Optimal load balancing algorithm for multi-cell LTE networks

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Abstract: Long Term Evolution (LTE) is a promising option for the 4th generation communications because of its higher data rates, lower latency and larger coverage. However, in a multi-cell LTE network, the network performance may be deteriorated by load imbalance. The unbalanced load among multiple cells leads to a higher delay and a higher packet drop rate in the over-loaded cell, or an underutilization of resources in the under-loaded cell. In order to solve this problem, we propose a practical load balancing algorithm to find the optimal handover operations between the overloaded cell and possible target cells. The simulation results demonstrate that the proposed algorithm can reduce network overload and increase the network bandwidth efficiency.

Keywords: Smart Grid, LTE, Load Balancing, Resource Allocation

1. Introduction

The 3rd Generation Partnership Project (3GPP) Long Term Evolution (LTE) is a promising option for the 4th generation communications because of its higher data rate, lower latency and larger coverage. The LTE Release 9 shows that LTE provides up to 300Mbps download rate and 75Mbps upload rate [1]. The specification also defines Orthogonal Frequency Division Multiple Access (OFDMA) as the access technique for the downlink and Single Carrier FDMA (SC-FDMA) for the uplink [2]. OFDMA shows robustness against multi-path fading, high spectral efficiency, and bandwidth scalability, and SC-FDMA enables users to save energy. The additional crucial technique applied in LTE is Multiple-Input-Multiple-Output (MIMO) that uses multiple transmitters and receivers to achieve a higher bit rate and an improved coverage [3].

In a multi-cell LTE network, it is critical to balance the load among the neighboring cells. The LTE network performance may be deteriorated by load imbalance among multiple cells. Figure 1 shows a network consists of multiple cells, each of which is controlled by the LTE base station, named Evolved Node B (eNB) in LTE system. Each user can receive the signal from more than one eNBs, including one serving eNB (SeNB) and one or multiple target eNBs (TeNBs). The SeNB represents the eNB which is serving the user, while the TeNBs represent the eNBs which can reach the users but are not serving the users. An imbalanced multi-cell LTE network will suffer from an increased delay, a decreased network throughput, and even packet drops. Therefore, the LTE multi-cell load balancing problem needs to be considered carefully.

The load balancing algorithm in the wireless cellular networks aims to find the optimal handover operations between the overloaded cell and possible target cells. The users in the overloaded cell are handed over to the under-loaded cells in order to improve the overall network performance in terms of the latency and throughput. However, it is quite challenging to appropriately distribute the load among multiple cells for improved network performance.

In this paper, we propose a load balancing scheme which can adapt to the network conditions, and achieve a better network performance by appropriately distributing the load
among the neighbouring cells. The proposed scheme can select a proper SeNB among the multiple TeNBs for each user based on the load difference and the Signal to Interference plus Noise Ratio (SINR). We conducted the performance evaluation in the network simulator OPNET [4]. The simulation results demonstrated that the proposed scheme can get a lower end-to-end delay.

This paper is organized as follows. We describe the related work in Section 2. In Section 3, we present the network models. In Section 4, we present the problem formulation and the proposed load-balancing algorithm. The experiment results are provided in Section 5 and the conclusions are drawn in Section 6.

2. Related Work

LTE load balancing problem has been investigated in the literature. Viering et al. presented a mathematical framework for quantitative study of self-optimizing wireless networks for LTE system, in which a self-optimizing network algorithm was proposed to adjust the cell-specific handover thresholds for load balancing [5]. Lobinger et al. proposed a handover off set based load balancing algorithm using the parameter “cell specific offset” to force users to handover from the overload eNB to the target eNB [6]. The main goal of the proposed algorithm is to find the optimal handover offset that allows the maximum number of users to change cell without any admission rejection at the target eNB [6]. A directional cell breathing based reactive congestion control algorithm was proposed in [7], where the coverage area of a cell can be dynamically extended towards a nearby loaded cell when it is under-loaded, or shrunk towards the cell center when it is over-loaded.

3. Network Models

In this section, we describe the channel model and several network parameters including resource block utilization ratio and average resource block utilization ratio.

3.1. Channel Model

We assume that each cell knows the instantaneous signal strength sending from the users through the control signals, such as Channel State Information (CSI). We divide the time into time slots with equal length $\tau$. We assume that the received SINR at eNB keeps unchanged during a time slot. The average received SINR at the base station of cell $k$ from user $i$ at time slot $t$ is given by [8]

$$SINR_{ik}(t) = \frac{\psi_i(t)/L_{ik}(t)}{\sum_{i\neq j} \psi_j(t)\rho_j/L_{jk}(t)+N},$$

where $\psi_i(t)$ represents the transmit power of the user $i$ at time slot $t$, $L_{ik}(t)$ represents the path loss from the user $i$ to the base station $k$ taking into account the distance between them, and $N$ represents the power of Additive White Gaussian Noise (AWGN). $\rho_k$ represents the resource blocks (RBs) utilization ratio of cell $k$, which is described in Section 3.2.

The data rate $S_{ik}(t)$ at time slot $t$ can be calculated using Shannon-Hartley theorem, which is given by [9]

$$S_{ik}(t) = x_{ik}(t)\frac{B}{N}\log_2(1+SINR_{ik}(t)), \quad (2)$$

where $B$ represents the total bandwidth for the eNB, $N$ is the total number of RBs for the eNB, and $x_{ik}(t)$ represents the number of RBs allocated to user $i$ by cell $k$ at time slot $t$, which can be determined using the LTE scheduling algorithm.

Therefore, the data rate depends on the channel condition between the users and the eNB. In other words, to send the same amount of the traffic, the user with a better channel condition will consume a less number of the resource blocks than the user with a worse channel condition.

3.2. Network Parameters

In order to investigate the different load distributions of different eNBs, we define the network parameters as follows. We use RB utilization ratio $\rho_k(t)$ to denote the ratio between the number of the allocated RBs and the total number of the RBs in cell $k$ at time slot $t$. A larger $\rho_k(t)$ indicates a higher percentage of RB utilization in cell $k$, and thus a higher level of load in cell $k$. Assuming that all cells have the same number of RBs, denoted by $N$. Then, $\rho_k(t)$ for cell $k$ at time slot $t$ can be written as [10]

$$\rho_k(t) = \frac{\sum_{i} x_{ik}(t)}{N}, \quad (3)$$

where $I$ represents the set of users in the whole network. $x_{ik}(t)$ is the number of RBs that cell $k$ allocates to user $i$ at time slot $t$ which can be determined by the LTE scheduling algorithm. We assume that the length of the time slot $\tau$ is much larger than the subframe duration (e.g., 1 ms).

The average RB utilization ratio of the whole network at time slot $t$ is given by

$$\rho(t) = \frac{1}{|K|}\sum_{k \in K} \rho_k(t), \quad (4)$$

where $K$ is the set of the cells in the network, and $|K|$ represents the number of the cells in the set $K$.

Handovers will be performed in case of overload. However, during the handover operation, the total load should not exceed the capacity of the eNB. We introduce a parameter $\xi(t)$ to indicate the load balancing level of the LTE network. The level of load balancing can be evaluated by the fairness index [11], which is given by

$$\xi(t) = \frac{\left(\sum_{k \in K} \rho_k(t)\right)^2}{|K|\left(\sum_{k \in K} \rho_k(t)^2\right)}, \quad (5)$$

where the value of load balancing index $\xi(t)$ is in the range $[1/N, 1]$. A larger $\xi(t)$ indicates a more balanced load distribution among cells, and vice versa. Particularly, $\xi(t) = 1$ represents that all cells have equal load at time slot $t$. 

4. Load Balancing Algorithm

4.1. Problem Formulation

In this section, we formulate the optimization problem for load balancing in the LTE network. On the one hand, we want to use a minimal resource to send all traffic in the network, which means that we want to minimize the average RB utilization ratio \( \rho(t) \). On the other hand, we want to evenly utilize the RBs among cells, which means that we want to maximize the load balancing level \( \xi(t) \). However, \( \rho(t) \) and \( \xi(t) \) depend on each other. Reducing \( \rho(t) \) may lead to load unbalancing, while increasing \( \xi(t) \) may cause a higher consumption of RBs. Considering the trade-off between the average RB utilization ratio \( \rho(t) \) and the load balancing level \( \xi(t) \), we introduce an aggregation parameter \( z \), which is defined as

\[
    z = \delta \rho(t) - (1 - \delta)\xi(t),
\]

where \( \delta \) is a weight representing the trade-off between the RB utilization and the load balancing level. If \( \delta \) is set to 1, the RB utilization will be the objective. On the contrary, if \( \delta \) is equal to 0, the load balancing level will be the objective. When \( \delta \) is between 0 and 1, the objective is the compromised value taking into account both RB utilization and load balance.

Therefore, the optimization problem is mathematically formulated as follows:

\[
    \text{Minimize}_{x_{lk}} \delta \rho(t) - (1 - \delta)\xi(t) \quad (7a)
\]

Subject to:

\[
    \sum_{i \in I} R_{i,k}(t)x_{i,k} \leq N, \quad \forall k \in K, \quad (7b)
\]

\[
    \sum_{k \in K} R_{i,k}(t) = 1, \quad \forall i \in I, \forall k \in K, \quad (7c)
\]

\[
    SINR_{i,k}(t) \geq SINR_{th}, \quad \forall i \in I, \forall k \in K. \quad (7d)
\]

The objective function (7a) represents the combination of the RB utilization and the load balancing level with the weight parameter \( \delta \) (\( 0 < \delta < 1 \)). Constraint (7b) shows that the number of resource blocks occupied by all users in a cell should not exceed the total number of resource blocks in the cell. Constraint (7c) specifies that each user can be served by only one eNB. Constraint (7d) represents that a user’s SINR \( SINR_{i,k}(t) \) should not be lower than the SINR threshold \( SINR_{th} \) in order to ensure acceptable data communications.

4.2. Practical Load-Balancing Algorithm

The optimization problem (7) is an integer programming. If we use exhaustive search to find the optimal solution, it is not suitable for delay-sensitive applications because the processing time will be unacceptable. In this section, we propose a practical load-balancing algorithm which provides a sub-optimal but much more efficient solution to the optimization problem (7). The principle of the proposed algorithm is that the user always chooses the eNB with the highest \( z \) value as its serving eNB.

The proposed practical load balancing algorithms executed at each SeNB. The SeNB first finds the set of the users, denoted by \( S \), that are served by it, and then sorts the users in the set \( S \) in an ascending order based on their SINR values. For each user in the set \( S \), SeNB will find a better TeNB if available, and hand over the user to the chosen TeNB. The process of finding the better TeNB is as follows. First, the SeNB finds the set of the target TeNBs, denoted by \( W \), for the current user. Then, for each TeNB in the set \( W \), the SeNB calculates the number of required resource blocks if the user is handed over to the TeNB, and compares the current value \( z_{SeNB} \) with the \( z \) value of the TeNB \( z_{TeNB} \). If the difference between \( z_{SeNB} \) and \( z_{TeNB} \) is larger than a threshold \( \rho_{th} \), the user will be handed over from the current SeNB to the TeNB. The same procedure repeats until all users in the set \( S \) have been processed by the SeNB. The proposed practical load balancing algorithm is presented in Algorithm 1 as follows.

<table>
<thead>
<tr>
<th>Algorithm 1: the proposed practical load balancing algorithm executed at a SeNB</th>
</tr>
</thead>
<tbody>
<tr>
<td>1: Find the set of users, denoted by ( S ), that are currently served by the SeNB.</td>
</tr>
<tr>
<td>2: Sort the users in the set ( S ) in an ascending order based on their SINR values.</td>
</tr>
<tr>
<td>3: for each user in the set ( S )do</td>
</tr>
<tr>
<td>4: Calculate ( z_{SeNB} ) value;</td>
</tr>
<tr>
<td>5: Find the set of TeNBs, denoted by ( W );</td>
</tr>
<tr>
<td>6: Collect measurements (e.g., CSI) from the user to each TeNB in the set ( W );</td>
</tr>
<tr>
<td>7: Obtain the number of the available resource blocks at each TeNB in the set ( W );</td>
</tr>
<tr>
<td>8: for each TeNB in the set ( W )do</td>
</tr>
<tr>
<td>9: if ( SINR_{th} &gt; SINR_{th} ) then</td>
</tr>
<tr>
<td>10: Calculate the number of required resource blocks after the handover;</td>
</tr>
<tr>
<td>11: Calculate the ( z_{TeNB} ) value;</td>
</tr>
<tr>
<td>12: if ( z_{TeNB} - \rho_{th} ) then</td>
</tr>
<tr>
<td>13: Perform handover from SeNB to TeNB;</td>
</tr>
<tr>
<td>14: end if</td>
</tr>
<tr>
<td>15: if ( z_{TeNB} - \rho_{th} ) then</td>
</tr>
<tr>
<td>17: end for</td>
</tr>
<tr>
<td>18: end for</td>
</tr>
</tbody>
</table>

The proposed practical load balancing algorithm is a sub-optimal solution to the optimization problem (7). The advantage of the proposed algorithm is that it can find the near-minimum objective at a much faster speed than the exhaustive search approach. The difficulty of finding the solution to the problem (7) lies in the processing order of the users at each SeNB. User assignments affect each other. Different processing order of users will lead to different results. In the proposed practical load balancing algorithm, the SeNB sorts the users in the set \( S \) in an ascending order based on their SINR values, and processes the user with the lowest SINR value first. This method tries to transfer the user with the worst channel condition to the TeNB with a better channel condition, which can improve the overall performance. At the same time, the users with better SINR keep in relatively good channel condition, because the proposed algorithm takes the fairness level into consideration.
If a handover is performed as long as the users’ $z$ value of TeNB $z_{TeNB}$ is smaller than that of SeNB $z_{SeNB}$, it may lead to hand over ping-pong effect [12], which means that a user switches its SeNB frequently and cannot reach a stable state. In order to prevent the handover ping-pong effect, we introduce an threshold value $\rho_{th}$ in the proposed algorithm. Only when $z_{SeNB} - z_{TeNB} > \rho_{th}$, a handover is triggered. The threshold $\rho_{th}$ should be chosen carefully. A smaller $\rho_{th}$ may still cause the ping-pong effect, while a larger $\rho_{th}$ would make the proposed algorithm perform worse.

5. Simulations

In this section, we perform simulations to evaluate the network performance of the proposed algorithm in terms of the algorithm efficiency and load balancing efficiency.

OPNET Modeler [13] is a network simulation tool. It provides a comprehensive environment for modeling and simulation of deployed wired and wireless networks. OPNET Modeler enables users to create customized models and to simulate various network scenarios. The wireless module is used to create models for wireless scenarios such as Wi-Fi and LTE. The Modeler is object-oriented and employs a hierarchical approach to model communications networks. It provides graphical user interfaces known as editors to capture the specifications of deployed networks, equipment, and protocols.

5.1. Simulation Setting

We simulate a city-scale LTE network of 5km x 5km. The network topology is shown in Figure 2. The scenario consists of 2 eNBs (eNodeB_1 and eNodeB_2), 8 users (node 1 – node 8) and a server. The eNBs are evenly located in the city area. Node_1_1 to Node_1_4 are initially connected to eNodeB_1, and Node_2_1 to Node_2_4 are connected to eNodeB_2. The IP addresses assigned to Node_1_1 to Node_2_4 are 192.168.7.2 to 192.168.7.9, respectively, and the IP address of the server is 192.168.4.2. The server connects to the Evolved Packet Core (EPC), the core network of the LTE system, by a link. Each user has the same configuration in terms of the applications and the traffic volumes. The SINR threshold $SINR_{th}$ is set to 10 dB, and the system bandwidth is set to 5 MHz. We set the weight parameter $\delta$ to 0.7, which means that we pay more attention to the utilization of network resources. The topology shown in Figure 2 is used in load balancing evaluation and end-to-end delay evaluation.

Table 1 shows the traffic parameters for the OPNET simulation. Each user sends the traffic every one second. In each time period (1 second), the user sends 500 packets, which follow a Poisson distribution, and the packet size is 1024 bytes. Table 2 shows the basic LTE parameters in OPNET simulation. We assume that eNodeB_1 is the overloaded cell, and each user of eNodeB_2 sends additional traffic according to a Poisson process with sending rate increased from 1000 kbps to 6000 kbps.

Table 1. Parameter settings of traffic.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Mean value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Packet size (bytes)</td>
<td>1024</td>
</tr>
<tr>
<td>Number of packets sent once</td>
<td>500</td>
</tr>
<tr>
<td>Initialization time (seconds)</td>
<td>200</td>
</tr>
<tr>
<td>Inter-request time (seconds)</td>
<td>1</td>
</tr>
</tbody>
</table>

Table 2. Simulation parameters for LTE.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Setting</th>
</tr>
</thead>
<tbody>
<tr>
<td>Uplink base frequency</td>
<td>1920 MHz</td>
</tr>
<tr>
<td>Downlink base frequency</td>
<td>2110 MHz</td>
</tr>
<tr>
<td>Uplink bandwidth</td>
<td>20 MHz</td>
</tr>
<tr>
<td>Downlink bandwidth</td>
<td>20 MHz</td>
</tr>
</tbody>
</table>

5.2. Simulation Results

5.2.1. Comparison between the Proposed Practical Load Balancing Algorithm and the Exhaust Search Approach

Figure 3 shows the comparison of $z$ value among the proposed practical load balancing algorithm, the exhaust search approach, and the default handover approach. The red, green, blue lines represent the proposed algorithm, the exhaust search, and the default handover, respectively. In LTE default handover approach, the user switches to the other eNB with the best channel condition as long as the SINR of the user is larger than the SINR threshold. The exhaust search approach finds the optimal solution to the optimization problem (7) at the price of extremely high computational complexity. The proposed practical load balancing algorithm is an efficient and lightweight algorithm, which performs close to the globally optimal result, as demonstrated from the small performance gap between the proposed algorithm and the exhaust search approach in Figure 3. We can see that the exhaust search approach and the proposed algorithm get much lower $z$ values than the default handover approach. The smaller $z$ values indicate that the proposed algorithm and exhaust search approach consume less resource and provide more balanced load for the LTE network than the default handover approach.
5.2.2. Load Balancing Evaluation

Figure 4 shows the comparison of average RB utilization ratio between the default handover approach and the proposed algorithm. The blue line represents the proposed scheme, while the red one represents the default handover scheme. With the increased sending rate, the average RB utilization ratios in both schemes go up. However, the average RB utilization ratio in the proposed algorithm increases slowly than the default handover scheme. The additional traffic injections into the four nodes (from Node_1_1 to Node_1_4) make the SINR of cell 1 worse. A worse cell would use more RBs to send the same amount of data. Through the handover in the proposed algorithm, the channel conditions of the users become better than before. The switched users with better SINR consume less resource, thus leading to a lower average RB utilization ratio compared to the default handover scheme.

Figure 5 shows the comparison of load balancing ratio between the proposed algorithm and the default handover scheme. It is observed from Figure 5 that the load balancing ratio for the default handover scheme decreases significantly when the sending rate is increased. The proposed algorithm achieves a higher load balancing ratio (e.g., a more balanced load distribution) than the default handover scheme. The reason is that the proposed algorithm can appropriately transfer a part of users from the over-loaded cell to the under-loaded cell to balance the load.

5.2.3. End-to-end Delay Evaluation

Figure 6 illustrates the end-to-end delay between UEs and eNBs. As shown in the figure, all traffic starts after around 100 seconds. After several seconds, the values of end-to-end delays become steady. In eNodeB_1, the values remain at around 0.025s, while in eNodeB_2, the delay values vary between 0.024s and 0.028s. That is due to different distances between the UEs and eNodeB, which leads to different channel conditions.

In order to evaluate the load distribution among cells, additional traffic is added to Node_2_1, Node_2_2, Node_2_3, and Node_2_4. The additional traffic is injected at the 3rd minute from the beginning, and we observe the change of end-to-end delay before and after the traffic injection. Figure 7 shows the end-to-end delays of the UEs after inserting the traffic flow. As shown in the Figure 7(a), the UEs’ delay at eNodeB_1 does not have too much change compared with the value before the traffic injection. Because the two cells are separated, the traffic injection in cell 2 does not affect the cell-1 network. Figure 7(b) shows the delays in default hand-over scheme, the delay values of all four nodes connected to eNodeB_2 start to increase from the third minute, and finally go above 0.04 second. The traffic injection makes the channel condition of Node_1_1, Node_1_2, Node_1_3 and Node_1_4 worse. However, the SINR values do not reach the thresholds. Therefore, the UEs do not take any hand-over action, so that the end-to-end delays are increased.

We implemented the proposed algorithm into the OPNET experiments. Figure 8 shows the delay values after implementing the proposed practical load balancing algorithm. When the additional traffic flow is added to eNodeB_2, the delay values of Node_2_1 to Node_2_4 start soaring, as shown in Figure 8(b). The increased delays are caused by the overload in eNodeB_2 when a large amount of traffic is inserted into it. When the time reaches 4 minute 20 seconds, the delay values stop increasing and start going down, as shown in Figure 8(b). This phenomenon occurred due to the fact that the proposed algorithm was triggered to hand over the UEs with worse channel conditions to the new SeNB, thus leading to delay reduction.
In this case, Node_2_3 switches its SeNB from eNodeB_2 to eNodeB_1. After the handover, the end-to-end delays become acceptable. Node_1_1 to Node_1_4 have fluctuations at this time which are caused by the handover. With more traffic injected, the end-to-end delays soar again. At the time of 5 minute 20 seconds, the SeNB of Node_2_1 is also switched to eNodeB_1, which leads to another delay reduction in Figure 8(b).

In summary, Node_1_1 to Node_1_4 are connected to eNodeB_1, and Node_2_1 to Node_2_4 are connected to eNodeB_2. Due to the additional traffic injection to eNodeB_2, Node_2_3 was handed over from eNodeB_2 to eNodeB_1 at time of 260 seconds, and Node_2_1 was handed over from eNodeB_2 to eNodeB_1 at time of 320 seconds. The handover is triggered by the proposed algorithm.

6. Conclusion

In this paper, we investigated the LTE load balancing problem. We proposed a practical load balancing algorithm for LTE networks to increase the load balancing ratio. We conducted experiments in OPNET. The experiments results demonstrated that the proposed algorithm can lead to a better load balance than the default handover scheme.

References


