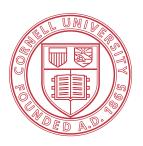
# A Hierarchy of Polyhedral Approximations of Robust Semidefinite Programs

Raphael Louca joint work with Eilyan Bitar

School of Electrical and Computer Engineering

Cornell University



# Robust Semidefinite Program

Consider the robust semidefinite program (SDP)

minimize 
$$c^{\top}x$$
  
subject to  $\sum_{i=1}^{k} \xi_i \mathcal{A}_i(x) \in \mathbf{S}^n_+, \quad \forall \ \xi \in \Xi,$ 

#### where

- $x \in \mathbf{R}^m$  is the decision variable
- $\mathcal{A}_i:\mathbf{R}^{m{m}} 
  ightarrow \mathbf{S}^{m{n}}$  is an affine function of x , and
- $\Xi \subseteq \mathbf{R}^k$  is the uncertainty set, a convex compact set

#### Some notation:

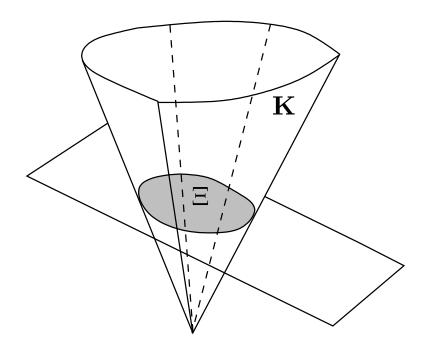
•  $\mathbf{S}^{n}\left(\mathbf{S}_{+}^{n}\right)$  : space of  $n \times n$  symmetric (symmetric PSD) matrices

# Uncertainty Model

The uncertainty set  $\Xi$  is defined

$$\Xi \triangleq \{ \xi \in \mathbf{R}^k \mid \xi_1 = 1, \ B\xi \in \mathbf{K} \},\$$

where  $\mathbf{K}$  is a proper cone, e.g.,



- positive orthant
- second-order cone
- positive semidefinite cone

# Robust SDPs : Challenges

The robust SDP is **NP-hard**, in general

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 subject to  $\sum_{i=1}^{k} \xi_i \mathcal{A}_i(x) \in \mathbf{S}_+^n, \quad \forall \ \xi \in \Xi,$ 

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The constraint holds if and only if

$$\min_{\xi \in \Xi} \ \lambda_{\min} \left( \sum_{i=1}^{k} \xi_i \mathcal{A}_i(x) \right) \ge 0.$$
concave in  $\xi$ 

 $\lambda_{\min}: \mathbf{S}^n \to \mathbf{R}$  is the minimum eigenvalue function

# A Robust LP Approach

Approximate the robust SDP with a robust linear program (LP)

by approximating the PSD cone by a polyhedral cone

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Robust LPs: Admit finite-dimensional reformulations as conic convex programs over the cone  $\mathbf{K}$  characterizing uncertainty set  $\Xi$ , e.g.

K	Polyhedral Cone	Second-order Cone	Semidefinite Cone
Robust LP	LP	SOCP	SDP

# Finite-Dim. Reformulations and Approximations

#### **Exact Reformulations**

- Ben-Tal, El-Ghaoui, Nemirovski, ['00]
  - Ξ is "Unstructured norm-bounded"

#### **Inner Approximations**

- Ben-Tal, El-Ghaoui, Nemirovski, ['00]
  - \(\pi\) is "Structured norm-bounded"
- Scherer and Hol, ['06]
  - $\Xi$  is described by polynomial matrix inequalities

#### Other related work

- Packard et al. ['93] Scherer ['05]

- Ben-Tal et al. ['02]
- El-Ghaoui et al. ['97] Dietz et al. ['08]
- Oishi et al. | '08|

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### Talk Outline

1. Introduction

2. Inner and Outer Polyhedral Hierarchies of the PSD Cone

3. Inner and Outer Hierarchies of Robust SDPs

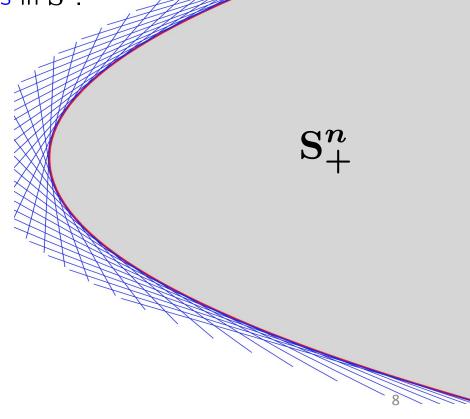
4. Application: Robust Resistance Network Design Problem

# Polyhedral Approximation of $S_{+}^{n}$

The positive semidefinite cone

$$\mathbf{S}_{+}^{n} = \bigcap_{u \neq 0} \left\{ X \in \mathbf{S}^{n} \mid u^{\top} X u \ge 0 \right\}$$

is an infinite intersection of half-spaces in  $S^n$ .



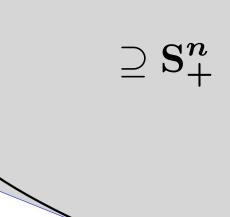
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• A finite intersection of half-spaces yields an outer polyhedral cone to  $\mathbf{S}_{+}^{n}$ .



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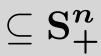
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• A finite intersection of half-spaces yields an outer polyhedral cone to  $\mathbf{S}^n_+$ 

• The dual of the outer polyhedral cone is an inner polyhedral cone to  $\mathbf{S}^n_+$ 

$$\mathbf{S}_{+}^{n} \subseteq \mathsf{polyhedral} \iff (\mathsf{polyhedral})^{*} \subseteq (\mathbf{S}_{+}^{n})^{*} = \mathbf{S}_{+}^{n}$$



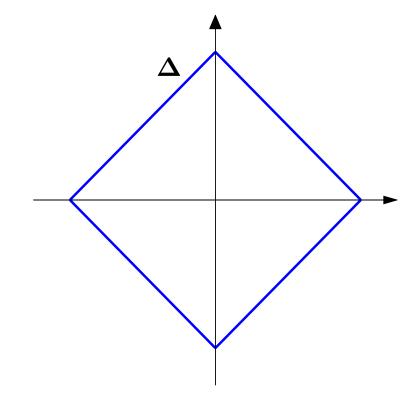
# Construction of Outer Polyhedral Approximations

Let  $\Delta$  denote the boundary of the  $\ell_1$  norm ball in  ${\bf R}^n$ 

$$\Delta := \{ x \in \mathbf{R}^n \mid ||x||_1 = 1 \}.$$

The PSD cone can be expressed as:

$$\mathbf{S}_{+}^{n} = \bigcap_{u \in \Delta} \{ X \in \mathbf{S}^{n} \mid u^{\top} X u \ge 0 \}$$



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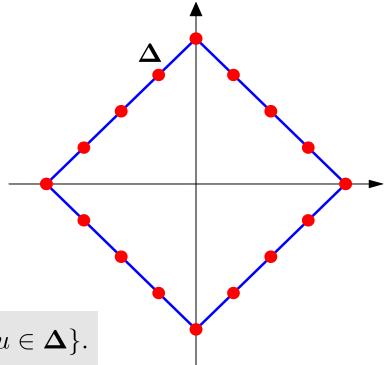
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• An outer polyhedral cone to  $\mathbf{S}^n_+$  arises by a discretization of  $\Delta$ , i.e.,

$$\mathbf{S}_{+}^{n} \subseteq \{X \in \mathbf{S}^{n} \mid u^{\top}Xu \ge 0, \text{ for some } u \in \mathbf{\Delta}\}.$$



Fix  $r \in \mathbb{N}$ . Consider the following discretization of  $\Delta \subseteq \mathbf{R}^n$ :

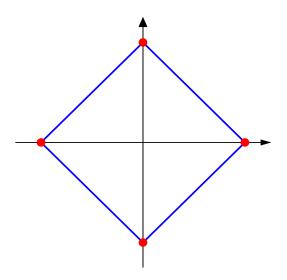
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Examples: n=2

$$\mathbf{\Delta_0} = \left\{ \begin{bmatrix} \pm 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ \pm 1 \end{bmatrix} \right\}$$



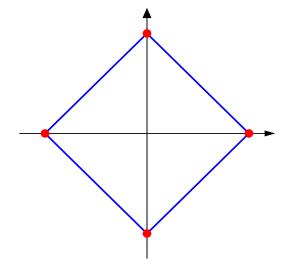
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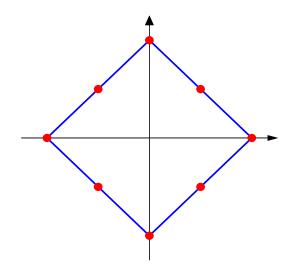
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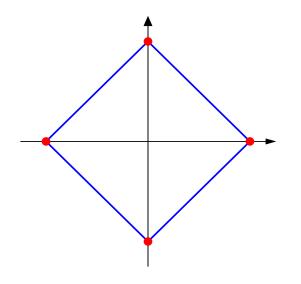
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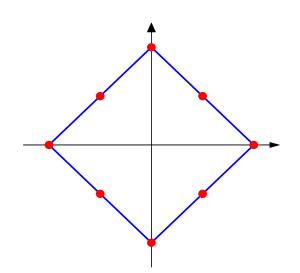
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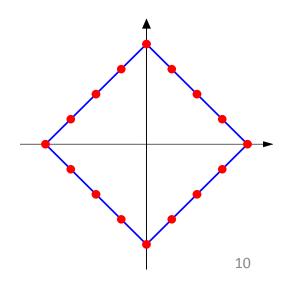
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**Remarks:** For any  $r \in \mathbb{N}$ , it holds that

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**Some notation:** The set  $\Delta_r$  has  $p_r$  elements denoted by:

$$u_1, u_2, \ldots, u_{p_r}$$

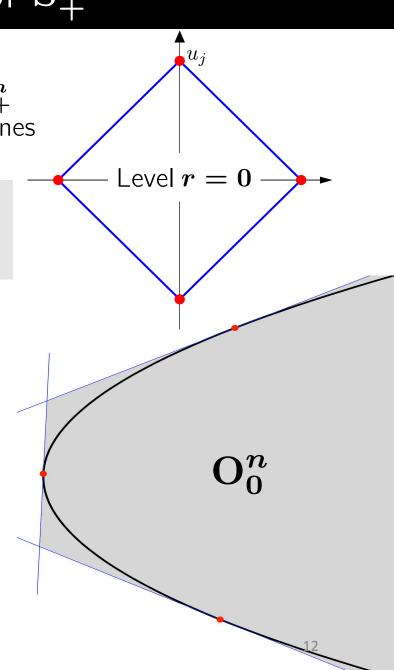
# Outer Polyhedral Hierarchies of $\mathbf{S}_{+}^{n}$

A hierarchy of outer polyhedral cones to  $\mathbf{S}_{+}^{n}$  arises by the following family of polyhedral cones

$$\mathbf{O}_{r}^{n} := \bigcap_{u \in \Delta_{r}} \{ X \in \mathbf{S}^{n} \mid u^{\top} X u \ge 0 \}$$

where  $r \in \mathbb{N}$ . In particular

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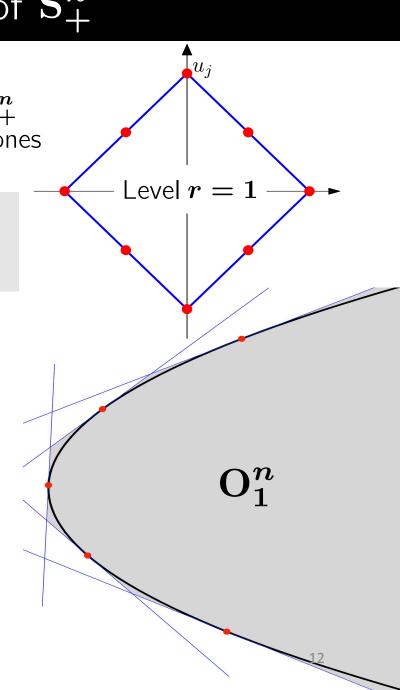
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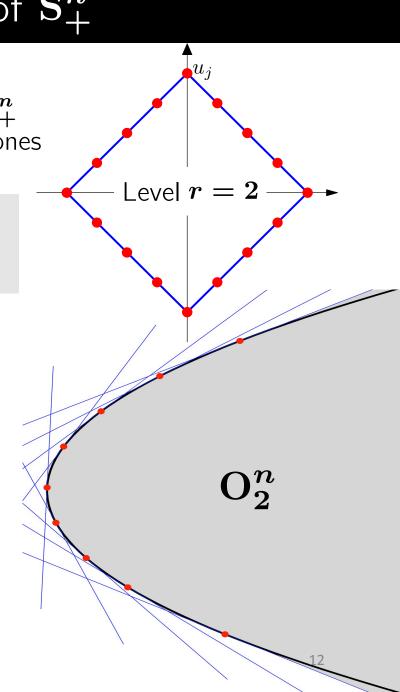
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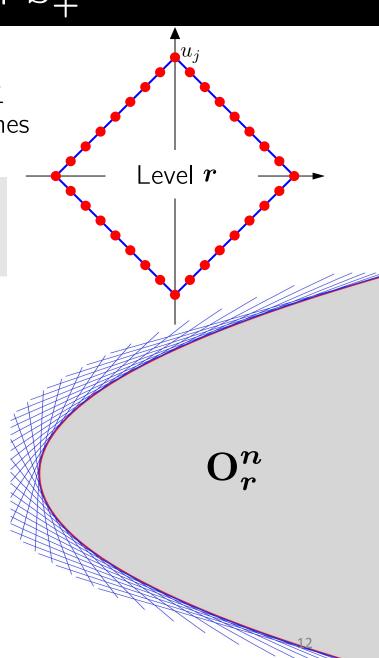
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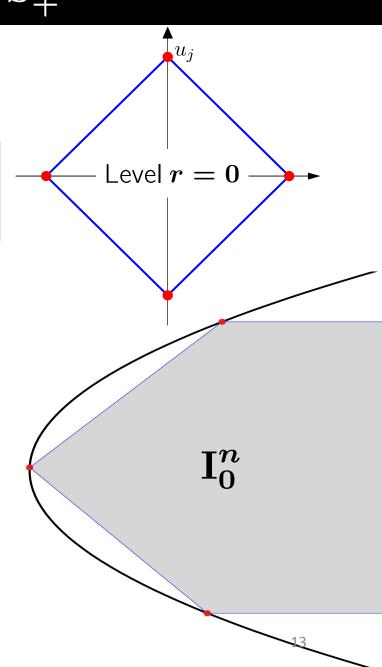


The dual cones to  $\mathbf{O}_r^n$  give a hierarch of  $\mathbf{S}_+^n$  inner polyhedral cones to  $\mathbf{S}_+^n$ 

$$\mathbf{I_r^n} = (\mathbf{O_r^n})^* = \operatorname{cone} \left\{ u_1 u_1^\top, \dots, u_{p_r} u_{p_r}^\top \right\}$$

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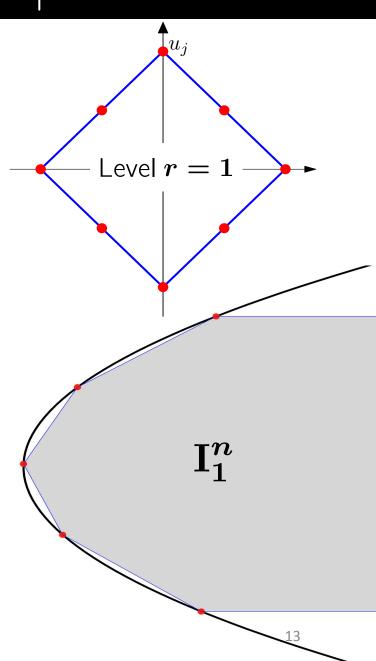


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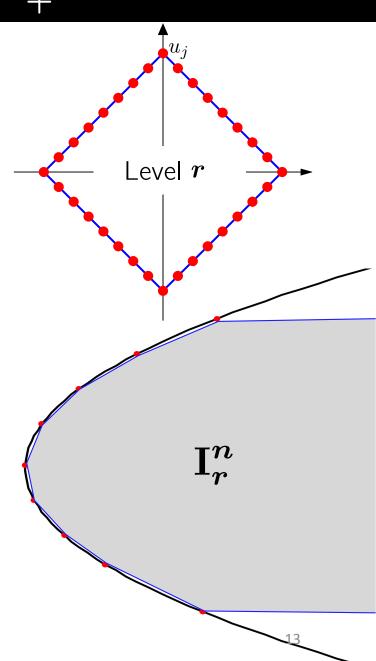


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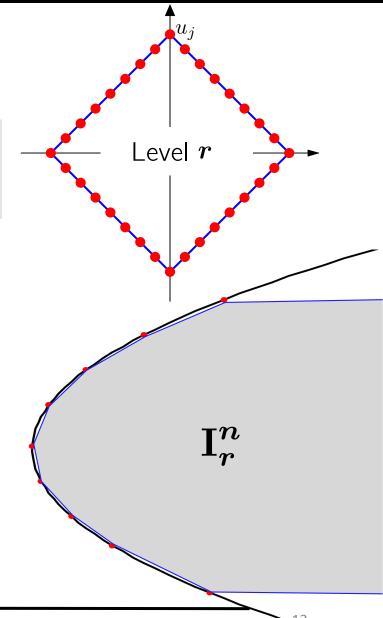
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#### **Examples**

- $\mathbf{I_0^n}$ : cone of nonnegative diagonal matrices
- $\mathbf{I_1^n}$ : cone of diagonally dominant matrices with nonnegative diagonal entries.



# Polyhedral Hierarchies of the PSD Cone

**Theorem:** For each level  $r \in \mathbb{N}$  ,

1. 
$$\mathbf{O}^n_r \supseteq \mathbf{O}^n_{r+1} \supseteq \mathbf{S}^n_+$$
 and

$$igcap_{i\in\mathbb{N}}\mathbf{O}_i^n\ =\ \mathbf{S}_+^n$$

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 and  $\operatorname{cl}\left(\left(\begin{array}{c} \mathbf{I}_r^n \end{array}\right)^{n}\right)$ 

$$\operatorname{cl}\left(\bigcup_{i\in\mathbb{N}}\mathbf{I}_{i}^{n}\right) = \mathbf{S}_{+}^{n}$$

[Braun, Fiorini, Pokutta, Steurer '12] "Approximation limits of linear programs (beyond hierarchies)."

"It's not possible to approximate SDPs arbitrarily well using small LPs"

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1. Introduction

2. Inner and Outer Polyhedral Hierarchies of the PSD Cone

3. Inner and Outer Hierarchies of Robust SDPs

4. Application: Robust Resistance Network Design Problem

# Outer Approximations to Robust SDP

Recall the outer and inner polyhedral cones approximating  $\mathbf{S}^n_+$ 

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For any  $r \in \mathbb{N}$ , the robust LP:

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# Finite-Dimensional Outer Approximation

The **hyperplane representation** of the outer polyhedral cones  $\mathbf{O}_r^n$ 

$$\mathbf{O}_{m{r}}^{m{n}} = \bigcap_{i=1}^{p_r} \mathsf{half-space}_i$$

and strong duality gives a finite-dimensional representation of the robust LP.

#### Finite-Dimensional Outer Approximation

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$$\mathbf{O}_{\boldsymbol{r}}^{\boldsymbol{n}} = \bigcap_{j=1}^{p_r} \{ X \in \mathbf{S}^n \mid u_j^\top X u_j \ge 0, \ u_j \in \boldsymbol{\Delta}_{\boldsymbol{r}} \}$$

and strong duality gives a finite-dimensional representation of the robust LP.

**Theorem :** The robust LP over  $\mathbf{O}_r^n$  admits an equivalent reformulation as a finite-dimensional conic linear program:

minimize 
$$c^{\top}x$$
  
subject to  $x \in \mathbf{R}^{m}, \mu_{j} \in \mathbf{R}_{+}, \lambda_{j} \in \mathbf{K}^{*}, \quad \forall j = 1, \dots, p_{r}$   
 $u_{j}^{\top} \mathcal{A}_{i}(x) u_{j} = \mu_{j} + e_{i}^{\top} B^{\top} \lambda_{j}, \quad \forall j = 1, \dots, k$   
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Its optimal value is a **lower bound** to the optimal value of the robust SDP.

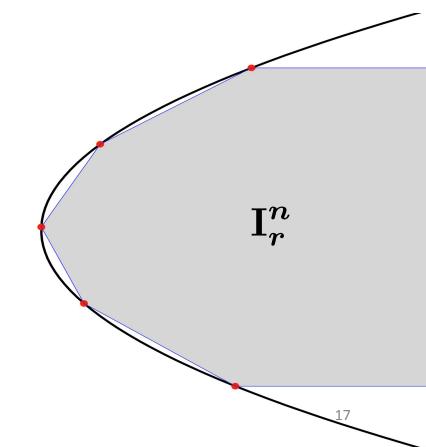
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### A Challenge with the Inner Approximation

The vertex representation of the inner polyhedral cone  $\mathbf{I}^{m{n}}_{m{r}}=(\mathbf{O}^{m{n}}_{m{r}})^*$ 

$$\mathbf{I}_{r}^{n} = \operatorname{cone}\left\{u_{1}u_{1}^{\top}, \dots, u_{p_{r}}u_{p_{r}}^{\top}\right\}$$

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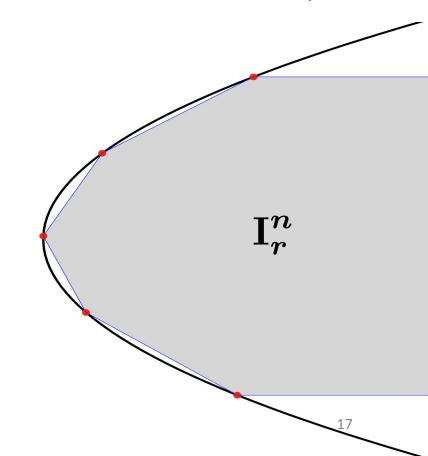
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#### Hyperplane representation of $\mathbf{I}_r^n$

• There exists  $q_r < \infty$  such that

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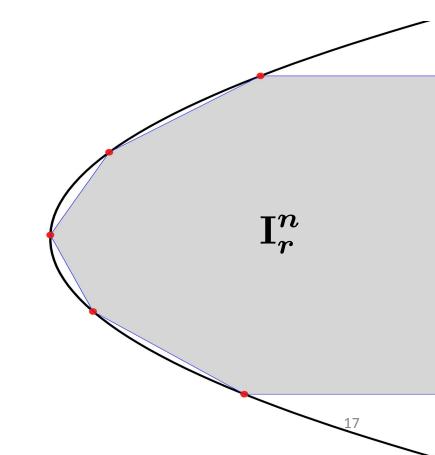
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• Yields a finite-dim. reformulation of robust LP over  $\mathbf{I}_r^n$ 



## Finite-Dimensional Outer Approximation

Let

$$\mathbf{I}_{r}^{n} = \bigcap_{j=1}^{q_{r}} \{X \in \mathbf{S}^{n} \mid \operatorname{tr}(H_{i}X) \ge 0\}$$

be the hyperplane representation of the inner polyhedral cone  ${f I}_r^n$ 

### Finite-Dimensional Outer Approximation

Let

$$\mathbf{I}_{r}^{n} = \bigcap_{j=1}^{q_{r}} \{X \in \mathbf{S}^{n} \mid \operatorname{tr}(H_{i}X) \ge 0\}$$

be the hyperplane representation of the inner polyhedral cone  $\mathbf{I}_r^n$ 

**Theorem:** The robust LP over  $\mathbf{I}_r^n$  admits an equivalent reformulation as a finite-dimensional conic linear program:

minimize 
$$c^{\top}x$$
  
subject to  $x \in \mathbf{R}^{m}, \mu_{j} \in \mathbf{R}_{+}, \lambda_{j} \in \mathbf{K}^{*}, \quad \forall j = 1, \dots, q_{r}$   

$$\operatorname{tr}(\mathcal{A}_{i}(x)H_{j}) = \mu_{j} + e_{i}^{\top}B^{\top}\lambda_{j}, \quad \forall j = 1, \dots, k$$

$$\forall j = 1, \dots, q_{r}$$

Its optimal value is an **upper bound** to the optimal value of the robust SDP.

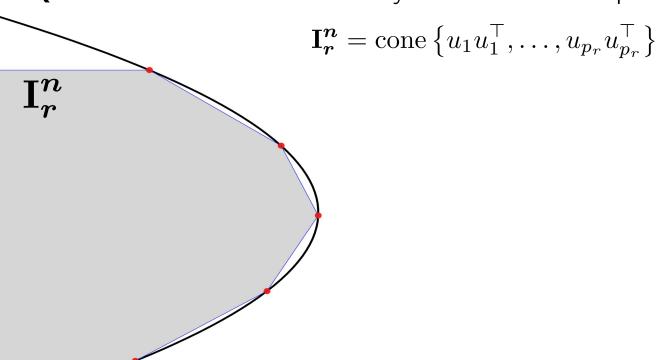
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**Question :** Can we work directly with the vertex representation of  $\mathbf{I}_r^n$ ?

$$\mathbf{I}_{r}^{n} = \operatorname{cone} \left\{ u_{1}u_{1}^{\top}, \dots, u_{p_{r}}u_{p_{r}}^{\top} \right\}$$

$$\exists \phi_{1}, \dots, \phi_{p_{r}} : \mathbf{R}^{k} \to \mathbf{R}_{+}, \text{ such that}$$

$$\sum_{i=1}^{k} \xi_{i} \mathcal{A}_{i}(x) \quad \forall \xi \in \Xi$$

$$\sum_{i=1}^{k} \xi_{i} \mathcal{A}_{i}(x) = \sum_{j=1}^{p_{r}} \phi_{j}(\xi)u_{j}u_{j}^{\top}, \quad \forall \xi \in \Xi.$$

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$$\Rightarrow \sum_{i=1}^{k} \xi_{i} \mathcal{A}_{i}(x) = \sum_{j=1}^{p_{r}} \phi_{j}^{\top} \xi, \quad \varphi_{j} \in \mathbf{R}^{k}$$

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Restriction to affine functions yields a finitedim. inner approximation to robust LP

Let

$$\mathbf{I}_{r}^{n} = \operatorname{cone}\left\{u_{1}u_{1}^{\top}, \dots, u_{p_{r}}u_{p_{r}}^{\top}\right\}$$

be the vertex representation of the inner polyhedral cone  $\mathbf{I}_r^n$ 

**Theorem:** The robust LP over  $\mathbf{I}_r^n$  admits an finite-dimensional inner approximation as a conic linear program:

minimize 
$$c^{\top}x$$
  
subject to  $x \in \mathbf{R}^{m}, \mu_{j} \in \mathbf{R}_{+}, \lambda_{j} \in \mathbf{K}^{*}, \quad \forall j = 1, \dots, p_{r}$   

$$\mathcal{A}_{i}(x) = \sum_{j=1}^{p_{r}} e_{i}^{\top}(\mu_{j}e_{1} + B^{\top}\lambda_{j})u_{j}u_{j}^{\top}, \quad \forall i = 1, \dots, k$$

Its optimal value is an **upper bound** to the optimal value of the robust SDP.

#### Talk Outline

1. Introduction

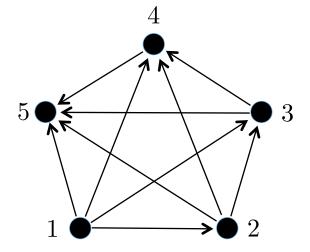
2. Inner and Outer Polyhedral Hierarchies of the PSD Cone

3. Inner and Outer Hierarchies of Robust SDPs

4. Application: Robust Resistance Network Design Problem

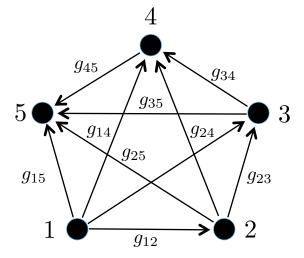
## Robust Resistance Network Design Problem

Given a circuit topology and a set  $\mathcal{I} = \{Q\xi \mid \xi \in \Xi\}$  of external currents



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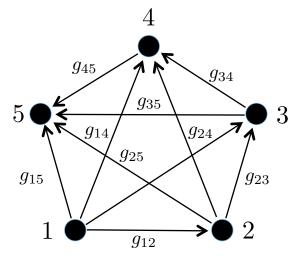


**Objective:** Choose a conductance  $g_{ij}$  for each line (i, j) such that:

$\displaystyle                                    $	maximal dissipation over $\mathcal{I}$				
subject to	$1^{\top}g \leq b$	budget constraint			
	$g \ge 0$	physical constraints			

## Robust Resistance Network Design Problem

Given a circuit topology and a set  $\mathcal{I} = \{Q\xi \mid \xi \in \Xi\}$  of external currents



**Objective:** Choose a conductance  $g_{ij}$  for each line (i, j) such that:

$$\begin{array}{ll} \underset{(\tau,g)}{\operatorname{minimize}} & \tau \\ \text{subject to} & \mathbf{1}^\top g \leq \omega \\ & g \geq 0 \\ & \left[ \begin{matrix} \tau & Q \xi \\ Q \xi & M \mathrm{diag}(g) M^\top \end{matrix} \right] \succeq 0, \quad \forall \; \xi \in \Xi \\ & \text{incidence matrix} \end{array}$$

$$\Xi = \{ \xi \in \mathbf{R}^6 \mid \|\xi\|_2 \le 2, \ \xi_1 = 1 \}$$

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#### **Robust LP Hierarchies**

		Level $r$ in Hierarchy				
		0	1	2	3	4
Lower Bound	$(\mathbf{O}^{\boldsymbol{n}}_{\boldsymbol{r}})$					
Upper Bound I	$(\; \mathbf{I^{n}_{r}}\;)$					
Upper Bound II	$(\ \mathbf{I}^n_r\ )$					

#### **Comparisons:**

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#### **Robust LP Hierarchies**

	Level $r$ in Hierarchy				
	0	1	2	3	4
Lower Bound $(\mathbf{O}^n_r)$	0				
Upper Bound I $(\mathbf{I}_{r}^{n})$	$\infty$				
Upper Bound II ( $\mathbf{I}_{r}^{n}$ )	$\infty$				

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#### **Robust LP Hierarchies**

	Level $r$ in Hierarchy				
	0	1	2	3	4
Lower Bound $(\mathbf{O}^n_r)$	0	0			
Upper Bound I $(\mathbf{I}_{r}^{n})$	$\infty$	4.75			
Upper Bound II ( $\mathbf{I}^{m{n}}_{m{r}}$ )	$\infty$	6.72			

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#### **Robust LP Hierarchies**

	Level $r$ in Hierarchy				
	0	1	2	3	4
Lower Bound $(\mathbf{O}^n_r)$	0	0	2.25		
Upper Bound I $(\mathbf{I}_{r}^{n})$	$\infty$	4.75	3.15		
Upper Bound II ( $\mathbf{I}^{m{n}}_{m{r}}$ )	$\infty$	6.72	4.94		

#### **Comparisons:**

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#### **Robust LP Hierarchies**

	Level $r$ in Hierarchy				
	0	1	2	3	4
Lower Bound $(\mathbf{O}^n_r)$	0	0	2.25	2.34	2.36
Upper Bound I $(\mathbf{I}_r^n)$	$\infty$	4.75	3.15	comp. e	expensive
Upper Bound II ( $\mathbf{I}_{r}^{n}$ )	$\infty$	6.72	4.94	4.56	4.55

#### **Comparisons:**

$$\Xi = \{ \xi \in \mathbf{R}^6 \mid \|(\xi_2, \xi_3)\|_2 \le 1, \|(\xi_4, \xi_5, \xi_6)\|_2 \le 1, \xi_1 = 1 \}$$

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#### **Robust LP Hierarchies**

	Level $r$ in Hierarchy				
	0	1	2	3	4
Lower Bound $(\mathbf{O}^n_r)$	0	1.65	3.66	4.19	4.24
Upper Bound I $(\mathbf{I}_{r}^{n})$	$\infty$	6.35	5.26	comp. e	xpensive
Upper Bound II ( $\mathbf{I}_{r}^{n}$ )	$\infty$	8.02	6.80	6.61	6.51

#### **Comparisons:**

• **Ben-Tal et. al ['00]** – (Upper Bound to Robust SDP)

 $\infty$ 

Scherer, Hol ['06] – (Upper Bound to Robust SDP)
 4.27

## Polytopic Uncertainty

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#### **Robust LP Hierarchies**

	Level $r$ in Hierarchy				
	0	1	2	3	4
Lower Bound $(\mathbf{O}^n_r)$	0	3.40	8.17	8.17	8.17
Upper Bound I $(\mathbf{I}_{r}^{n})$	$\infty$	8.96	8.44	comp. e	expensive
Upper Bound II ( $\mathbf{I}^{m{n}}_{m{r}}$ )	$\infty$	8.96	8.44	8.34	8.26

#### **Comparisons:**

- Nemirovski, El-Ghaoui ['00] (Not Applicable)
- Scherer, Hol ['06] (Upper Bound to Robust SDP)
   8.22

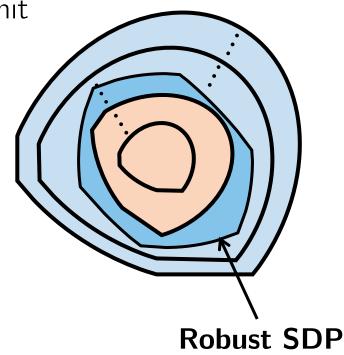
### Summary and Future Research

Developed computationally tractable inner and outer hierarchies

to robust SDPs that are exact in the limit

• **Approach**: Developed inner and outer polyhedral hierarchies to  $\mathbf{S}^n_+$ 

 Challenges: Impractical for moderate levels in the hierarchy!

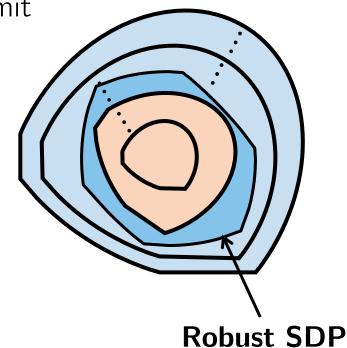


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Developed computationally tractable inner and outer hierarchies

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- **Approach**: Developed inner and outer polyhedral hierarchies to  $\mathbf{S}^n_+$
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#### **Future Research**

• Adaptively improve the polyhedral approx. of  $\mathbf{S}_{+}^{n}$  by using the guidance of the objective function!

# Questions?

Thank you! Raphael Louca e-mail: rl553@cornell.edu