

Measurement Study of Low-bitrate Internet Video Streaming

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In Proceedings of ACM SIGCOMM Workshop
on Internet Measurement
November 2002



Introduction

- Many studies of Internet performance
 - Paxson, Mogul, Caceres...
 - Across countries, many sites
 - Well-connected (often schools on backbone)
- But few look at it from the point of dialup user
 - About 50% of home users dialup
 - Peak, but will remain majority for 3-5 years
 - ISP cannot always do 56 kb/s



Introduction

- Most studies involve TCP
 - 90-99% of traffic on Internet is TCP
- But MM prefers UDP
 - (Why?)
- Also, TCP uses ACK-based scheme
 - MM protocols prefer NACK to scale (why?)
- Video studies have done few paths



Introduction

- Video streaming experiment
 - Seven month long
 - MPEG-4 (low-bandwidth) over UDP
 - Over dialup
 - 600 major cities
 - 50 States



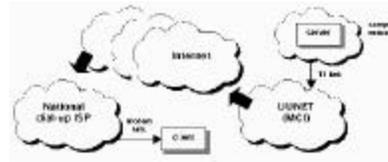
Outline

- Introduction (done)
- Methodology ←
 - Setup
 - Streaming
 - Client-Server architecture
- Results
- Analysis
- Summary



Setup

- Clients connected to each long-distance
- Server was in NY



- 3 ISPs in all 50 states
- 1813 different access points
- 1188 major U.S. cities



Setup

- Dialer
 - Connect to ISP with PPP connection
 - *traceroute* from sender → receiver and receiver → sender
 - Parallel paths
- Detect when modem connection was bad
 - If r is target bitrate, p is packet loss
 - If $B_p < 0.9r$ then bad (toss)
 - If B_p is $> 15\%$ then bad (toss)
- Good data was time-stamped
 - Day of week plus 3 eight hour slots each data
 - At least one from each day for each state for each slot



Streaming

- MPEG-4 stream
 - 2 ten-minute QCIF (176x144) streams
 - S_1 14 kbps (Nov-Dec 1999)
 - S_2 25 kbps (Jan-May 2000)
- Server split into 576 byte packets
 - With overhead S_1 16 kbps and S_2 27.4 kbps
 - About 6 packets/sec (for S_2)
- To remove jitter, had delay buffer
 - (*What is this?*)
- Chose 2.7 seconds (1.3 ideal in pilots stud)
 - (*Why might this be a bad idea?*)



Client-Server Architecture

- Server
 - Multithreaded
 - Bursts of packets (340-500 ms)
- Client
 - Recover lost packets through NACK
 - Collect RTT delay
 - Based on NACK
 - (*When might this not work well?*)
 - Probes every 30 seconds if loss $< 1\%$
- Evaluated for 9 months
 - Whew!



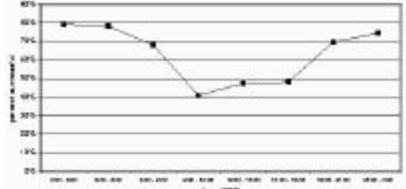
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Results

- Two datasets
 - D_1 16,783 connections, 8429 successful
 - D_2 17,465 connections, 8423 successful
 - To get MPEG-4, need 2 attempts on avg



- Time of day matters

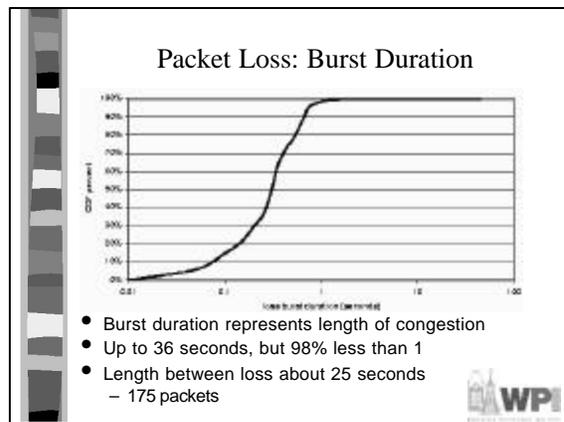
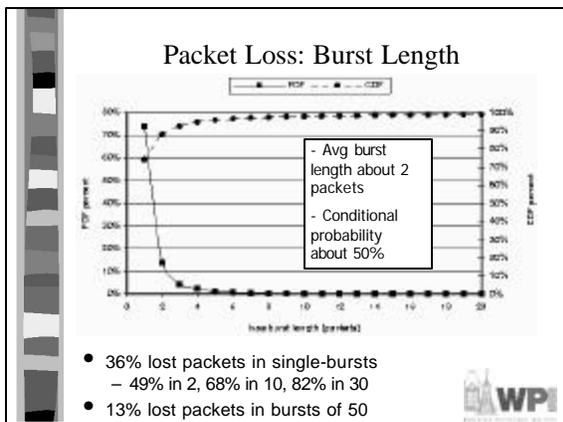
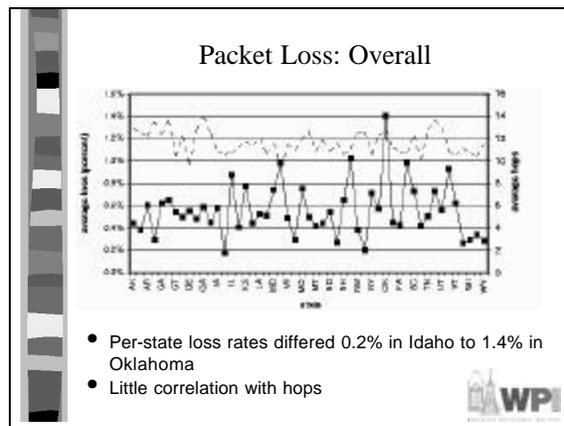
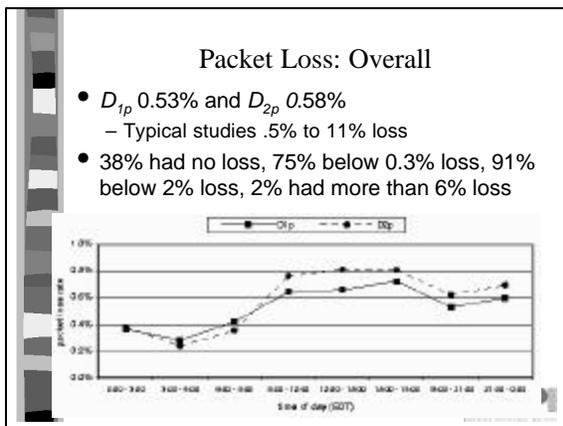
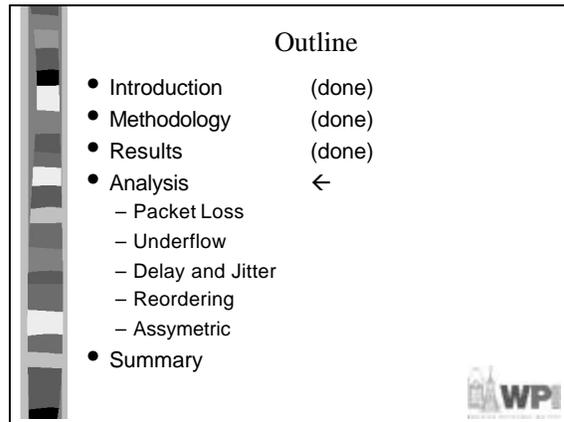
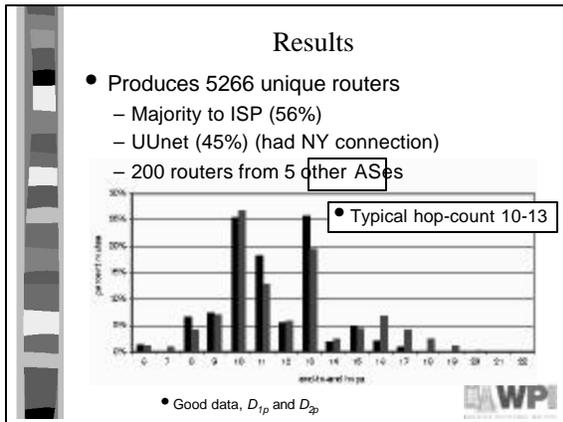


Results



- D_1 had 962 dialup points, 637 cities
- D_2 had 880 dialup points, 575 cities





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 - Packet Loss (done)
 - Underflow ←
 - Delay and Jitter
 - Reordering
 - Assymmetric
- Summary



Video Quality

- No user studies, no PSNR
 - Do not provide insight into network
- Instead, consider *underflow event*
 - When there is no frame to play
- Consider repair?
 - No standardized techniques to conceal loss
 - Techniques range from simple to complex
 - Performance depends upon:
 - Motion in video
 - Type of frame from packet (I, P, B)
 - Don't want this to be a study evaluating repair
 - Every packet loss may cause an *underflow event*



Video Quality

- Too much delay can cause underflow
 - Retransmitted packet will still be late
- Too much jitter can cause underflow
 - Retransmitted or original packet late
- Two types of late
 - Completely late (of no use)
 - Partially late (can help decode other frames in GOP)
- GOP: IPPPPPPPPP

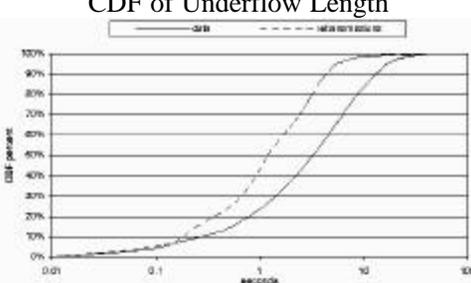


Underflow Results from Delay and Jitter

- For D_{1p} and D_{2p} , 431,000 lost packets
 - 160,000 found after deadline (37%) so no NACK
 - 257,000 (94%) sent NACK and recovered
 - 9,000 recovered late
 - 4000 (about 50%), "rescue" about 5 frames
 - 5,000 never recovered
- Jitter caused 1,100,000 underflow events
 - 98% of underflow events
 - 73% if don't use retransmission
- (*How to improve these numbers?*)



CDF of Underflow Length



Retransmit: 25% late by 2+, 10% by 5+, 1% by 10+

Jitter: 25% by 7+, 10% by 13+, 1% by 27

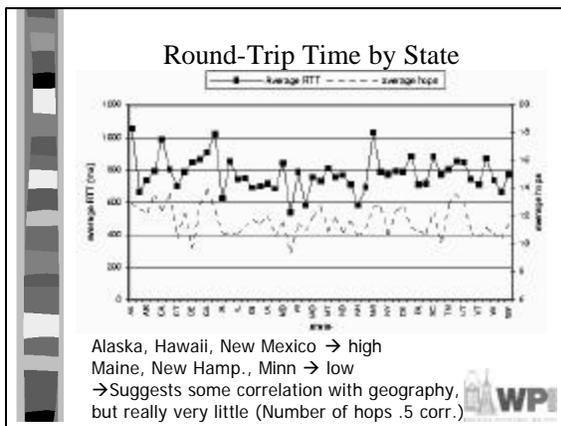
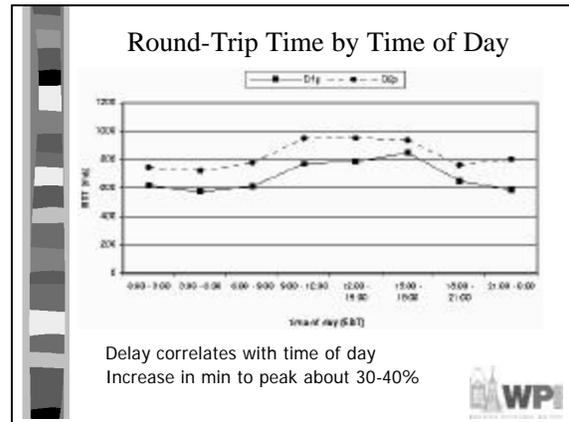
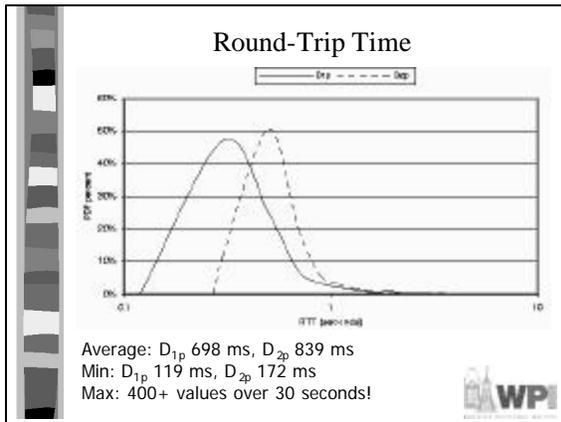
Buffer of 13 seconds would recover 99% of retransmissions and 84% of jitter



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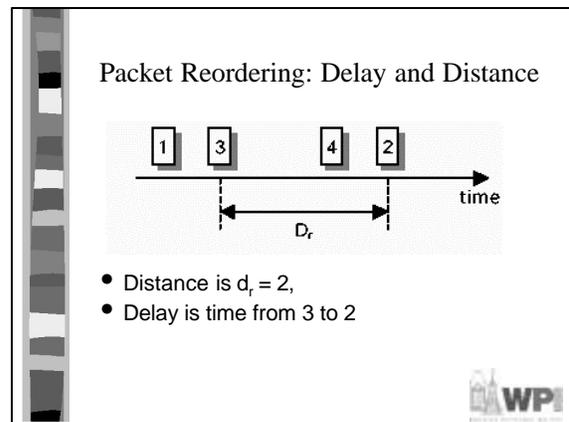
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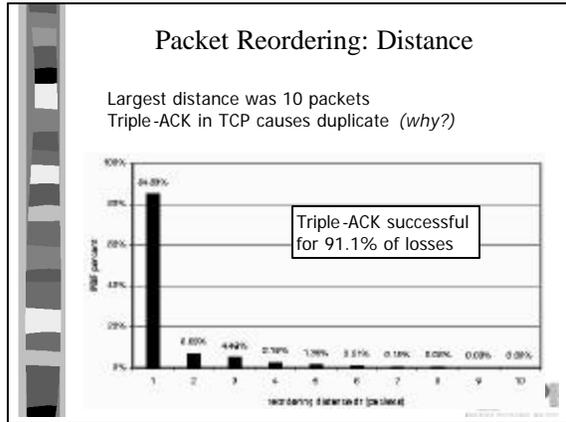
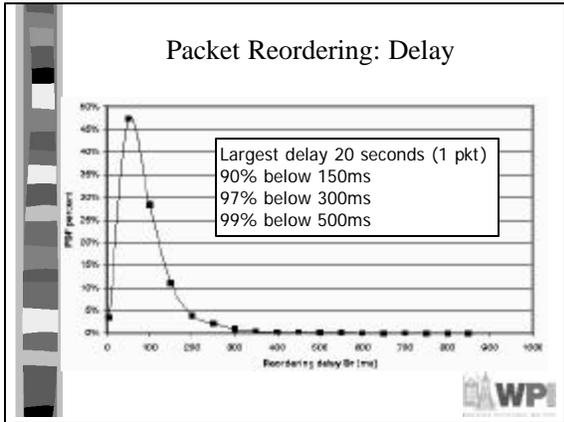




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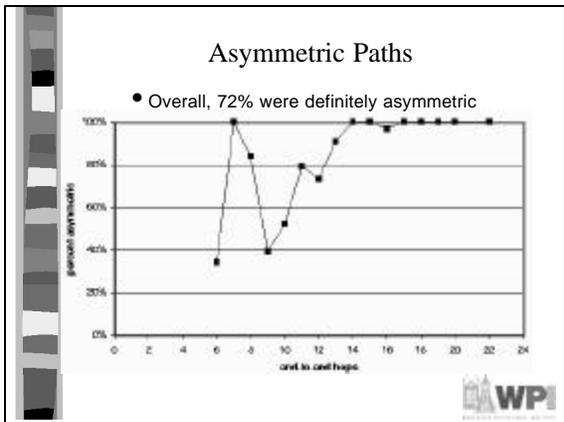
- ### Packet Reordering: Overview
- Gap in sequence numbers indicates loss
 - (When might this fail?)
 - For D_{1p}^a , 1 in 3 missing packets arrived out of order
 - Simple streaming protocol with NACK could waste bandwidth
 - Average
 - was 6.5% of missing
 - 0.04% of sent packets
 - Of 16,952 sessions, 9.5% have at least 1
 - ½ of sessions from ISP a
 - No correlation with time of day





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- ### Asymmetric Paths
- If number of hops from sender → receiver different than receiver → sender
 - then *asymmetric*
 - If number of hops from sender → receiver same as receiver → sender
 - then *probably symmetric*



- ### Conclusion
- Internet packet loss is bursty
 - Jitter worse than packet loss or RTT
 - RTTs on the order of seconds are possible
 - RTT correlated with number of hops
 - PacktlLoss not correlated with number of hops or RTT
 - Most paths asymmetric