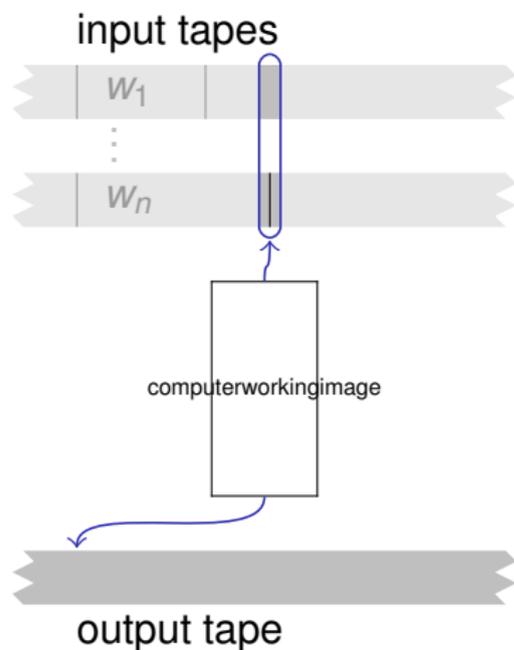


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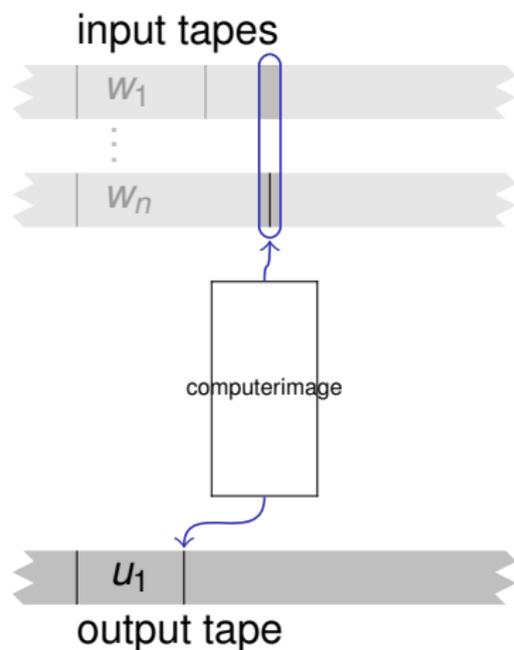


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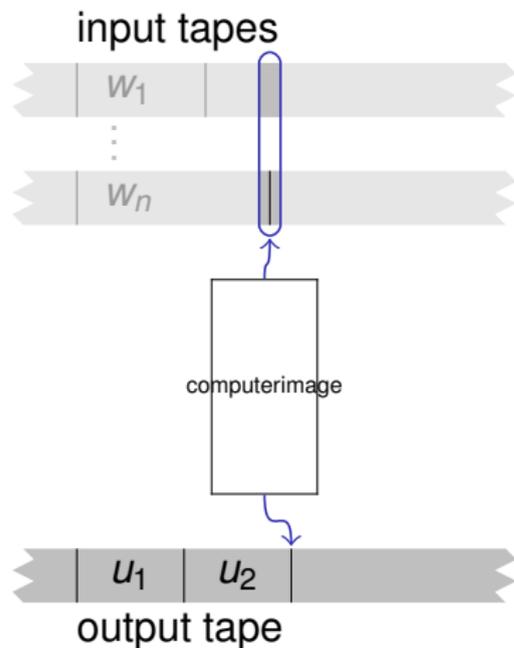


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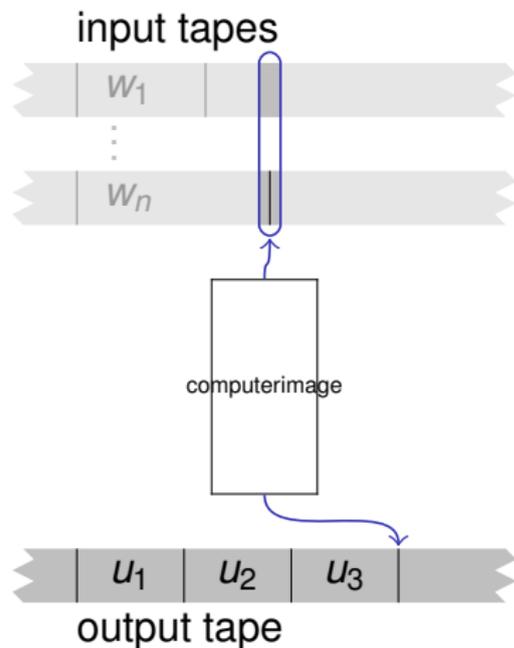


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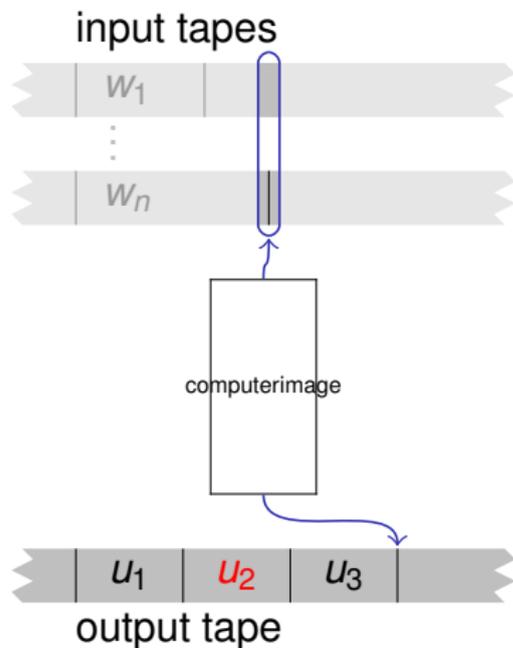


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How Well Can the Cardinality Function Be Enumerated?

Observation

For fixed n , the cardinality function $\#_A^n$

- can be 1-enumerated by Turing machines only for **recursive** A , but
- can be $(n + 1)$ -enumerated for **every** language A .

Question

What about 2-, 3-, 4-, ..., n -enumerability?

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How Well Can the Cardinality Function Be Enumerated by Turing Machines?

Cardinality Theorem (Kummer, 1992)

If $\#_A^n$ is n -enumerable by a Turing machine, then A is recursive.

Weak Cardinality Theorems ()

- *If χ_A^1 is 1-enumerable by a Turing machine, then A is recursive.*
- *If $\#_A^2$ is 2-enumerable by a Turing machine, then A is recursive.*

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- 1** *If χ_A^n is n -enumerable by a Turing machine, then A is recursive.*
- 2** *If $\#_A^2$ is 2-enumerable by a Turing machine, then A is recursive.*
- 3** *If $\#_A^n$ is n -enumerable by a Turing machine that never enumerates both 0 and n , then A is recursive.*

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How Well Can the Cardinality Function Be Enumerated by Finite Automata?

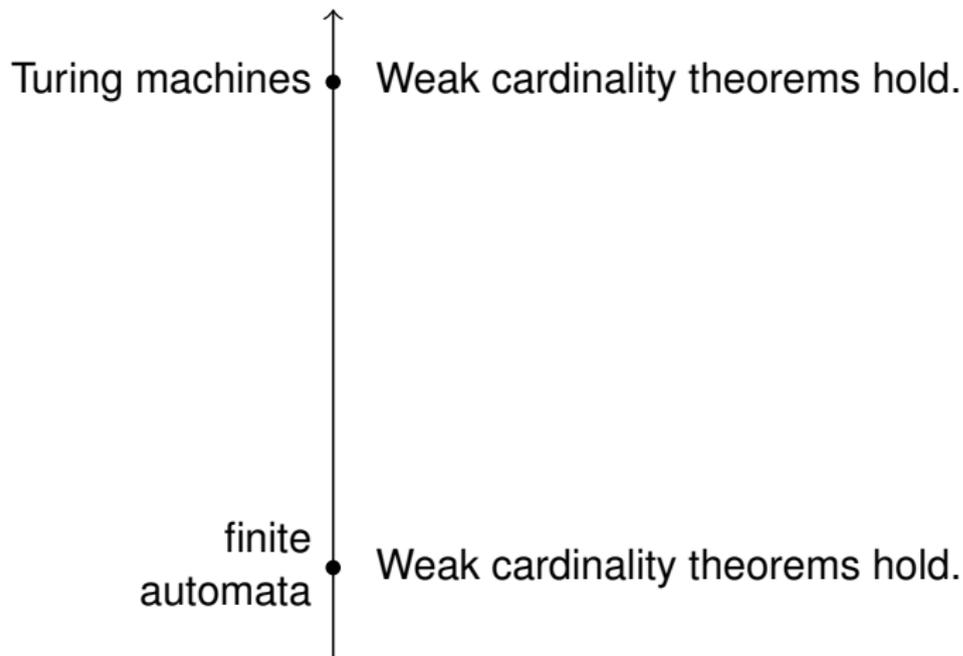
Conjecture

If $\#_A^n$ is n -enumerable by a **finite automaton**, then A is **regular**.

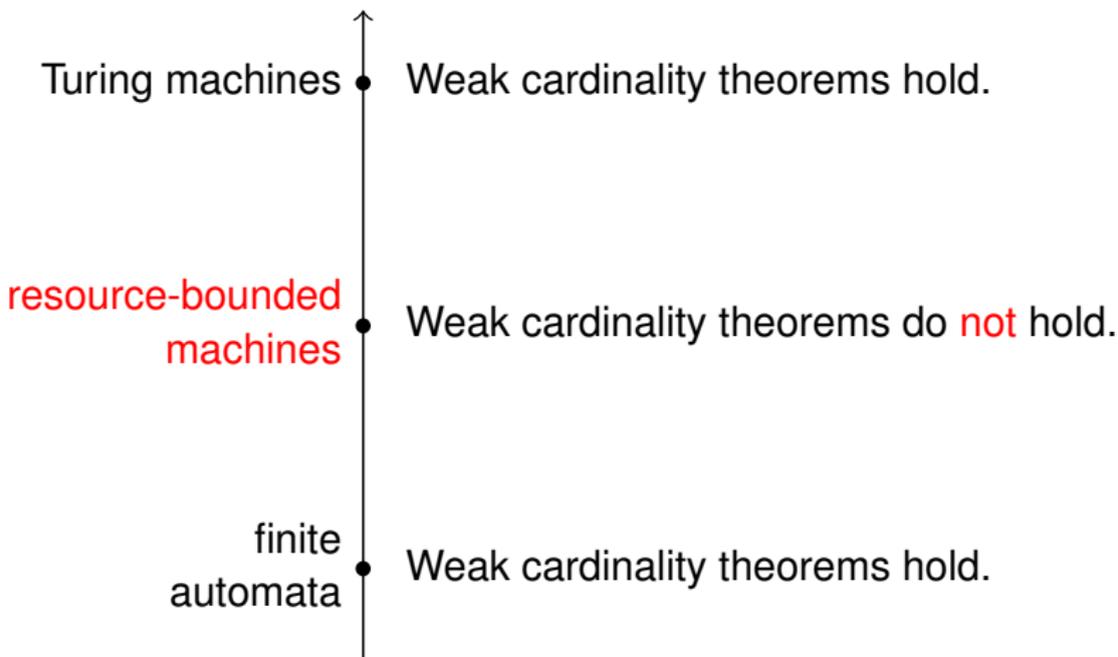
Weak Cardinality Theorems (2001, 2002)

- 1 If χ_A^n is n -enumerable by a **finite automaton**, then A is **regular**.
- 2 If $\#_A^2$ is 2-enumerable by a **finite automaton**, then A is **regular**.
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Cardinality Theorems Do Not Hold for All Models



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Why?

First Explanation

The weak cardinality theorems hold both for recursion and automata theory **by coincidence**.

Second Explanation

The weak cardinality theorems hold both for recursion and automata theory, **because they are instantiations of single, unifying theorems**.

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The weak cardinality theorems hold both for recursion and automata theory **by coincidence**.

Second Explanation

The weak cardinality theorems hold both for recursion and automata theory, **because they are instantiations of single, unifying theorems**.

The second explanation is correct.

The theorems can (almost) be unified using first-order logic.

Outline

What Are Elementary Definitions?

Definition

A relation R is **elementarily definable in a logical structure \mathcal{S}** if

- 1 there exists a first-order formula ϕ ,
- 2 that is true exactly for the elements of R .

Example

The set of even numbers is elementarily definable in $(\mathbb{N}, +)$ via the formula $\phi(x) \equiv \exists z . z + z = x$.

Example

The set of powers of 2 is not elementarily definable in $(\mathbb{N}, +)$.

Characterisation of Classes by Elementary Definitions

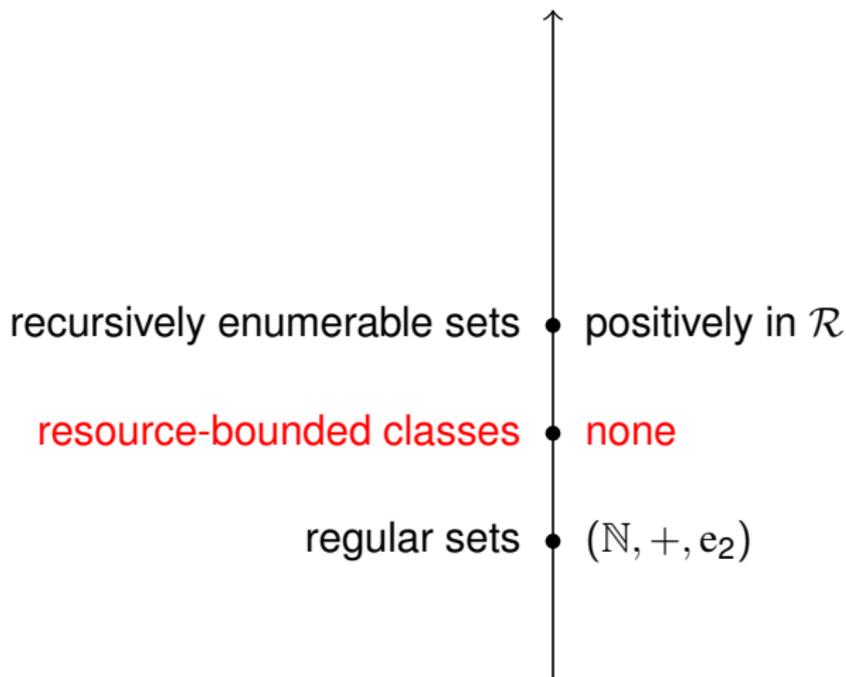
Theorem (Büchi, 1960)

There exists a logical structure $(\mathbb{N}, +, e_2)$ such that a set $A \subseteq \mathbb{N}$ is *regular* iff it is *elementarily definable* in $(\mathbb{N}, +, e_2)$.

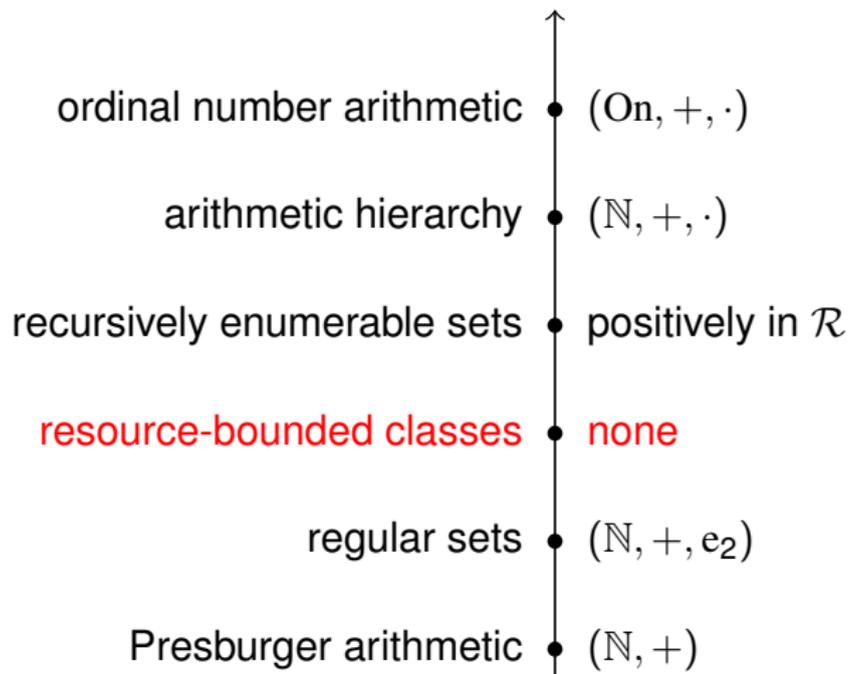
Theorem

There exists a logical structure \mathcal{R} such that a set $A \subseteq \mathbb{N}$ is *recursively enumerable* iff it is *positively elementarily definable* in \mathcal{R} .

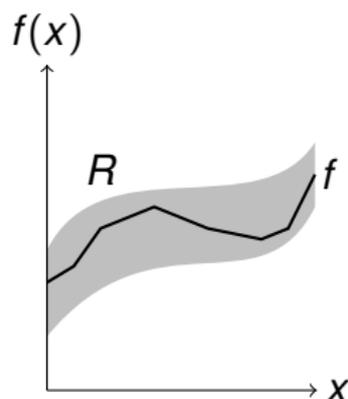
Characterisation of Classes by Elementary Definitions



Characterisation of Classes by Elementary Definitions



Elementary Enumerability is a Generalisation of Elementary Definability



Definition

A function f is **elementarily m -enumerable** in a structure \mathcal{S} if

- 1 its graph is contained in an **elementarily definable** relation R ,
- 2 which is **m -bounded**, i.e., for each x there are at most m different y with $(x, y) \in R$.

The Original Notions of Enumerability are Instantiations

Theorem

A function is m -enumerable by a *finite automaton* iff it is elementarily m -enumerable in $(\mathbb{N}, +, e_2)$.

Theorem

A function is m -enumerable by a *Turing machine* iff it is positively elementarily m -enumerable in \mathcal{R} .

The First Weak Cardinality Theorem

Theorem

Let S be a logical structure with universe U and let $A \subseteq U$. If

- 1 S is well-orderable and
- 2 χ_A^n is elementarily n -enumerable in S ,

then A is elementarily definable in S .

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Corollary

If χ_A^n is n -enumerable by a finite automaton, then A is regular.

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Corollary (with more effort)

If χ_A^n is n -enumerable by a Turing machine, then A is recursive.

The Second Weak Cardinality Theorem

Theorem

Let S be a logical structure with universe U and let $A \subseteq U$. If

- 1 S is well-orderable,
- 2 every finite relation on U is elementarily definable in S , and
- 3 $\#_A^2$ is elementarily 2-enumerable in S ,

then A is elementarily definable in S .

The Third Weak Cardinality Theorem

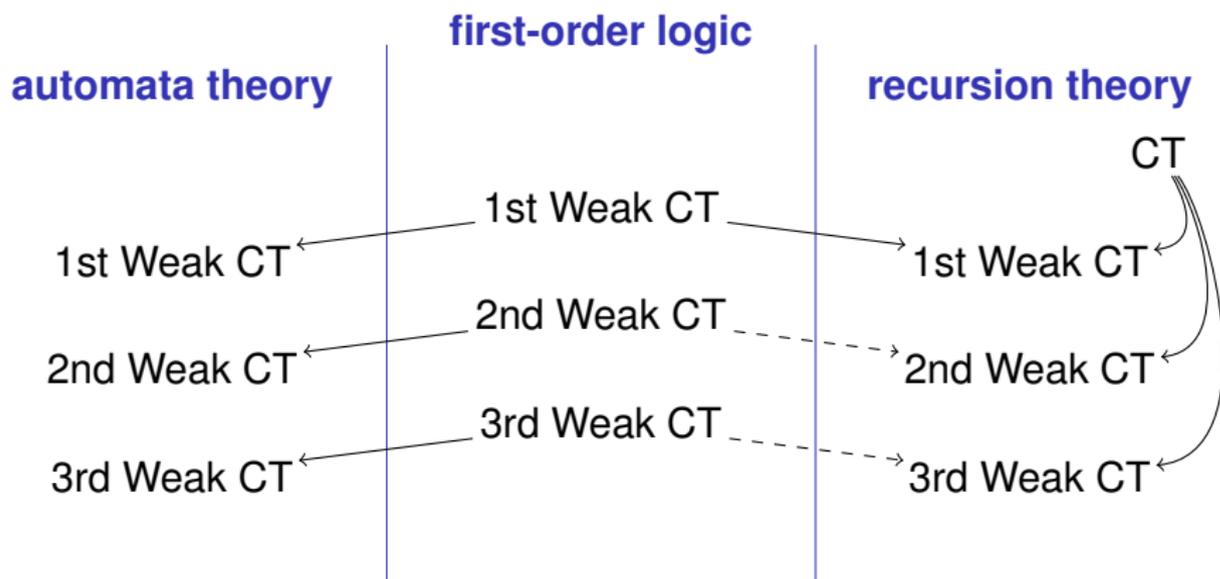
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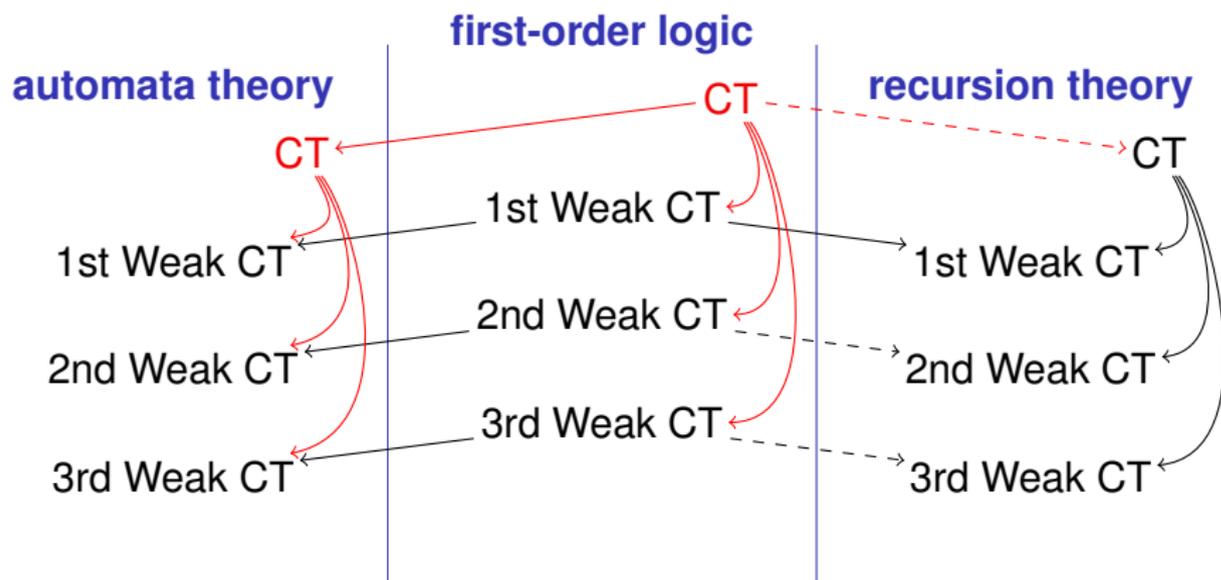
- 1 S is well-orderable,
- 2 every finite relation on U is elementarily definable in S , and
- 3 $\#_A^n$ is elementarily n -enumerable in S via a relation that never 'enumerates' both 0 and n ,

then A is elementarily definable in S .

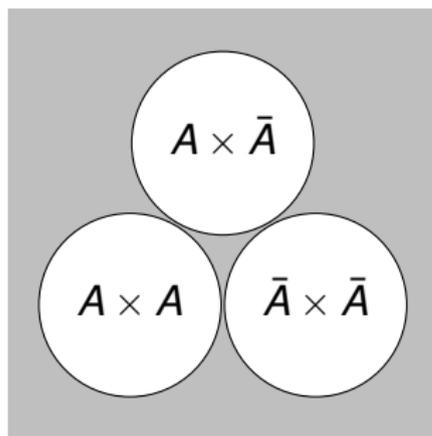
Relationships Between Cardinality Theorems (CT)



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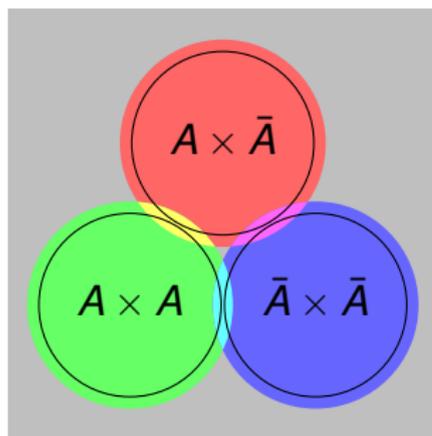
Theorem

Let S be a well-orderable logical structure in which all finite relations are elementarily definable.

If there exist elementarily definable supersets of $A \times A$, $A \times \bar{A}$, and $\bar{A} \times \bar{A}$ whose intersection is empty, then A is elementarily definable in S .

Note

The theorem is no longer true if we add $\bar{A} \times A$ to the list.



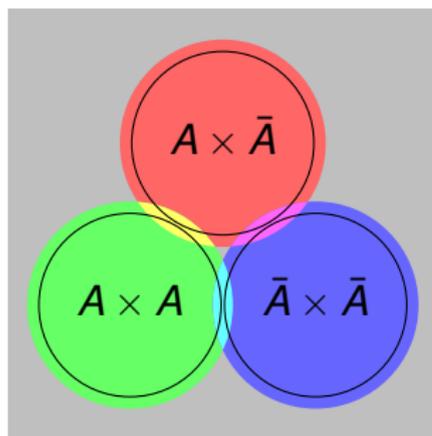
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Summary

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- The weak cardinality theorems for first-order logic **unify** the weak cardinality theorems of automata and recursion theory.
- The logical approach yields weak cardinality theorems for **other computational models**.
- Cardinality theorems are **separability theorems** in disguise.

Open Problems

- Does a cardinality theorem for first-order logic hold?
- What about non-well-orderable structures like $(\mathbb{R}, +, \cdot)$?