

mlin: Rethinking and Rebooting gprof for the Multicore Era

nino Garcia, Donghwan Jeon, Chris Louie, Michael B. Taylor

Computer Science & Engineering Department
University of California, San Diego

Activating a “gprof for parallelization”

How effective are programmers at picking the right parts of a program to parallelize?

User study* we performed at UC San Diego (UCSD IRB #100056)

First and second year CS graduate students

Users parallelize their programs and submit to job queue for timing

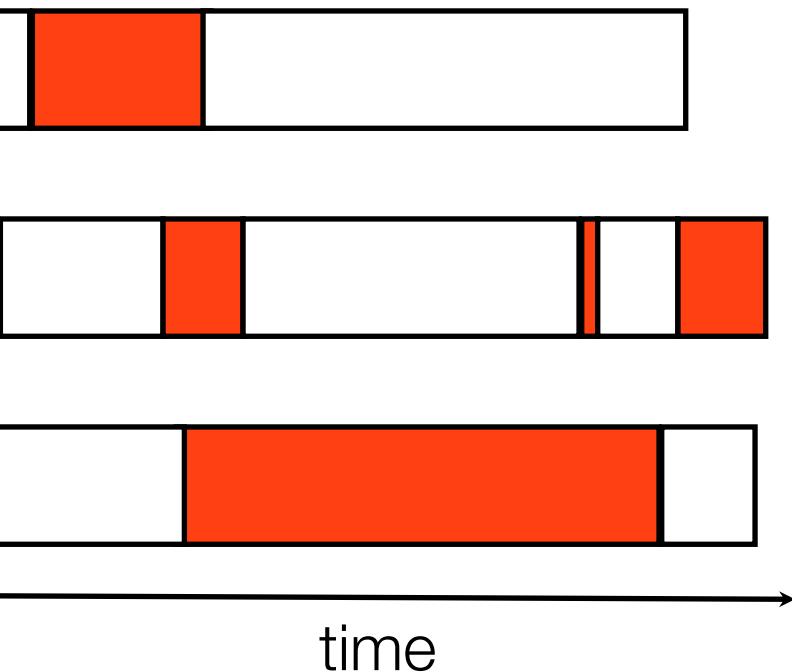
32-core AMD machine, Cilk++, access to gprof

Students were graded based on effectiveness of their parallel speedups
students told serial optimization would not help their grade

*Disclaimer: No graduate students were harmed in the making of this slide.

Case Study: Results

examined student's activities to determine result of efforts



■ Fruitless Parallelization Effort

Significant fraction of fruitless effort because of three basic problems

Low Parallelism: Region was not parallel enough

Low Coverage: Region's execution time was too small

gprof answers the question:
“*What parts of this program should I
spend time **optimizing**?*”

Kremlin answers the question:
“*What parts of this program should I
spend time **parallelizing**?*”

Kremlin's Usage Model

Usage model inspired by gprof

make CC=kremlin-cc

/tracking lolcats

kremlin tracking --personality=openmp

1. Compile instrumented binary

2. Profile with sample input

3. Run analysis tool to create report

file (lines)

imageBlur.c (49-58)

imageBlur.c (37-45)

setInterpPatch.c (26-35)

alcSobel_dx.c (59-68)

alcSobel_dx.c (46-55)

Self-P

145.3

145.3

25.3

126.2

126.2

Cov (%)

9.7

8.7

8.9

8.1

8.1

mlin's Key Components

al
c
de

- *Hierarchical Critical Path Analysis (HCPA)*

Parallelism Discovery

“What’s the potential parallel speedup of each part of this program?”

- *Self-Parallelism*

▶ Estimates ideal parallel speedup of a specific region

- *Planning Personalities*

Parallelism Planning

“What regions must I parallelize to get the maximum benefit on this system?”

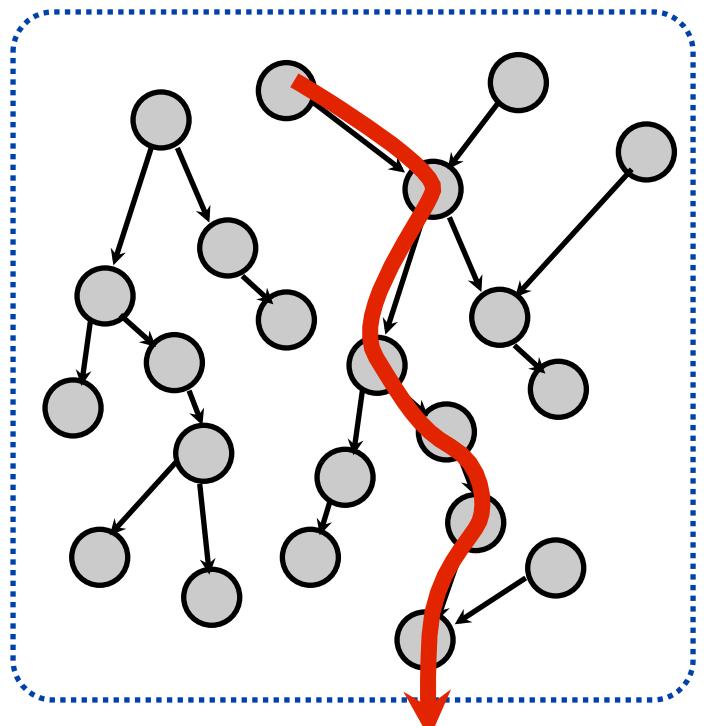
Parallelization

Developing an Approach for Parallelism Discovery

Testing Technique: 1980's-era Critical Path Analysis (CPA)

Finds critical path through the dynamic execution of a program

Mainly used in research studies to quantify limits of parallelism



$$\text{parallelism} = \frac{\text{work}}{\text{critical path length}}$$

work \approx # of instrs

instruction
→ data or control dependence

critical path (cp)

Benefits of CPA as a Basis for a Parallelism Discovery

Estimates program's potential for parallelization under relatively optimistic assumptions

Closer approximation to what human experts can achieve
versus pessimistic static analysis in automatic parallelizing compilers

It is predictive of parallelism after typical parallelization formations

e.g., Loop interchange, loop fission, locality enhancement

roving CPA with Hierarchical CPA (HCPA)

A is typically run on an entire program

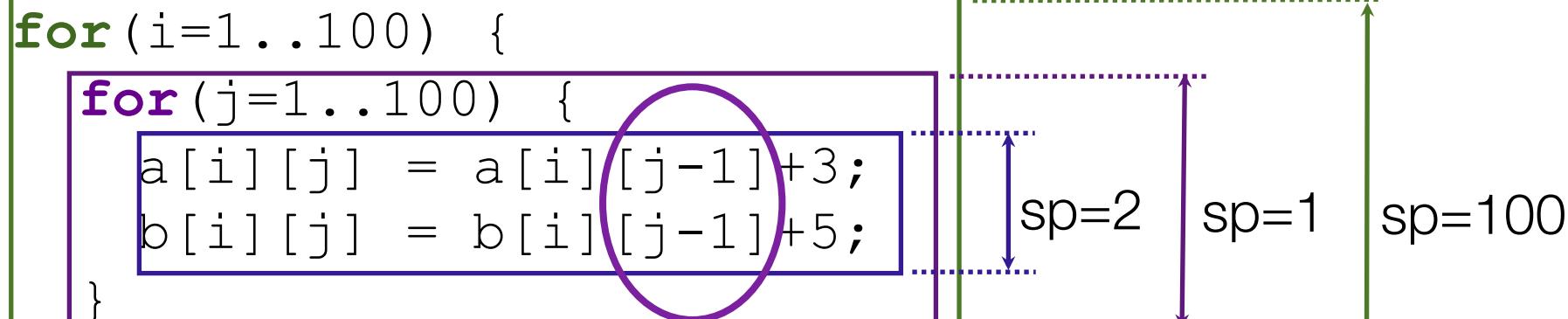
Not helpful for identifying specific regions to parallelize

Doesn't help evaluate execution time of a program if only a subset
the program is parallelized

archical CPA is a region-based analysis

Self-Parallelism (sp) identifies parallelism in specific regions

Provides basis for estimating parallel speedup of individual regions



PA Step 1: Hierarchically Apply CPA

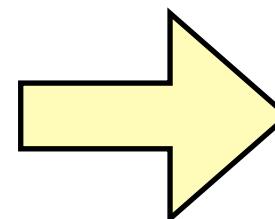
Goal: Introduce localization through *region-based* analysis

Shadow-memory based implementation

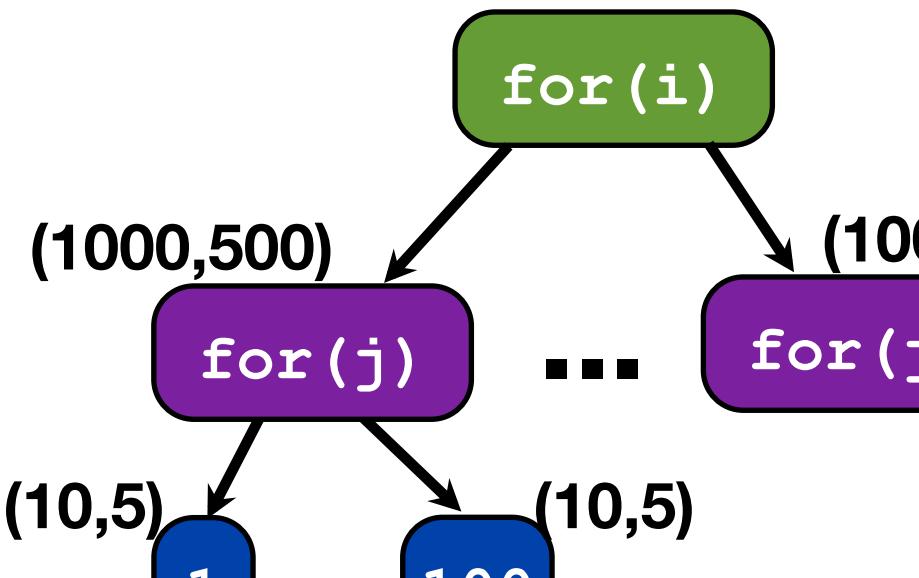
Performs CPA analysis on every program region

Single pass: Concurrently analyzes multiple nested regions

```
1..100) {  
    j=1..100) {  
        i[j] = a[i][j-1]+3;  
        i[j] = b[i][j-1]+5;
```



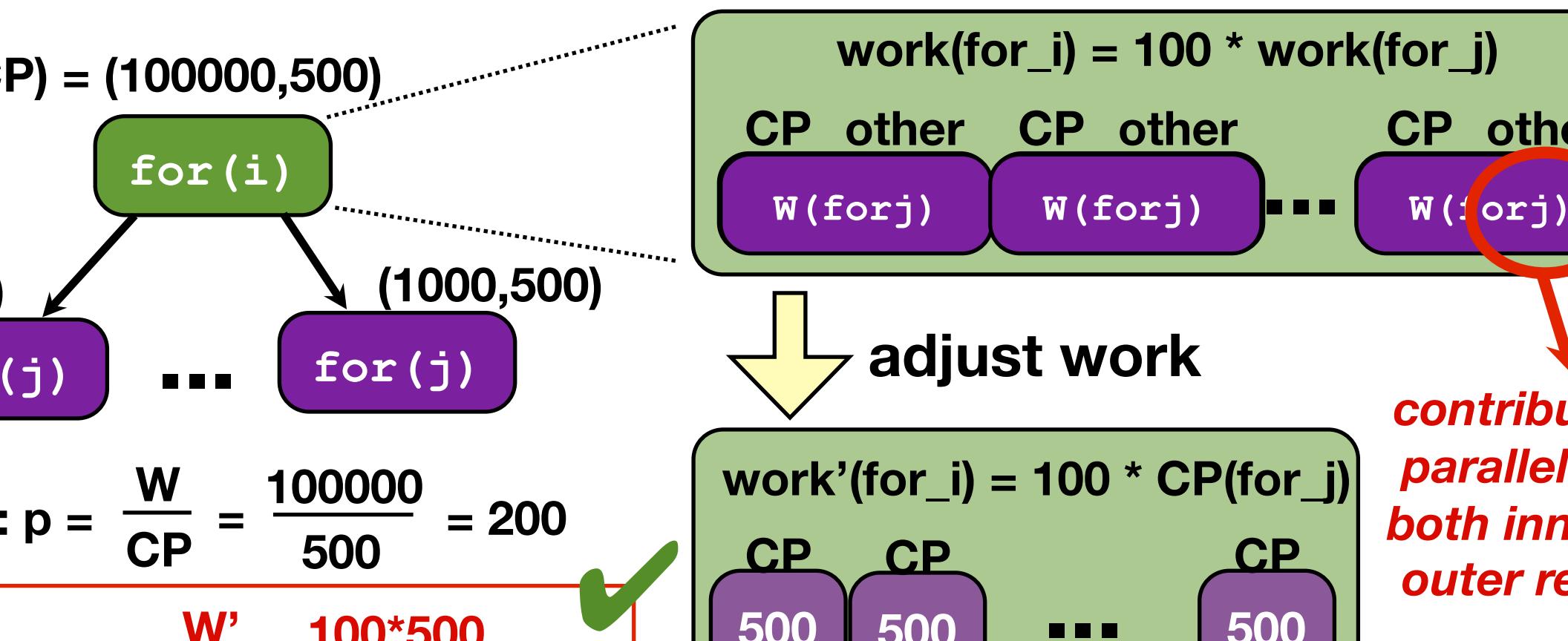
(work, cp length) = (100000,500)



$$p = \frac{W}{CP} = \frac{100000}{500} = 200 \times \text{X}$$

PA Step 2: Calculate Self-Parallelism

Goal: Eliminate effect of nested parallelism in parallelism calculation
approximate self-parallelism using only HCPA output
Subtracts" nested parallelism from overall parallelism



PA Step 3: Compute Static Region Data

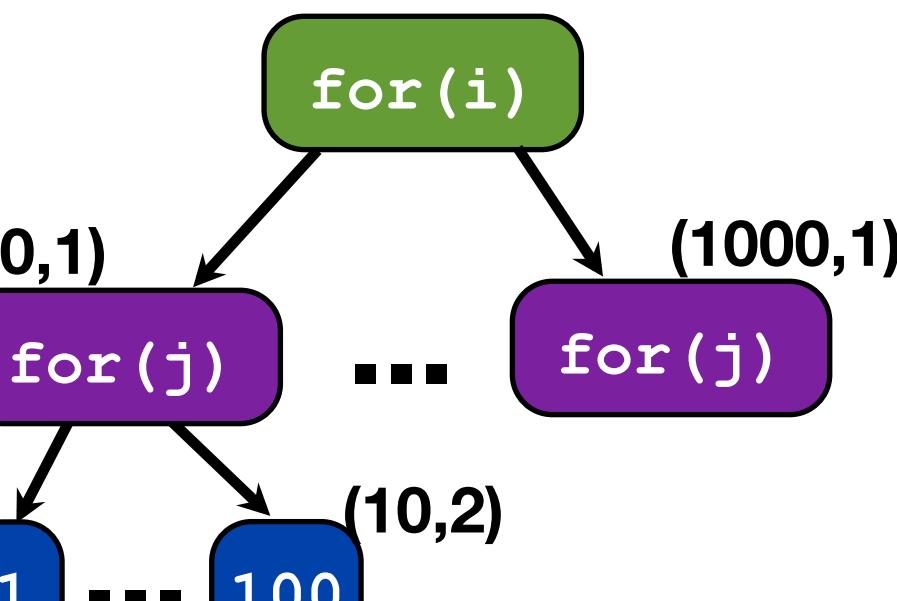
Goal: Convert dynamic region data to static region output

Merge dynamic nodes associated with same static region

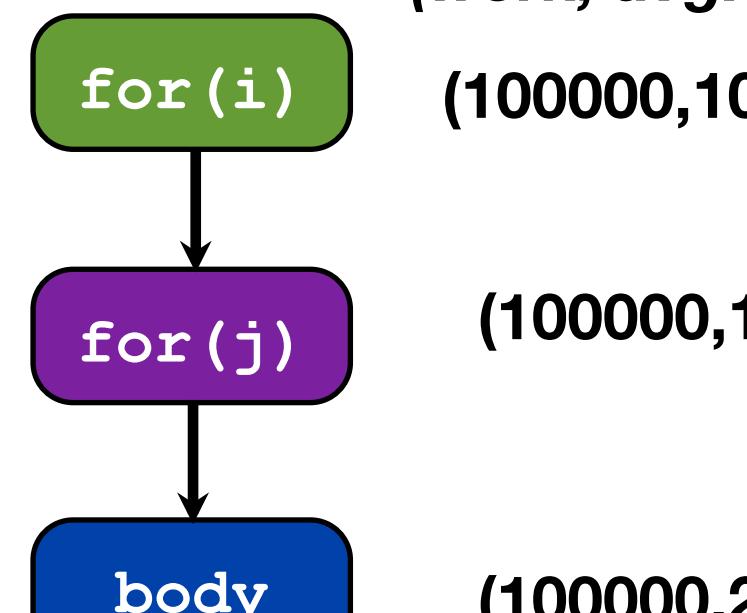
Work: Sum of work across dynamic instances

Self-Parallelism: Weighted Average across dynamic instances

work, sp) = (100000,100)



merge
dynamic



Other Details on Discovery in Our Paper

mlin handles much more complex structures than just nested loops: finds parallelism in arbitrary code including recursion

f-parallelism metric is defined and discussed in detail in the paper

compression technique used to reduce size of HCPA output

ating a Parallelization Plan

al: Use HCPA output to select best regions for target system

unning personalities allow user to incorporate system constraints

Software constraints: What types of parallelism can I specify?

Hardware constraints: Synchronization overhead, etc.

Planning algorithm can change based on constraints

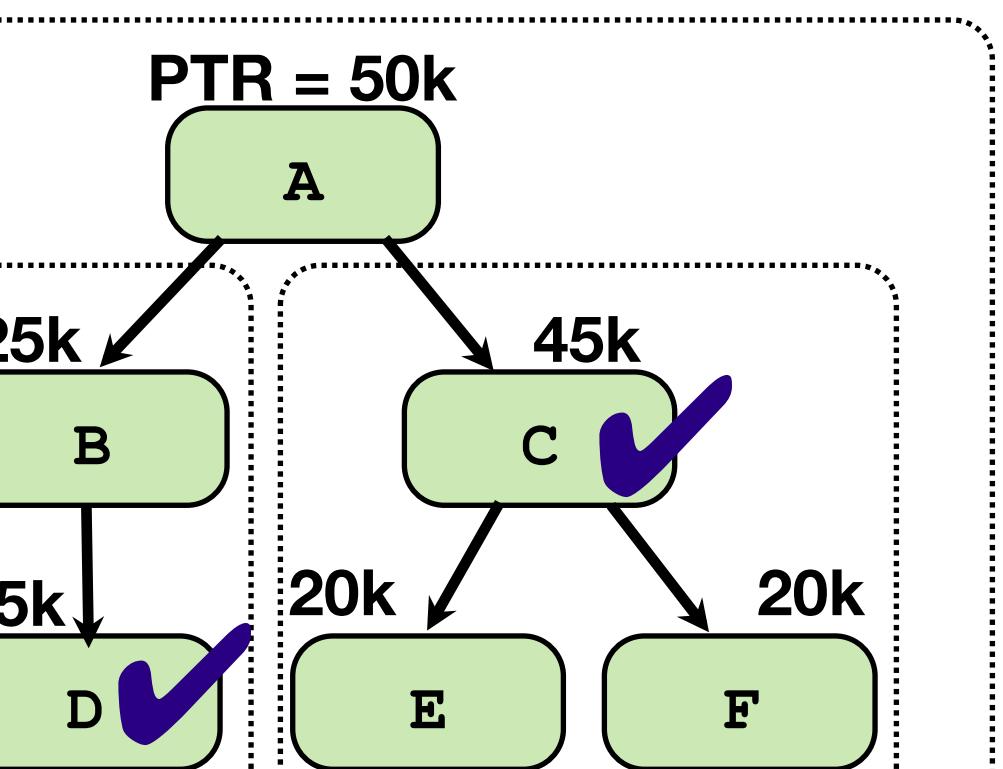
OpenMP Planner

Based on OpenMP 2.0 specification

Focused on loop-level parallelism

Disallow nested parallelism because of overhead

Planning algorithm based on dynamic programming



parallelized time reduction = W - (W / SP)

Region	Work	SP
A	100k	2
B	50k	2
C	50k	10
D	50k	10
E	25k	5

Evaluation

Methodology:

Ran Kremlin on serial versions; targeting OpenMP

Parallelized according to Kremlin's plan

Gathered performance results on 8 socket AMD 8380 Quad-core

Compared against third-party parallelized versions (3rd Party)

Benchmarks: NAS OpenMP and SpecOMP

Have both serial and parallel versions

Wide range of parallel speedup (min: 1.85x, max: 25.89x) on 32 cores

How much effort is saved using Kremlin?

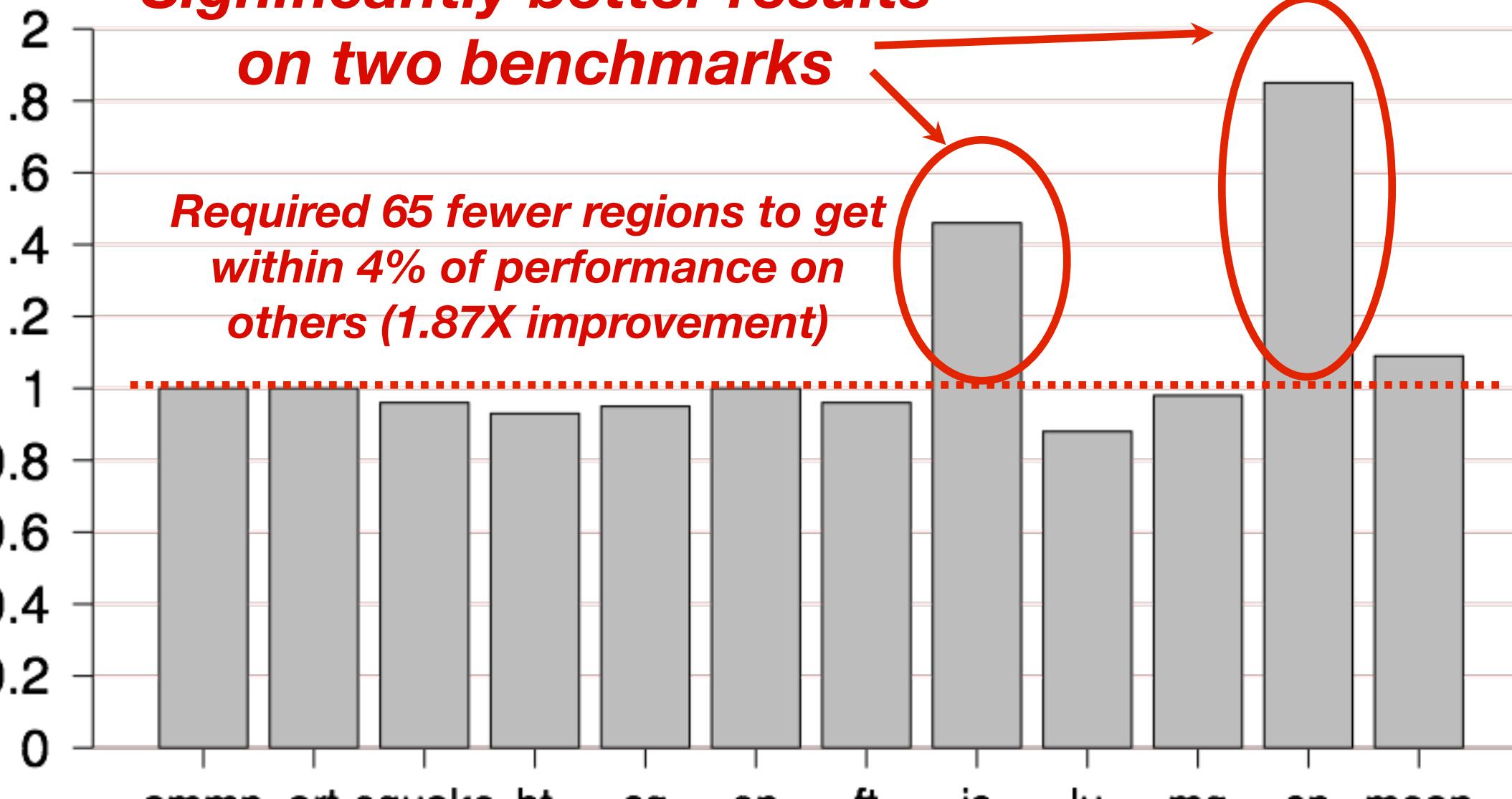
		# of Regions Parallelized		
Suite	Benchmark	3rd Party	Kremlin	Reduction
SciOMP	art	3	4	0.75x
	ampp	6	3	2.00x
	equake	10	6	1.67x
NPB	ep	1	1	1.00x
	is	1	1	1.00x
	ft	6	6	1.00x
	mg	10	8	1.25x
	cg	22	9	2.44x
	lu	28	11	2.55x
	bt	54	27	2.00x
	sp	70	58	1.21x
	Overall	211	134	1.57x

How good is Kremlin-guided performance?

Compared performance against expert, third-party version

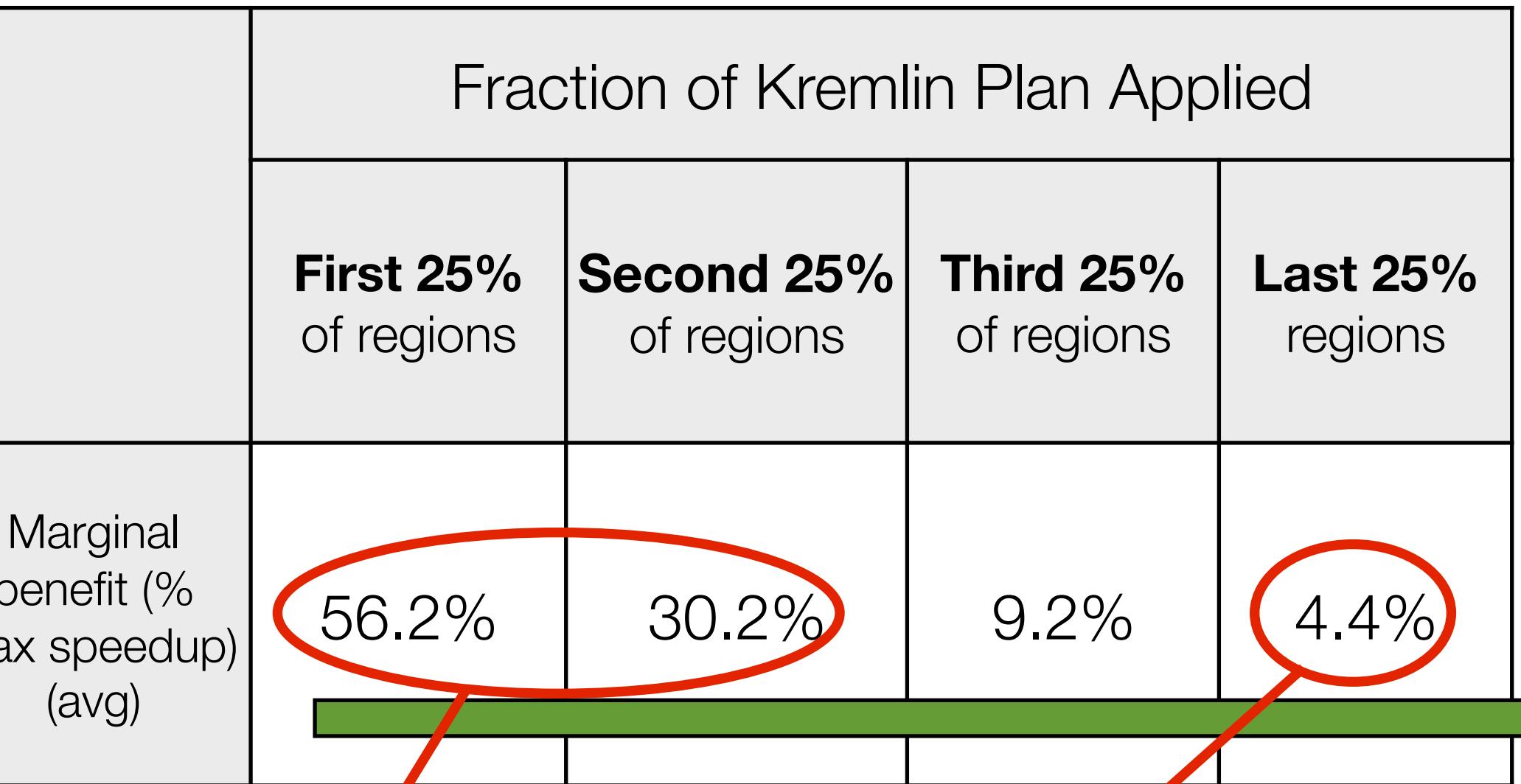
**Significantly better results
on two benchmarks**

**Required 65 fewer regions to get
within 4% of performance on
others (1.87X improvement)**



Does Kremlin pick the best regions first?

Determined what % of speedup comes from first {25,50,75,100} recommended regions



Inclusion

Kremlin helps a programmer determine:

'What parts of this program should I spend time parallelizing?'

Three key techniques introduced by Kremlin

Hierarchical CPA: How much total parallelism is in each region?

Self-Parallelism: How much parallelism is only in this region?

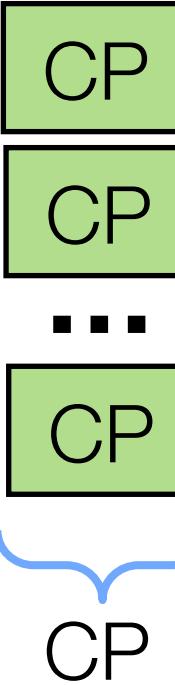
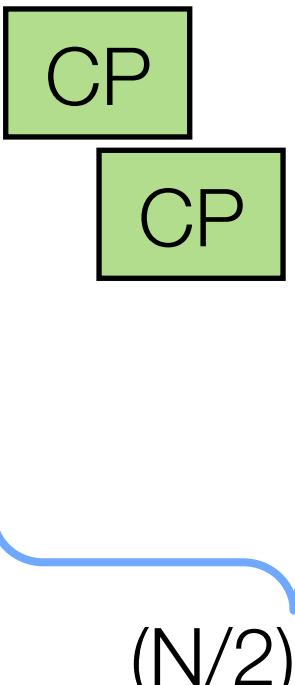
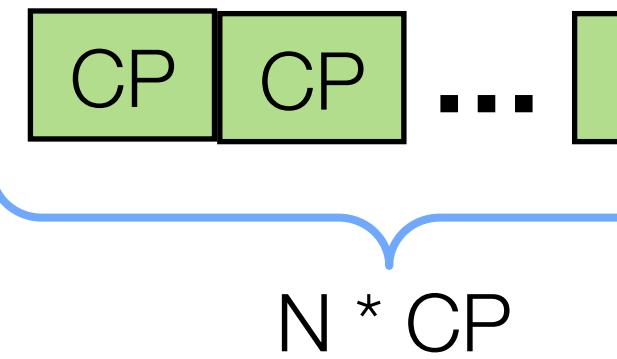
Planning Personalities: What regions are best for my target system?

Impressive results

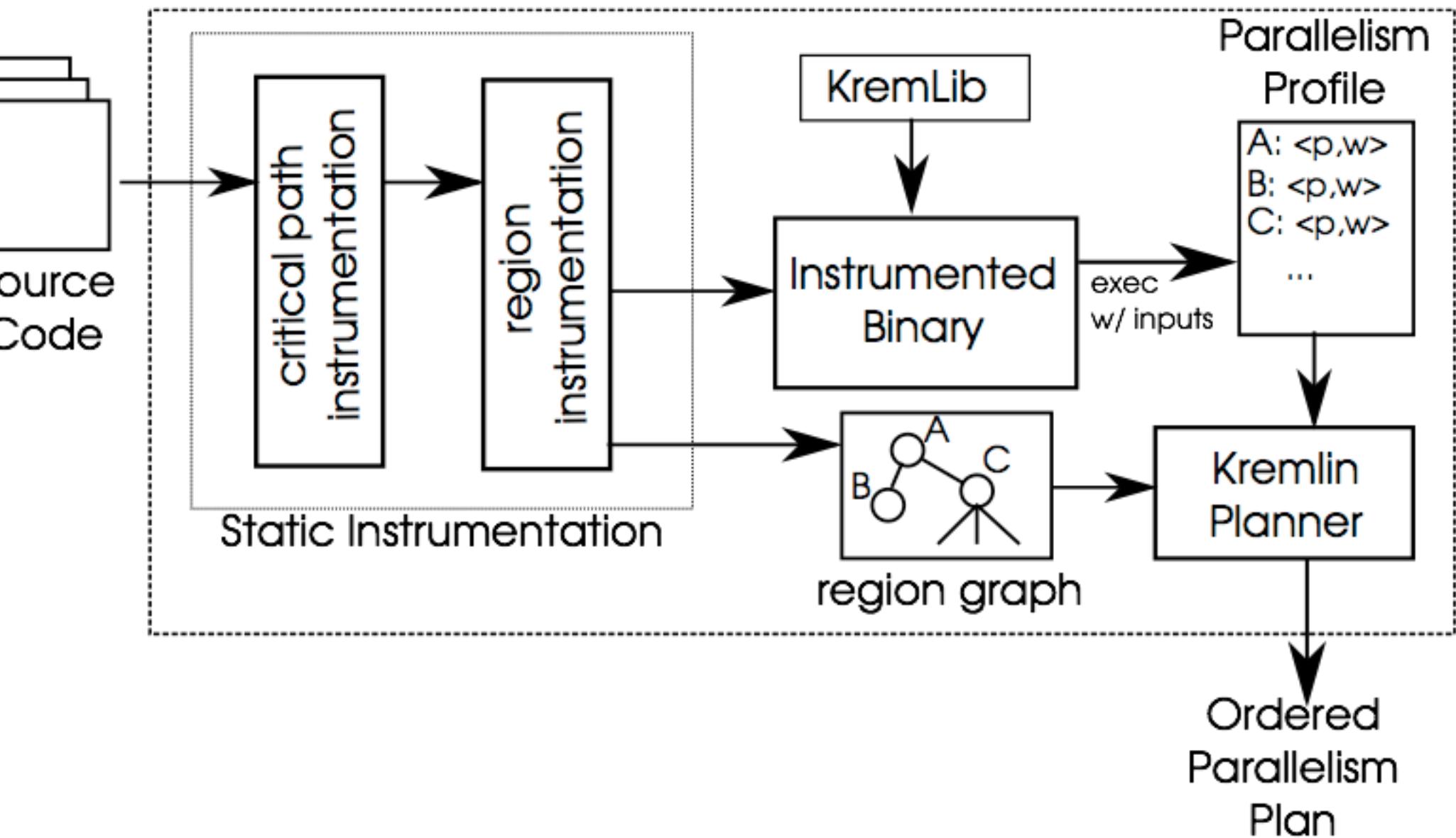
.57x average reduction in number of regions parallelized

Greatly improved performance on 2 of 11 benchmarks; very close on others

Parallelism for Three Common Loop Types

Type	DOALL	DOACROSS	Serial
Path (CP)			
Time (s)	$N * CP$	$N * CP$	$N * CP$
Efficiency	$\frac{N * CP}{CP} = N$	$\frac{N * CP}{(N/2) * CP} = 2.0$	$\frac{N * CP}{N * CP} = 1.$

Kremlin System Architecture



Interpreting the Parallelism Metric

