



Helsinki University of Technology

S-72.333 Postgraduate Seminar on Radio Communications

Multiple Access Methods

Er Liu

liuer@cc.hut.fi

Communications Laboratory

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Content of presentation

+ Protocol Classification

+ Conflict-Free Access Protocols

▣ Static Allocation

- TDMA
- FDMA
- CDMA

▣ Dynamic Allocation

- Polling
- Token passing

+ Contention Access Protocol

▣ Static Resolution

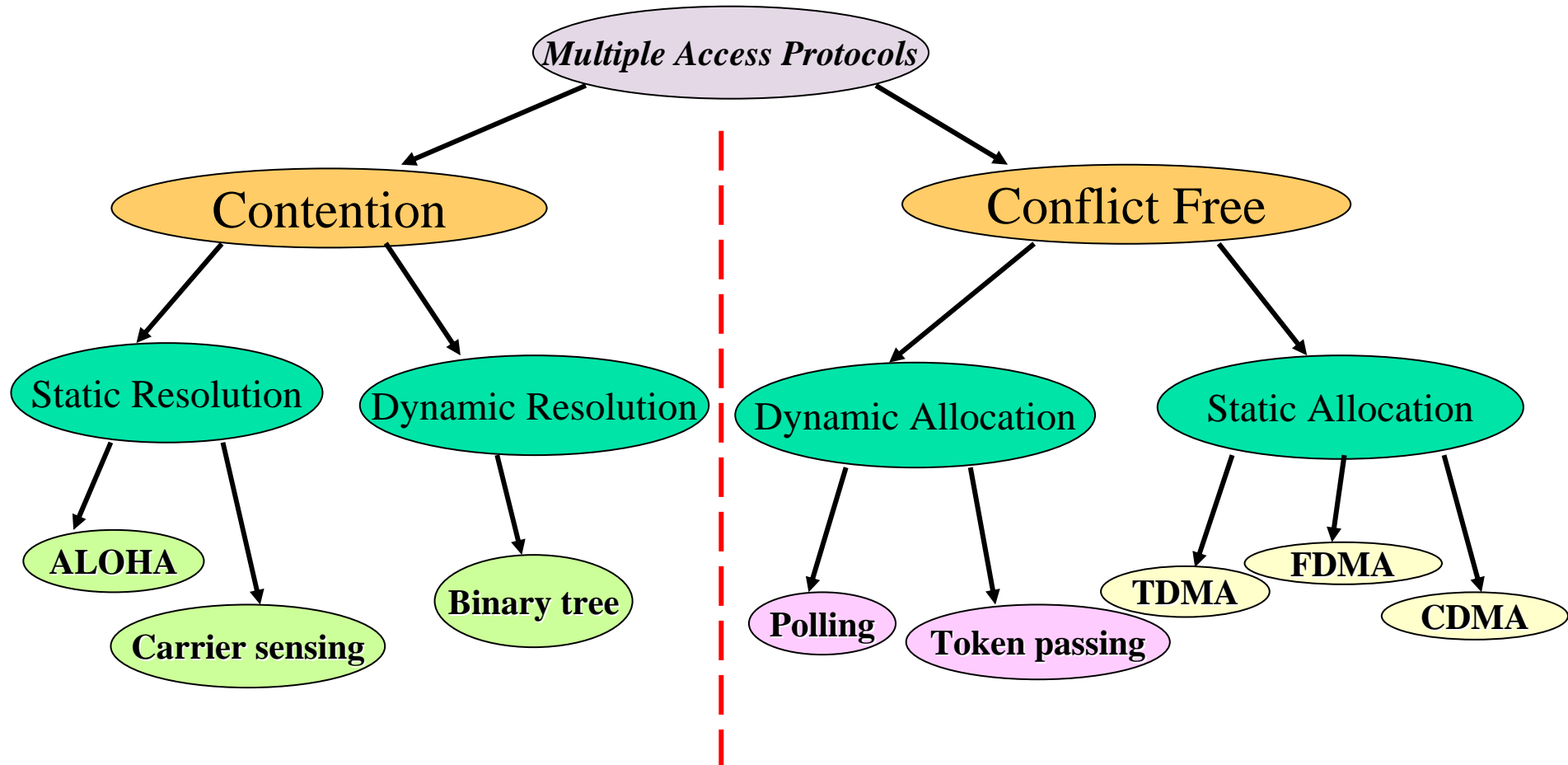
- Aloha
- Carrier sensing protocols

▣ Dynamic Resolution

- Binary Tree



Protocol Classification





Conflict-Free Access Protocols

- ✚ Static Allocation = Fixed Assignment:
 - ➔ *I.e. fixed allocated channel resource*
 - ➔ *Resource can be frequency, or time, or both*
 - ➔ *Predetermined basis to a single user*

- ✚ Basic access methods:
 - ➔ *FDMA – Frequency-Division Multiple Access*
 - ➔ *TDMA – Time-division Multiple Access*
 - ➔ *CDMA – Code-Division Multiple access*

- ✚ Some other formats:
 - ➔ *Combination of the basic access methods*
 - ➔ *Implemented with various multi-user access algorithm*

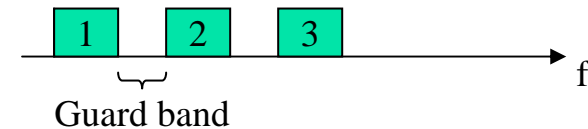
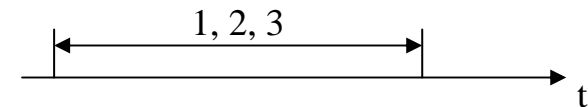


Frequency-Division Multiple Access

- ❏ Built upon *Frequency-division multiplexing* scheme
- ❏ It is the simplest and oldest form of multiplexing
- ❏ A fixed subchannel is assigned to a user terminal and is retained until released by the user
- ❏ At receiver, the user terminal filters the designated channel out of the composite signal

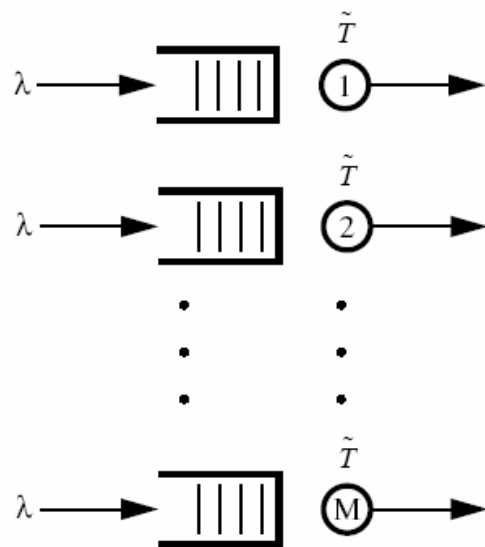
- ❏ Currently used in
 - Cellular mobile telephone
 - VHF & UHF land-mobile radio system
 - Satellite networks

- ❏ Characteristics:
 - Efficient when information is steady flow
 - Inefficient when data are sporadic



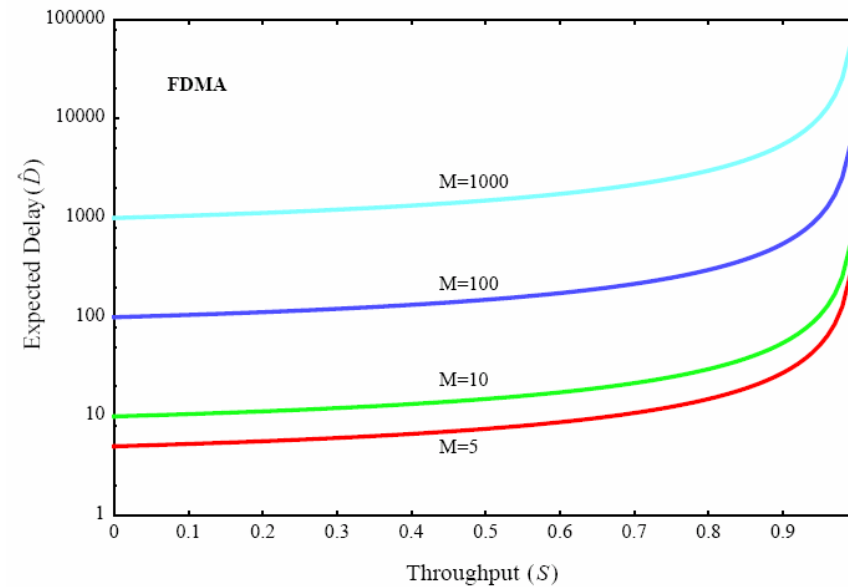


FDMA Performance



System model:

- M independent queues
- $M/G/1$ queueing system



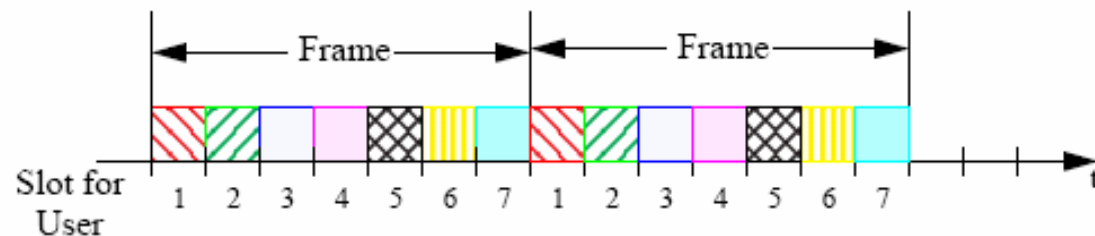
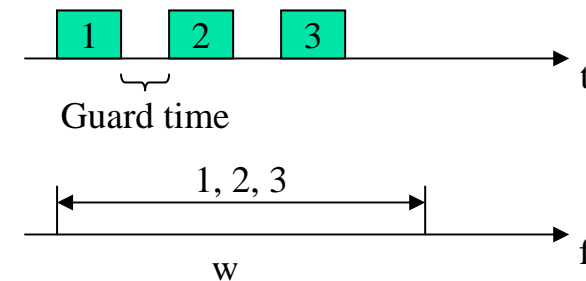
Throughput-Delay

$$\hat{D} = \frac{D}{P/R} = \left[1 + \frac{S}{2(1-S)} \right] M = M \frac{2-S}{2(1-S)} = \frac{M}{2} \left(1 + \frac{1}{1-S} \right)$$



Time-Division Multiple Access

- ❏ Built upon *Time-division multiplexing* transmission format
- ❏ Deterministic allocation of time interval – *time-slots*
- ❏ Time slots are organized into frames
 - ➔ T1 channel : 1.544 Mbits/s
 - ➔ Multiplexing 24 PCM encoded voice channel
 - ➔ Each channel is 64 kbits/s
 - ▶ Each channel sampled at an 8-kHz rate
 - ▶ Each sample is encoded into 8 bits
- ❏ Used in new digital cellular network
 - ➔ Europe (GSM)
 - ➔ Japan (JDC)
 - ➔ America (IS-54)





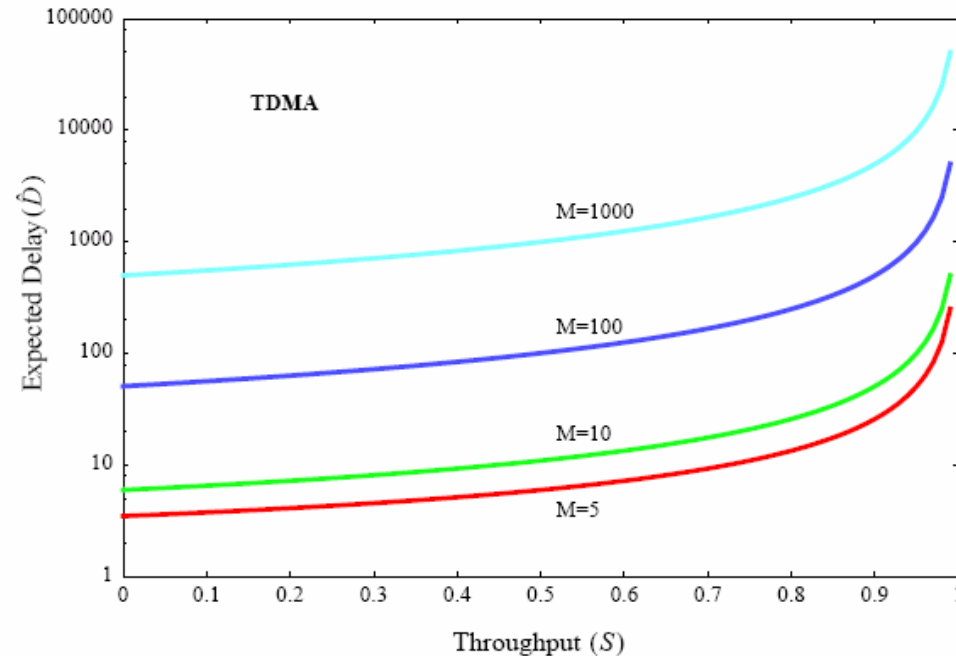
TDMA Performance

System model:

- M independent queues
- $M/D/1$ queueing system

Throughput-Delay

$$\hat{D} = \frac{D}{T} = 1 + \frac{M}{2(1-S)}$$



Comparison between FDMA and CDMA

$$D_{FDMA} = D_{TDMA} + \frac{P}{R} \left[\frac{M}{2} - 1 \right] \geq D_{TDMA}$$



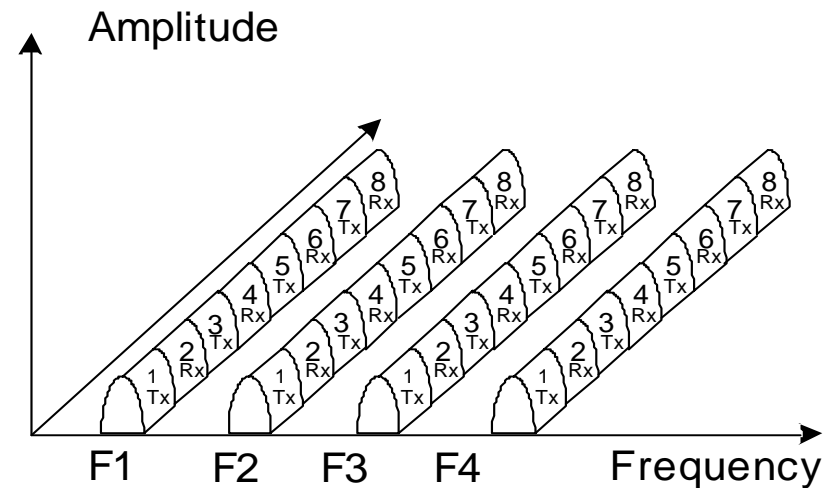
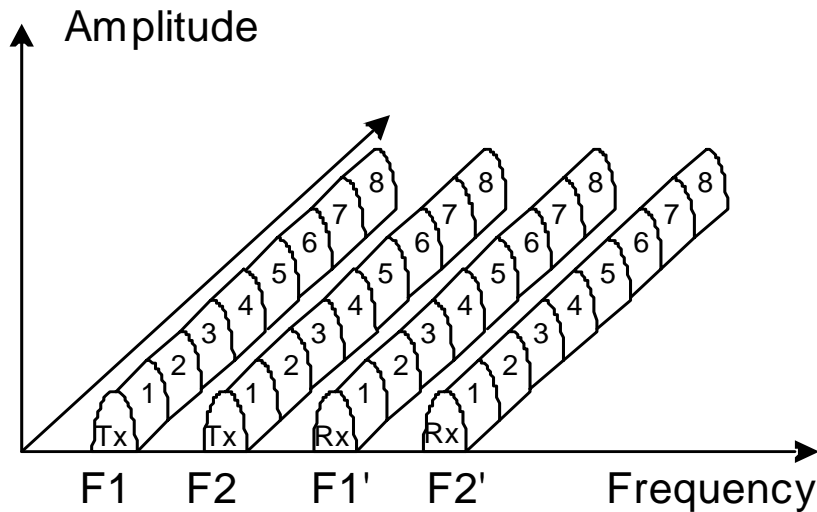
Hybrid of TDMA and FDMA

✚ TDMA and FDMA can be implemented together to get optimized function and performance

✚ For example

☐ *TDMA/FDMA*

☐ *TDMA/TDD/FDMA*



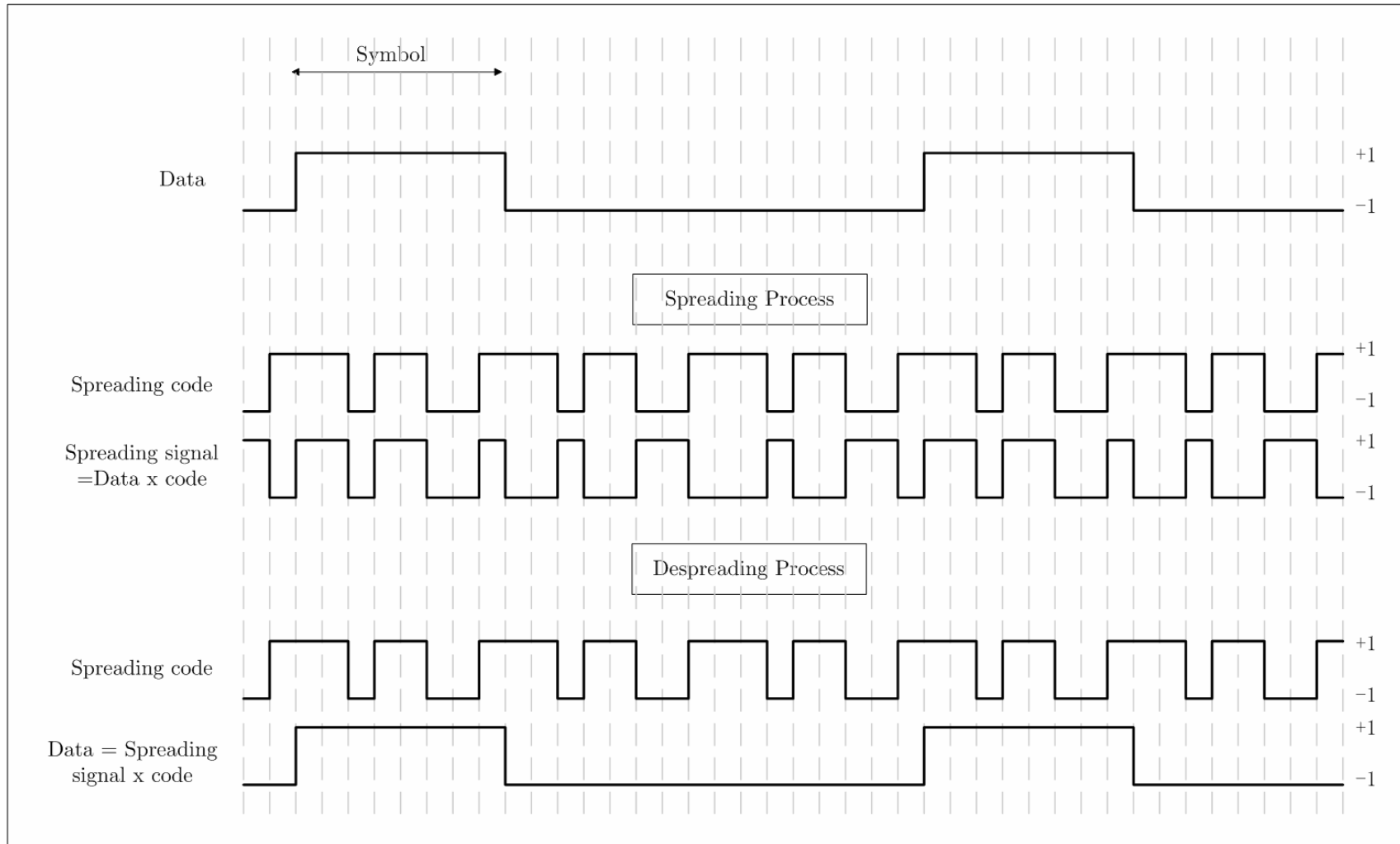


Code-Division Multiple Access

- ✚ Hybrids combination of FDMA and TDMA
- ✚ Characteristics
 - ☞ Multiple users
 - ☞ Simultaneously operating
 - ☞ Entire bandwidth of time-frequency domain
 - ☞ Separated by distinct user-signal codes (Spread spectrum)
- ✚ Two common CDMA
 - ☞ Direct sequence (DS) CDMA
 - ☞ Frequency Hopping (FH) CDMA
- ✚ Spreading code
 - ☞ ML (Maximum length) code
 - ☞ Gold code
 - ☞ Walsh-Hadamard code

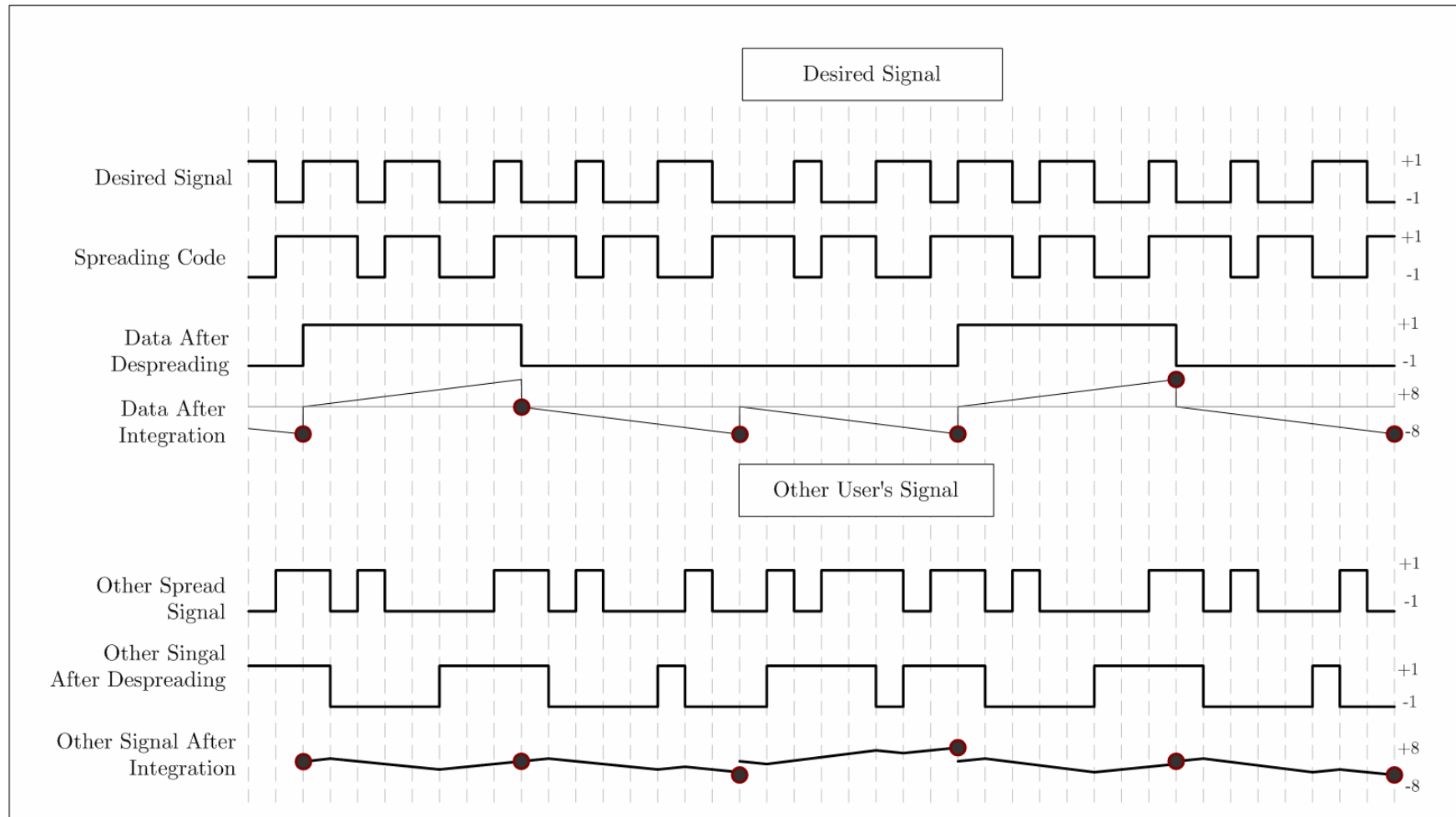


CDMA --- Spreading





CDMA --- De-spreading





Dynamic Allocation Protocols

- ✚ Dynamic Allocation = **On demand**
- ✚ Realized via **reservation** schemes
- ✚ Basic access methods:
 - ☐ Polling
 - ☐ Token passing
- ✚ Others
 - ☐ MSAP
 - ➔ *MiniSlotted Alternative Priority*
 - ☐ BRAM
 - ➔ *Broadcast Recognition Access Method*



Polling Techniques (1)

Centralized control

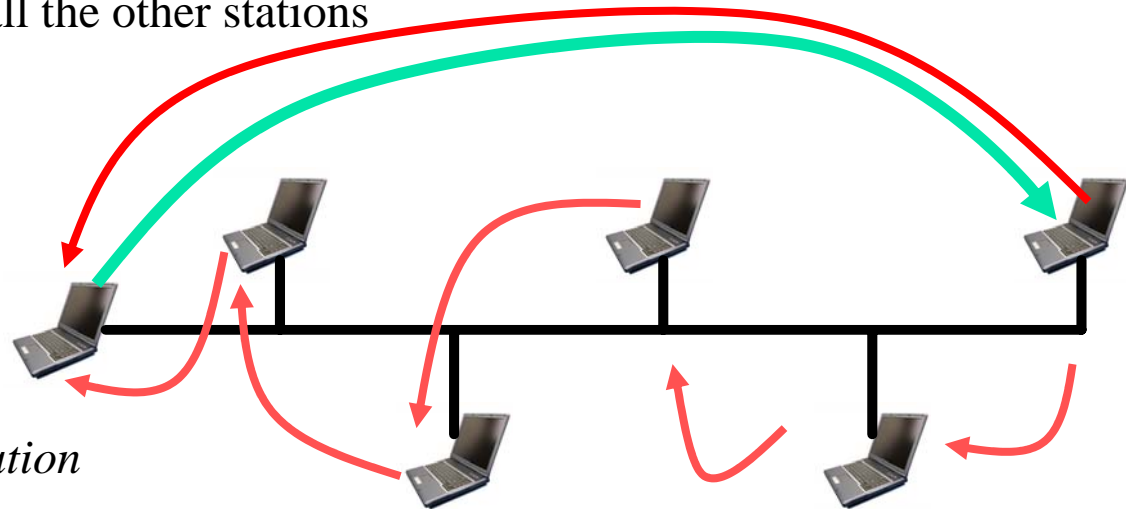
- One station is equipped as a controller
- Periodically polling all the other stations

Classification

- Hub polling
- Roll polling

Polling procedure

- Hub polling
 - From furthest station
 - Polled station starts sending if it has something to transmit
 - If not, a negative response is detected by the controller
 - The polled station transmits the poll message to its neighbor in upstream (control)
 - Control message finally is regained by the controller

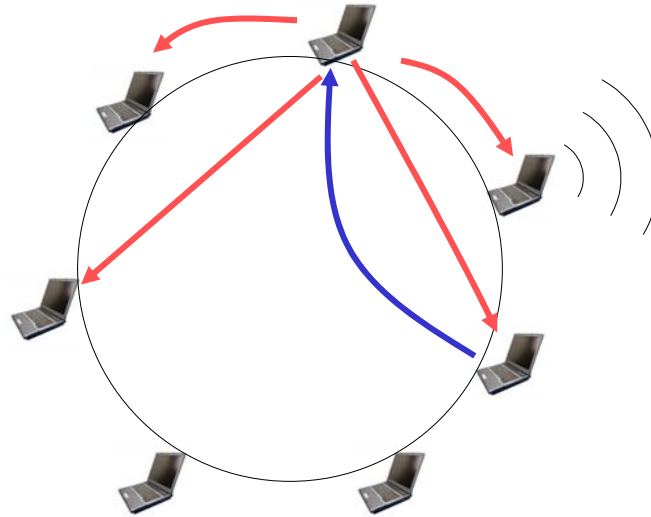




Polling Techniques (2)

Roll polling

- Controller keeps a polling list, giving the order in which the terminals are polled
- Polled station starts sending if it has something to transmit
- If not, a negative reply is detected by the controller
- Controller then polls the next terminals in the sequence
- Initial exchange of short messages required (between a station and the controller)





Characteristics and performance

- ✚ Polling protocols are efficient in systems
 - ☞ Propagation delay is small
 - ☞ Overhead is low
 - ☞ Number of stations shouldn't be large (proportional to overhead)
- ✚ Polling protocols are inefficient
 - ☞ Lightly loaded
 - ☞ Part of stations have data to transmit
 - ☞ Subdivide stations into subsets (variations)
- ✚ Hub polling overhead is much smaller than that of roll polling
- ✚ Applications:
 - ☞ Widely used in dedicated telephone networks for data communications
 - ☞ Generally not been adopted in existing mobile data network or WLAN



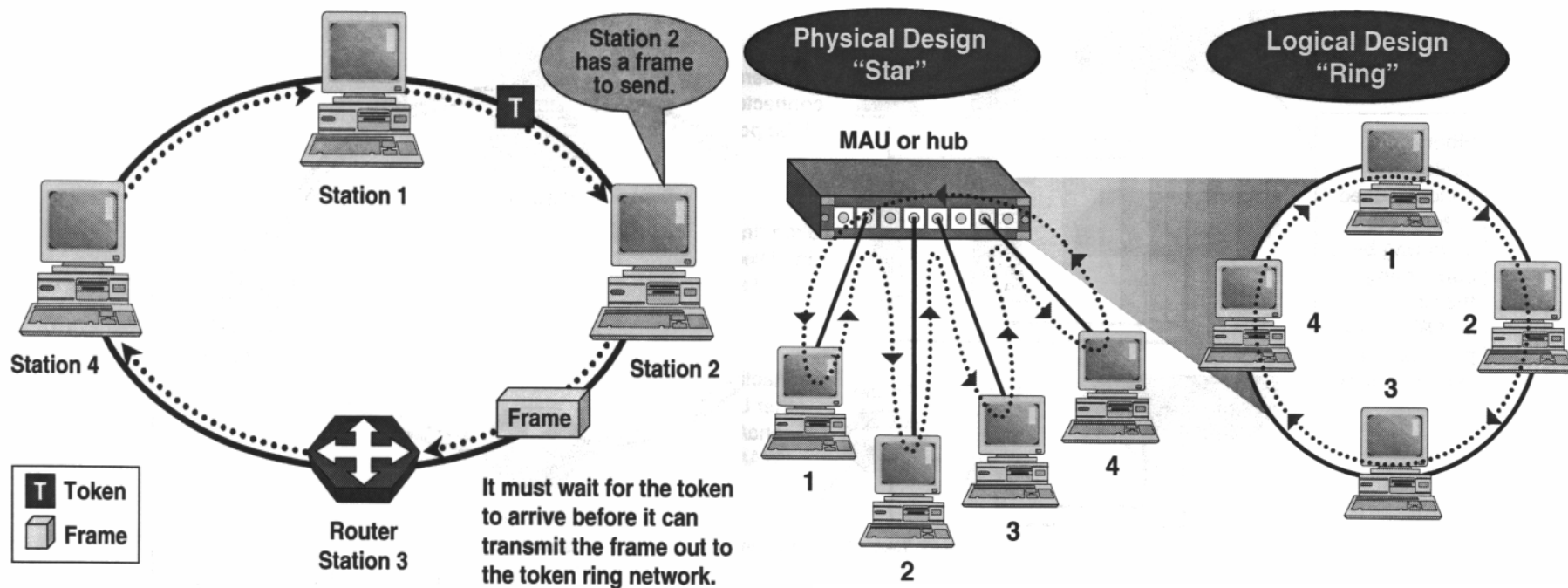
Token Passing Protocol (1)

- ✚ Two logical topologies
 - ☐ Bus
 - ☐ Ring
- ✚ Token ring is originally developed by IBM, specified in IEEE 802.5
- ✚ Token ring protocol operation
 - ☐ Networks move a small frame, called a token, around the network
 - ☐ Possession of the token grants the right to transmit.
 - ☐ If the node, receiving the token, has no information to send, it passes the token to the next station.
 - ☐ If the node, possessing the token does have information to transmit
 - ➔ *Seize the token*
 - ➔ *Alter 1 bit of the token*
 - ➔ *Append the information to be transmitted, and send to the next station in the ring*
 - ☐ The intended destination station flips the recognized address and frame-copied bits in frame status field in the frame, and sends the modified frame back out to the ring



Token Passing Protocol (2)

- When information reaches the sending station again, it examines and removed the frame from the ring
 - The source station then transmits a new token
- + Physically "star" topology, logically "ring" topology





Token Ring Characteristics

- ✚ Token passing networks are deterministic, so the maximum propagation time is possibly calculated, more predictable than Ethernet
- ✚ Priority schemes can be deployed to improve the efficiency
 - ☞ User-designated, high priority station can use network more frequently
 - ☞ $\text{Priority}_{\text{Station}} \geq \text{Priority}_{\text{token}}$ can capture the token
- ✚ Several mechanisms for detecting and compensating for network fault
 - ☞ One station is selected as *active monitor*
 - ☞ It provides centralized source of timing information for other stations
 - ☞ Ring-maintenance function
 - ➔ Removal of continuously circling frames
 - ➔ Generation of the new token
- ✚ No collisions occur, contention-free!



Contention Protocols

✚ No guarantee to be successful

- ☞ Large users ->Contention-free schemes are impractical
- ☞ Resolution schemes are needed

✚ Static resolution

- ☞ Protocol actual behavior is not influenced by the dynamics of the system
- ☞ Examples:
 - ➔ *Aloha family*
 - ➔ *CSMA family*

✚ Dynamic resolution

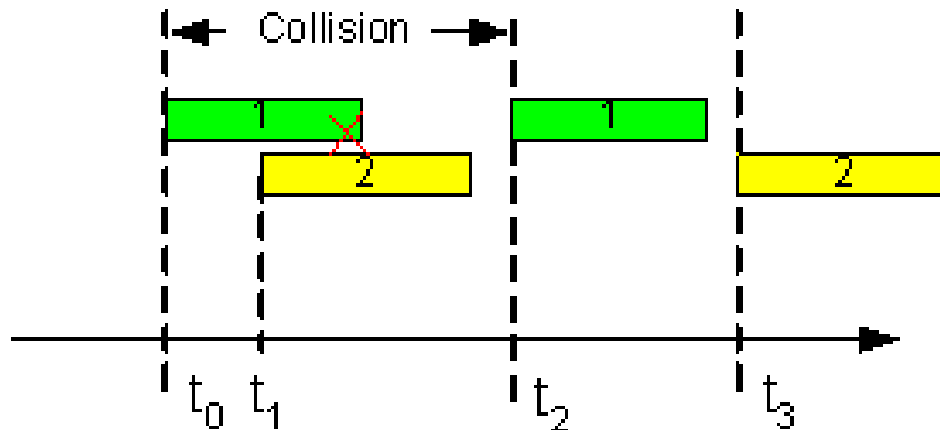
- ☞ Tracking and taking the advantage of the system changes
- ☞ Example:
 - ➔ *Binary-Tree CRP (collision Resolution Protocol)*



Pure Aloha

- ✚ It is the simplest contention protocol
- ✚ Whenever packet needs transmission:
 - 📄 Send without waiting
 - 📄 If collision occurs, then wait for a random time and resend, until successful

$$S = Ge^{-2G}$$

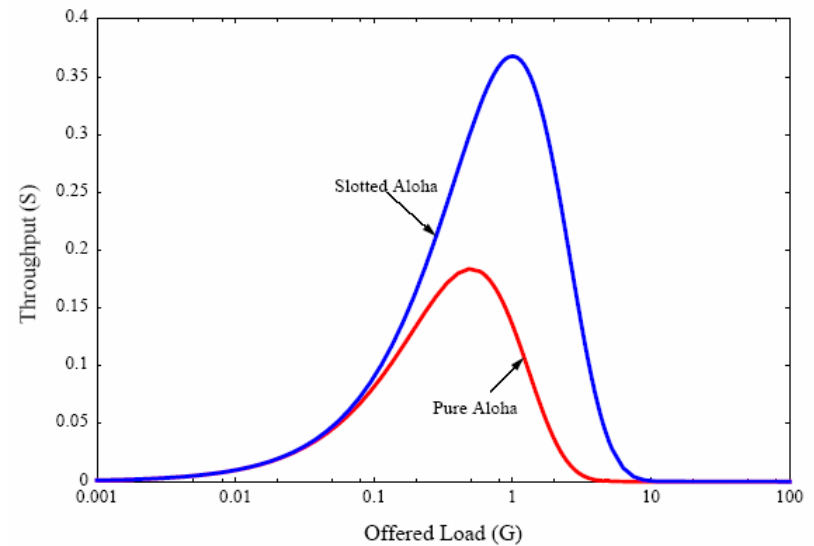
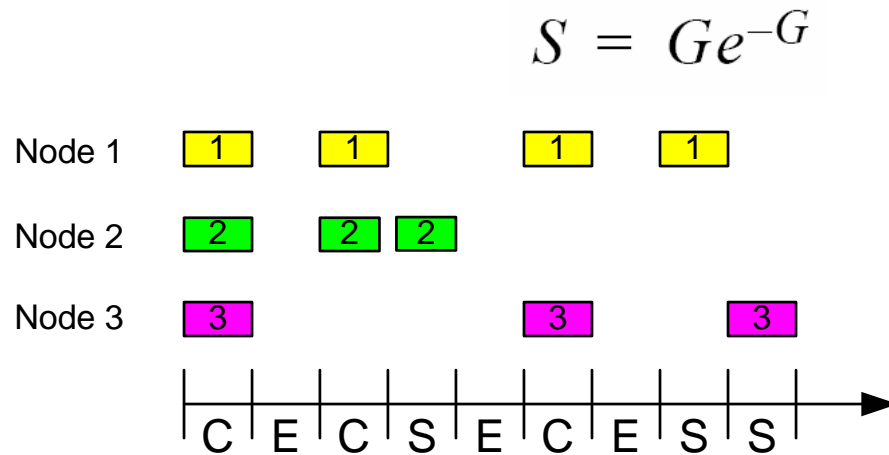


- ✚ Highly inefficient at large loads. Maximum utilization of 18% at a mean load of 0.5



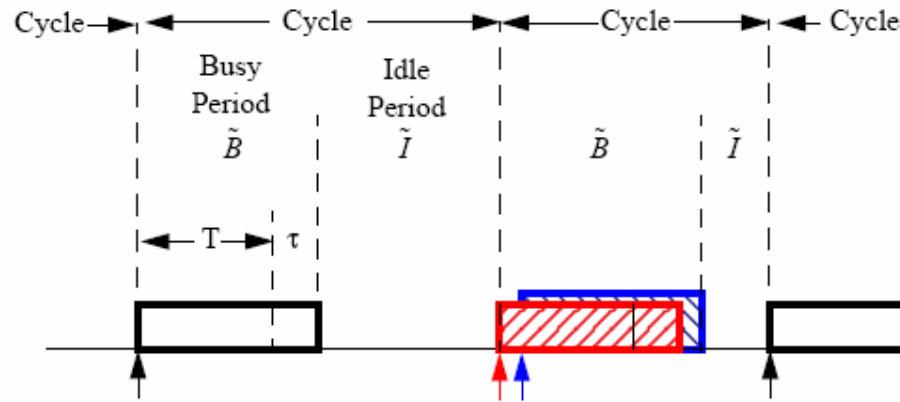
Slotted Aloha

- ✚ Time is divided into equal size slots (= packet Tx time)
- ✚ Node with new arriving packet: transmit at beginning of next slot
- ✚ If collision: retransmit packet in future slots with probability p , until successful.
- ✚ Maximum utilization of 36% at a mean load of 1 transmission/slot





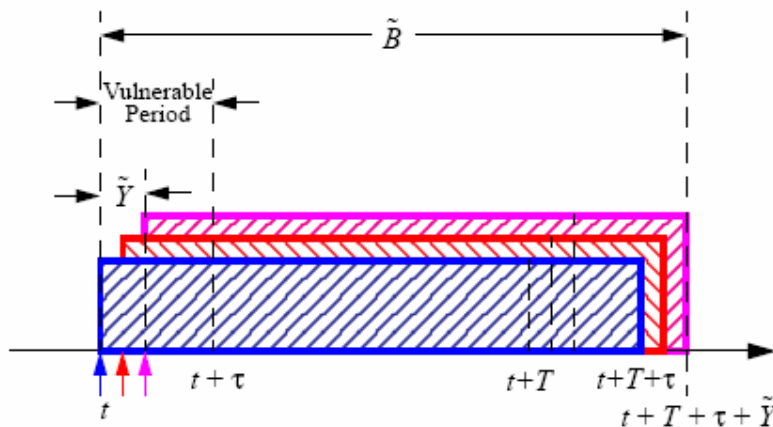
Nonpersistent (NP) CSMA



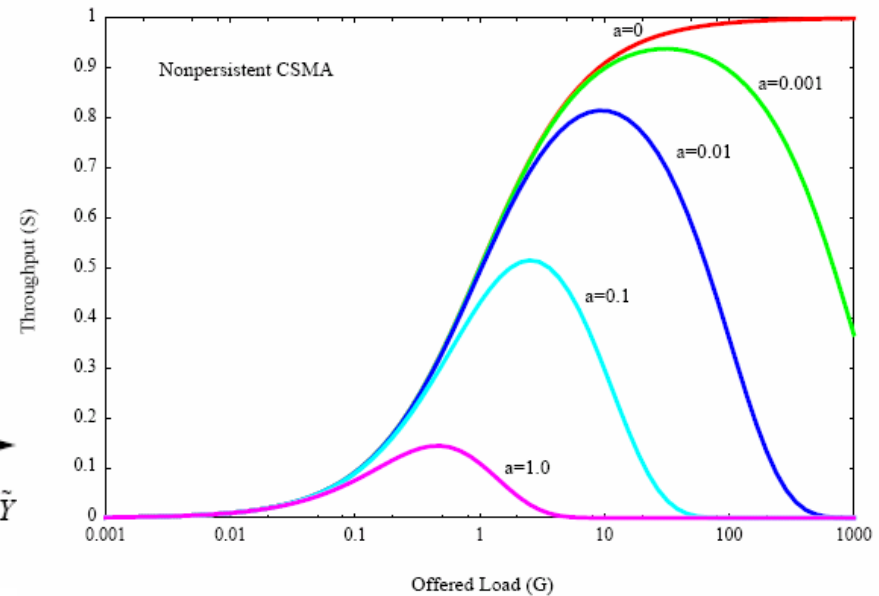
(a) Cycle Structure

Impolite -> "listen before talk"

$$S = \frac{Ge^{-aG}}{G(1+2a) + e^{-aG}}$$

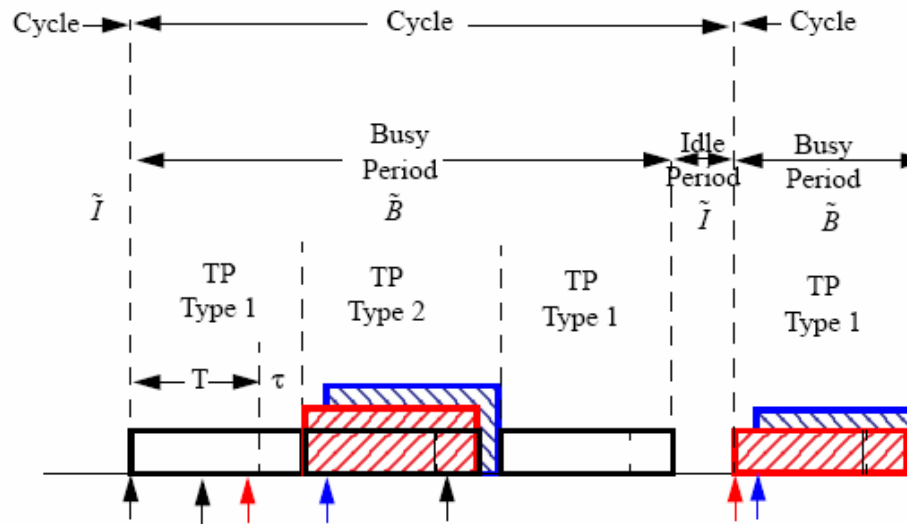


(b) Unsuccessful Transmission Period

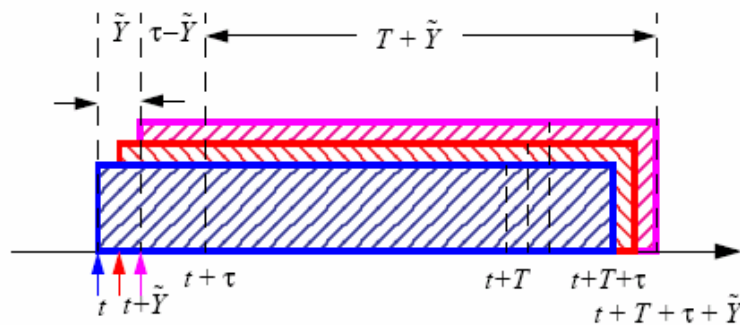




1-persistent CSMA



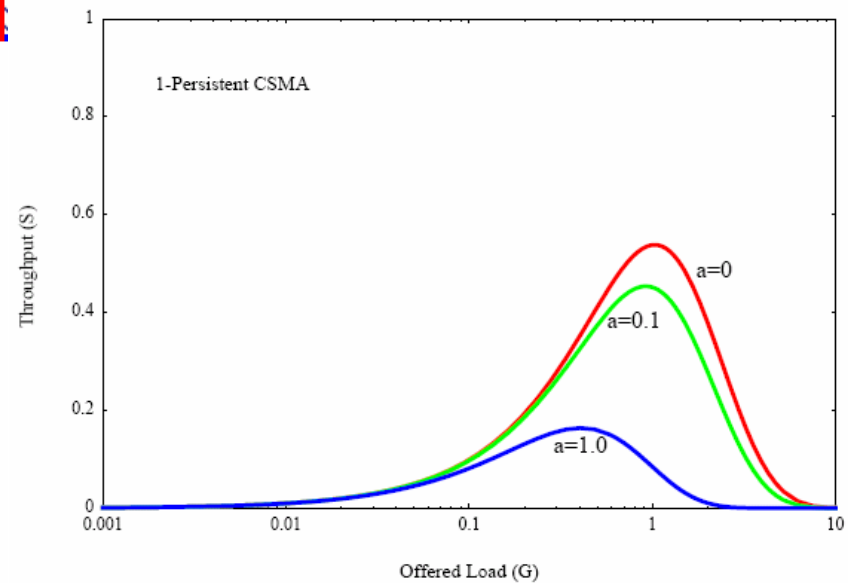
(a) Cycle Structure



(b) Unsuccessful Transmission Period

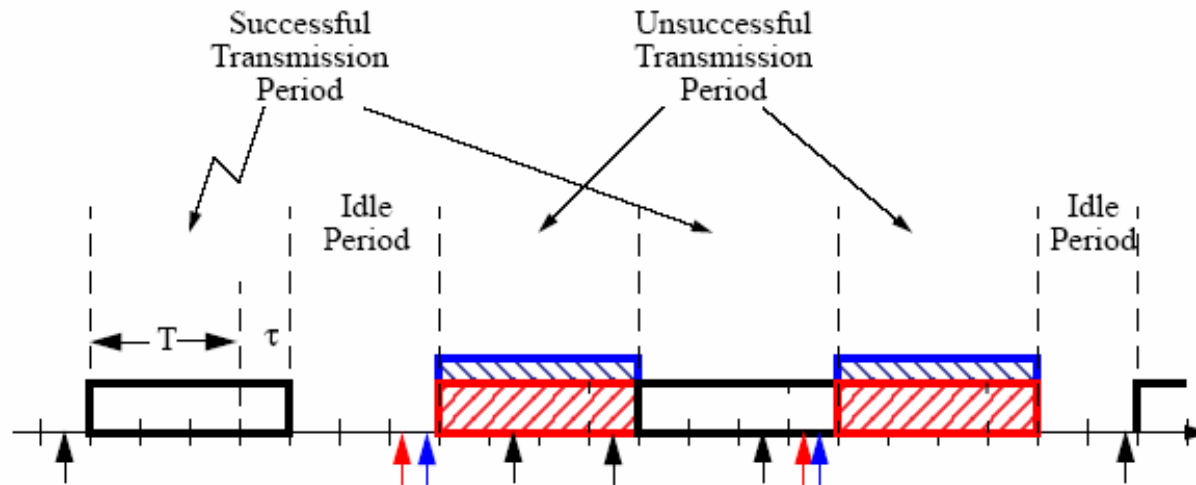
Transmit as soon as when it detect the channel is idle

$$S = \frac{Ge^{-G(1+2a)}[1 + G + aG(1 + G + aG/2)]}{G(1 + 2a) - (1 - e^{-Ga}) + (1 + Ga)e^{-G(1+a)}}$$

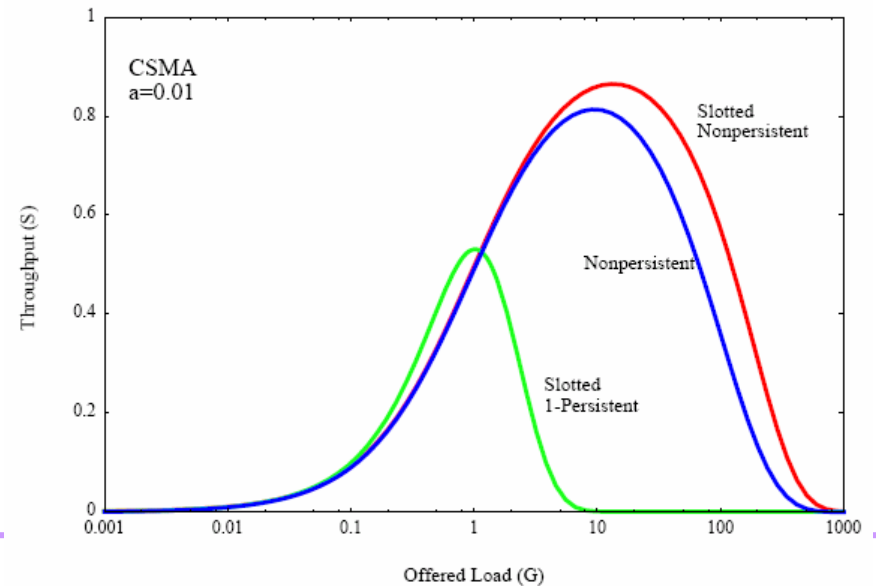




Slotted Carrier Sensing Protocols



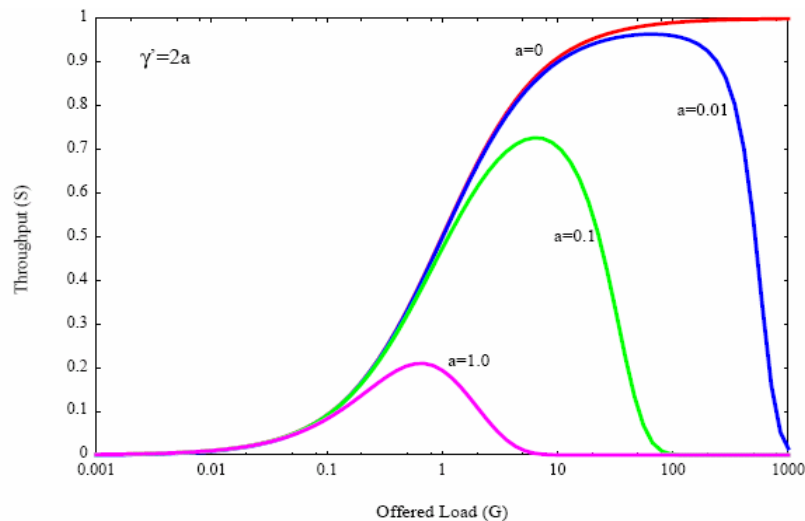
$$S = \frac{Ge^{-(1+a)G}[1 + a - e^{-aG}]}{(1+a)[1 - e^{-aG}] + ae^{-(1+a)G}}$$



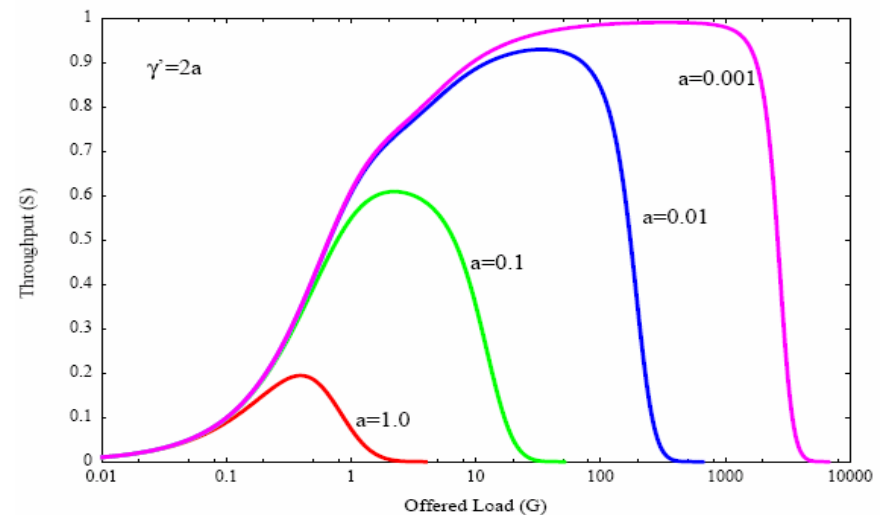


CSMA/CD

- Carrier sense multiple access with collision detection
- Same as CSMA except a collision is detected
 - Terminate transmission immediately
- Time slot and non/1-persistent concepts can also be applied



Slotted nonpersistent CSMA/CD



Slotted 1-persistent CSMA/CD

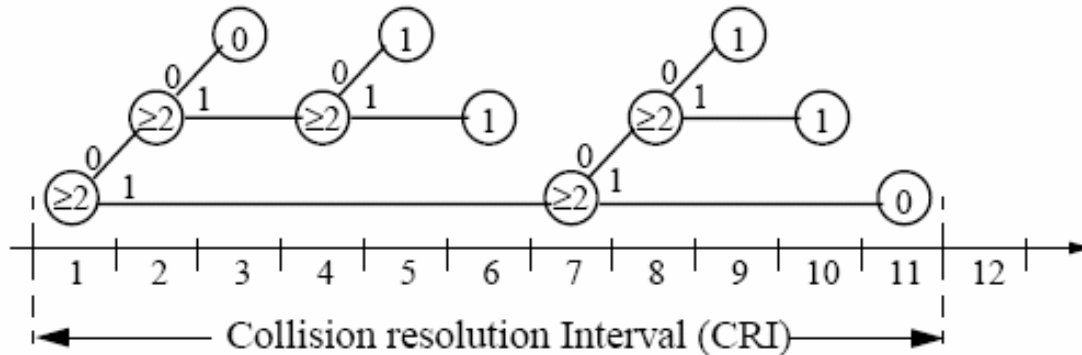


Collision Resolution Protocol (CRP)

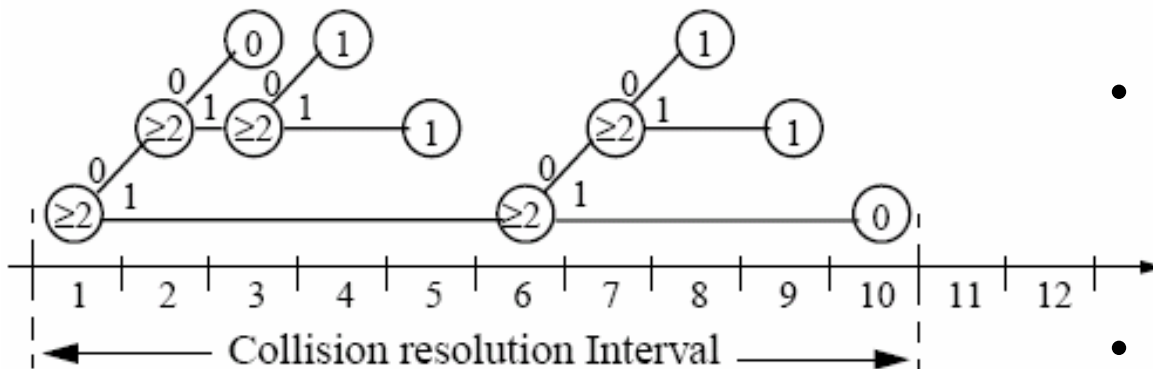
- ✚ Concentrate on **Resolving Collisions** as soon as they occurs
- ✚ Exploit feedback information to **control the retransmission**
- ✚ System model:
 - ☐ Similar to slotted Aloha
- ✚ Typical protocol:
 - ☐ Binary-Tree CRP



Binary-Tree Protocol (CRP)



Binary-Tree CRP

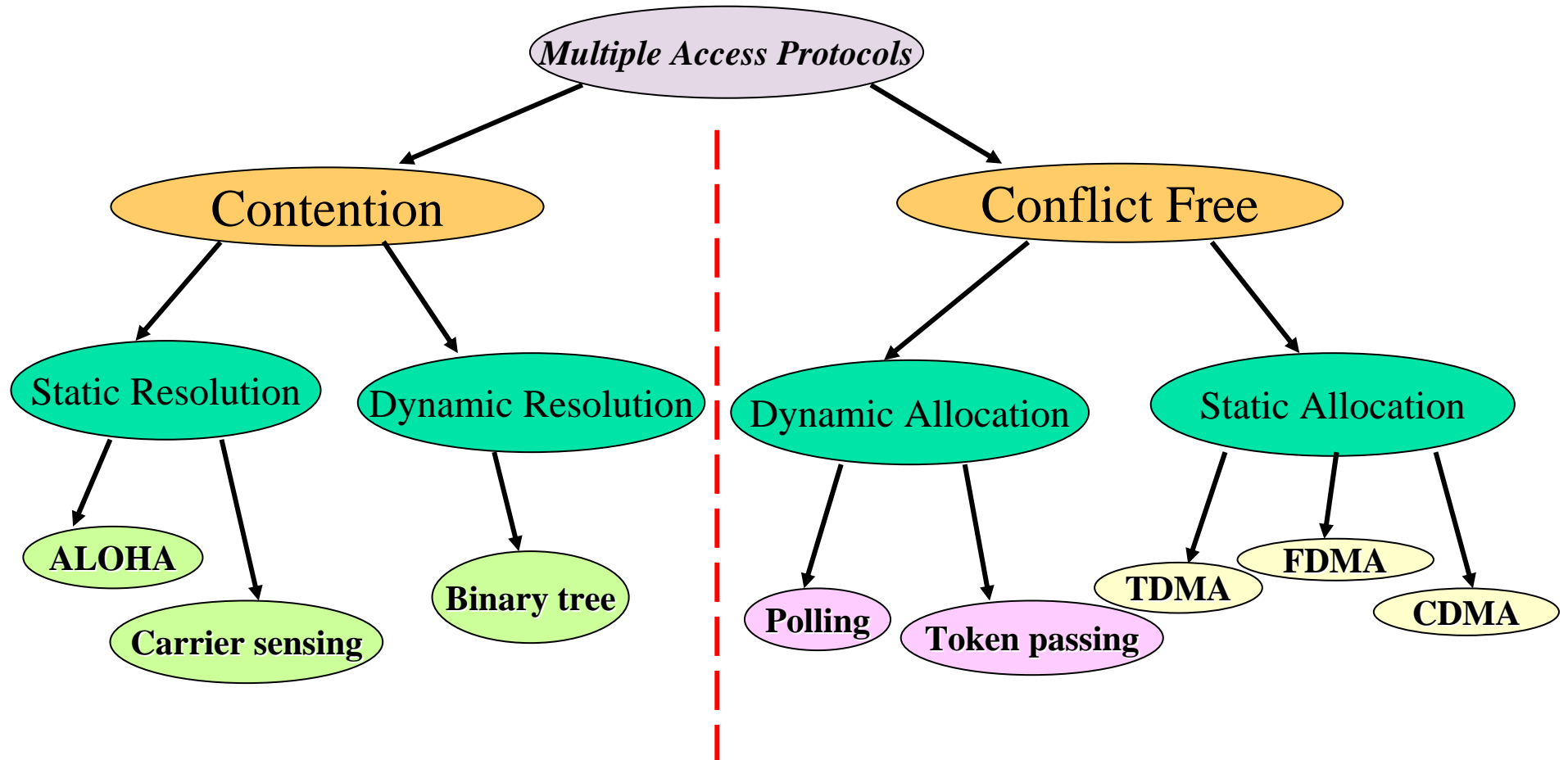


Enhanced Binary-Tree CRP

- When collision occurs, the users are divided in two subsets
- Divided method can be like flipping coin
- One subset should wait until the other set has finished transmission
- Within one subset, if collision occurs again, performance in the similar way again
- Enhanced method: "flipping beforehand"



Conclusion





Reference

[1] Raphael Rom, Moshe Sidi, “ Multiple Access protocols: Performance and Analysis”

www-comnet.technion.ac.il/rom/PDF/MAP.pdf

[2] Simon Haykin, Michael Moher, “Modern Wireless Communications”
ISBN 0-13-124697-6, Prentice Hall 2005



Homework

- ✚ Please explain what is bit-map protocol? What's the advantage and disadvantage of this protocol?

- ✚ What is FDDI, how it works. Please give a brief description on its MAC protocol structure and operation.



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Any questions?

Thanks!