

# Tightening LP Relaxations for MAP using Message-Passing

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Joint work with Talya Meltzer, Amir Globerson,  
Tommi Jaakkola, and Yair Weiss

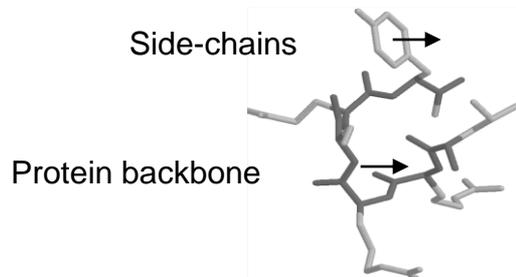


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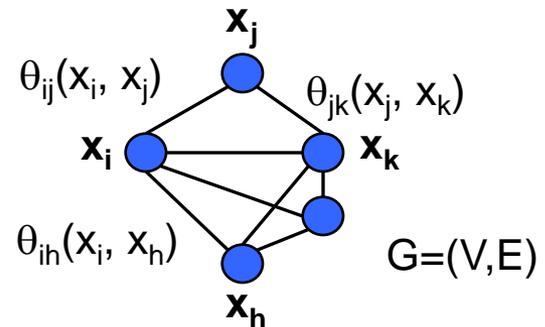


# MAP in Undirected Graphical Models

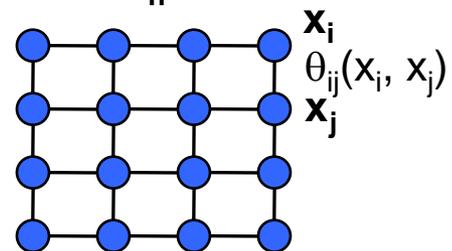
Real-world problems:



Protein design →



Stereo vision →



$$\Pr(x; \theta) \propto \exp \left( \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \right)$$

Find most likely

assignment:

$$x_{\text{map}} = \arg \max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j)$$

# How to solve MAP?

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- MAP is NP-hard
- Are real-world MAP problems really so hard?
- We give an algorithm which:
  - Improves approximation, using more computation
  - Problem-specific
  - If we *do* find best assignment, we know it
- Solves real-world problems *exactly*

# MAP as a linear program

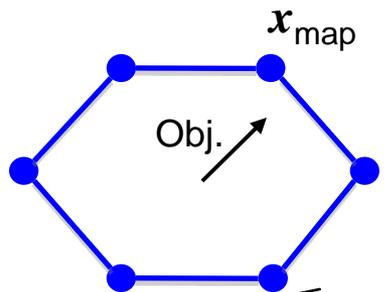
We can formulate the MAP problem as a linear

$$\max_{\mathbf{x}} \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) = \max_{\mathbf{x}} \sum_{\mu \in \mathcal{M}(G)} \sum_{(i,j) \in E} \delta(x_i, x_j) \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

where the variables  $\mu_{ij}$  are defined over edges.

The *marginal polytope* constrains the  $\mu_{ij}$  to be marginals of some distribution:

$$\mathcal{M}(G) = \{ \mu \mid \exists \Pr(\mathbf{x}; \theta) \text{ s.t. } \mu_{ij}(x_i, x_j) = \Pr(x_i, x_j; \theta) \}$$



Very many constraints!

Vertices correspond to assignments

# Relaxing the MAP LP

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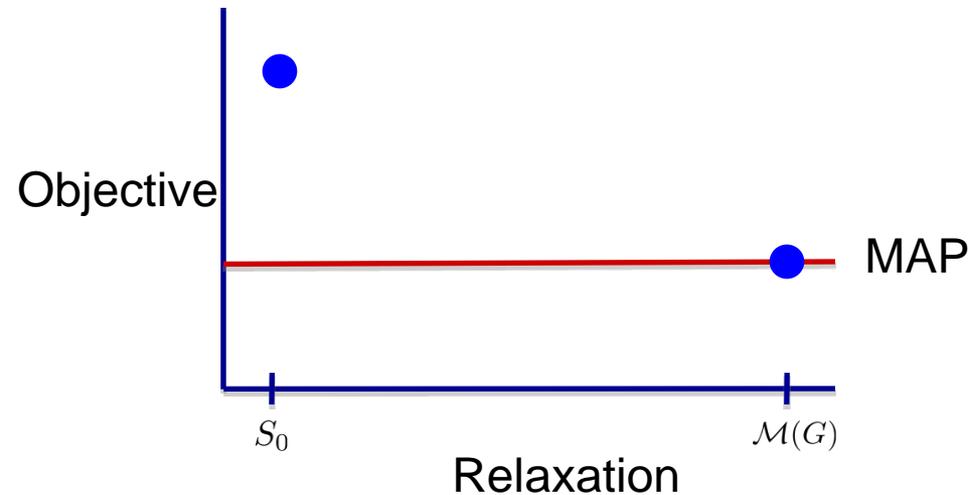
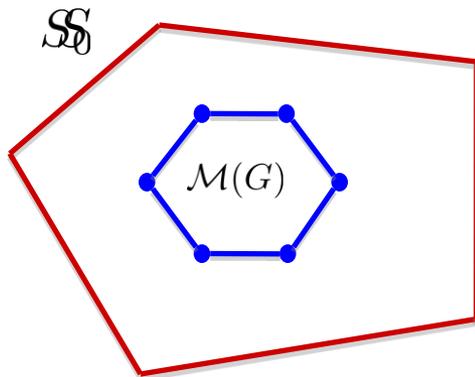
$$\max_{\mathbf{x}} \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) = \max_{\boldsymbol{\mu} \in \mathcal{M}(G)} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$



# Relaxing the MAP LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$

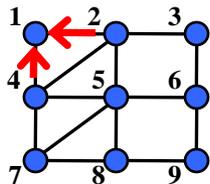
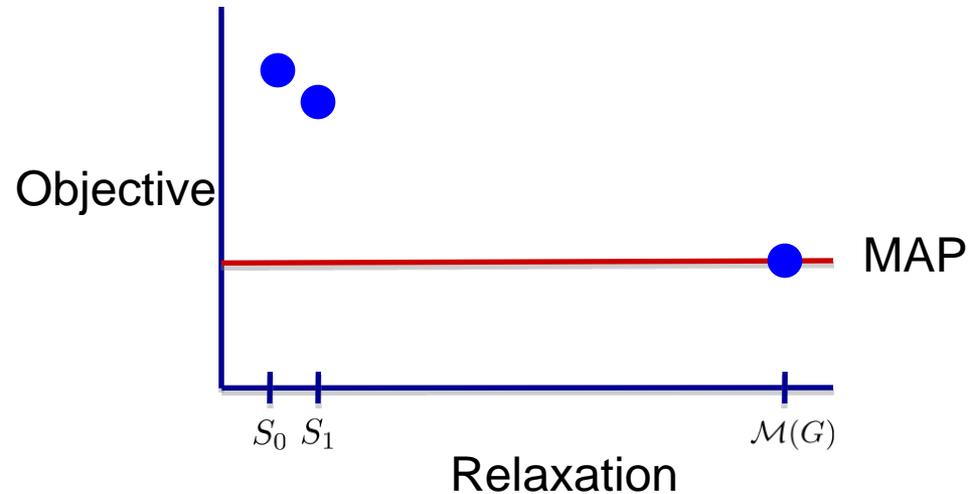
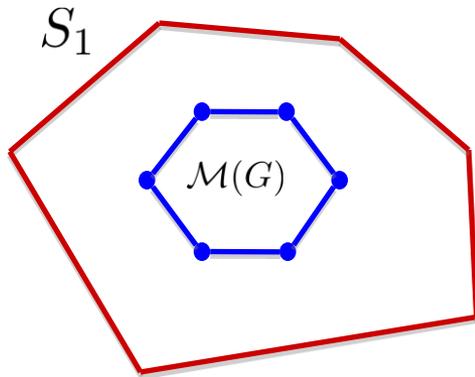


Simplest outer bound:  $\sum_{x_i, x_j} \mu_{ij}(x_i, x_j) = 1$

# Tightening the LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$



$$\sum_{x_2} \mu_{12}(x_1, x_2) = \sum_{x_4} \mu_{14}(x_1, x_4)$$

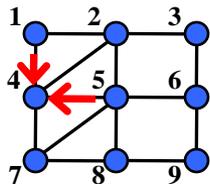
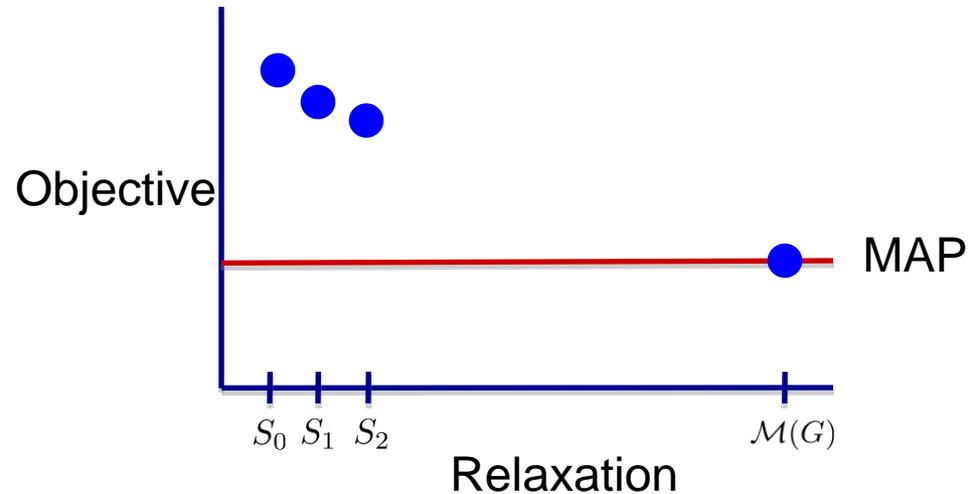
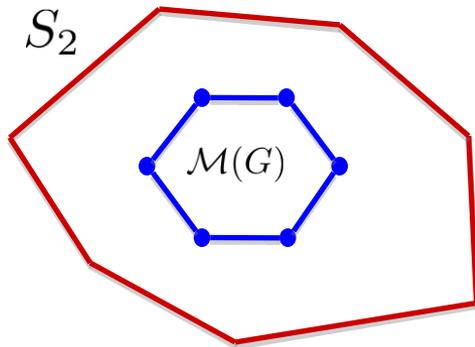


*Partial pairwise consistency*

# Tightening the LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$



$$\sum_{x_1} \mu_{14}(x_1, x_4) = \sum_{x_5} \mu_{45}(x_4, x_5)$$

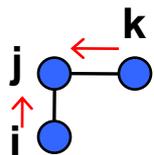
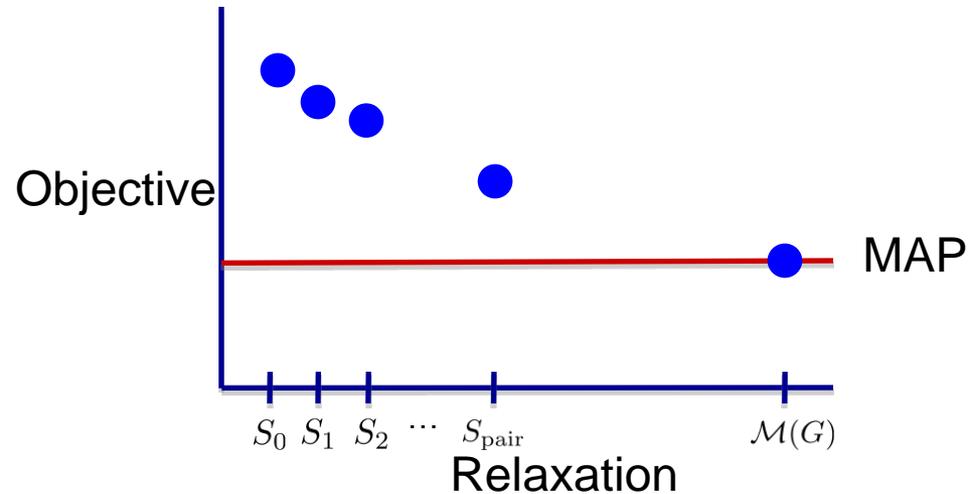
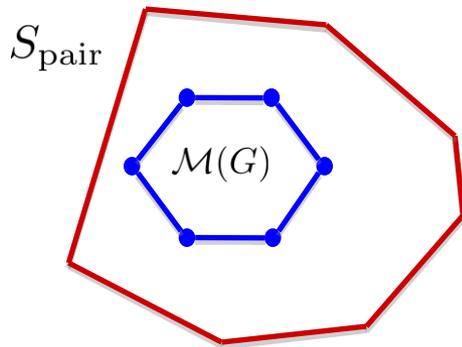


*Partial pairwise consistency*

# Tightening the LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$



$$\sum_{x_i} \mu_{ij}(x_i, x_j) = \sum_{x_k} \mu_{ij}(x_j, x_k)$$

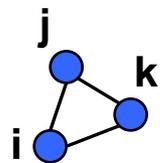
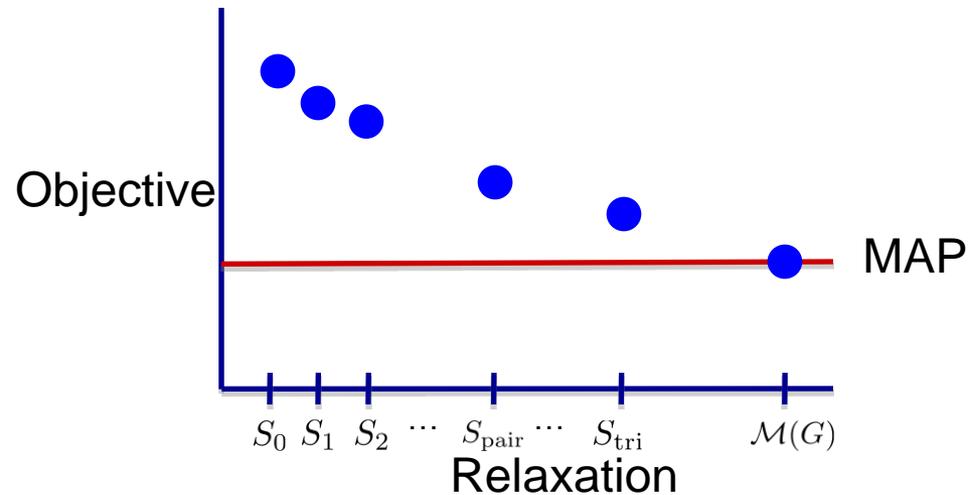
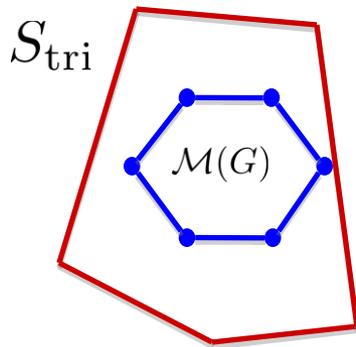


Pairwise consistency

# Tightening the LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$



$$\sum_{x_k} \mu_{ijk}(x_i, x_j, x_k) = \mu_{ij}(x_i, x_j)$$

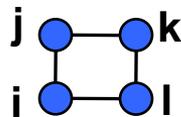
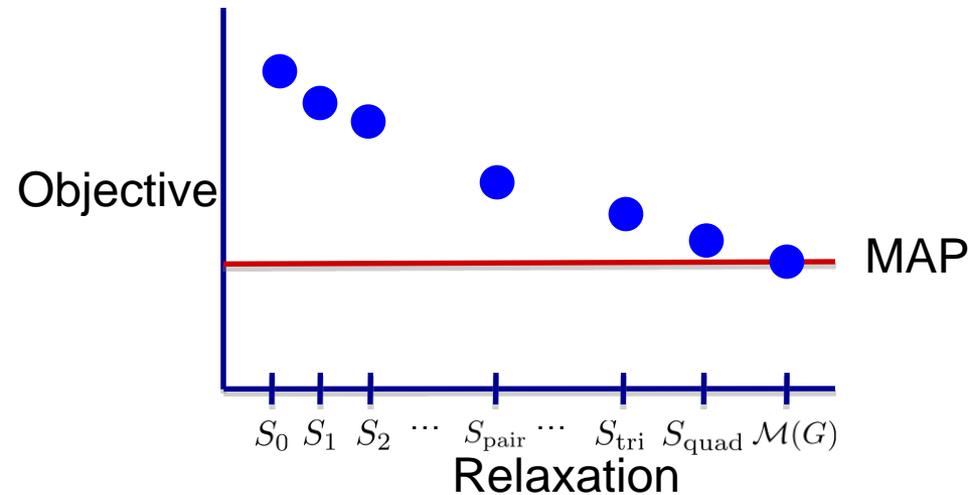
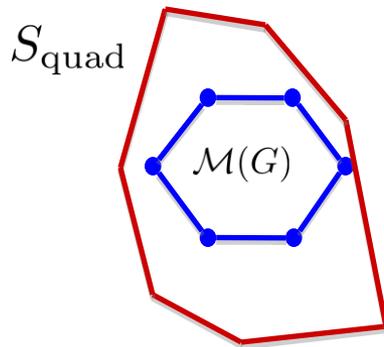


Triplet consistency

# Tightening the LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$



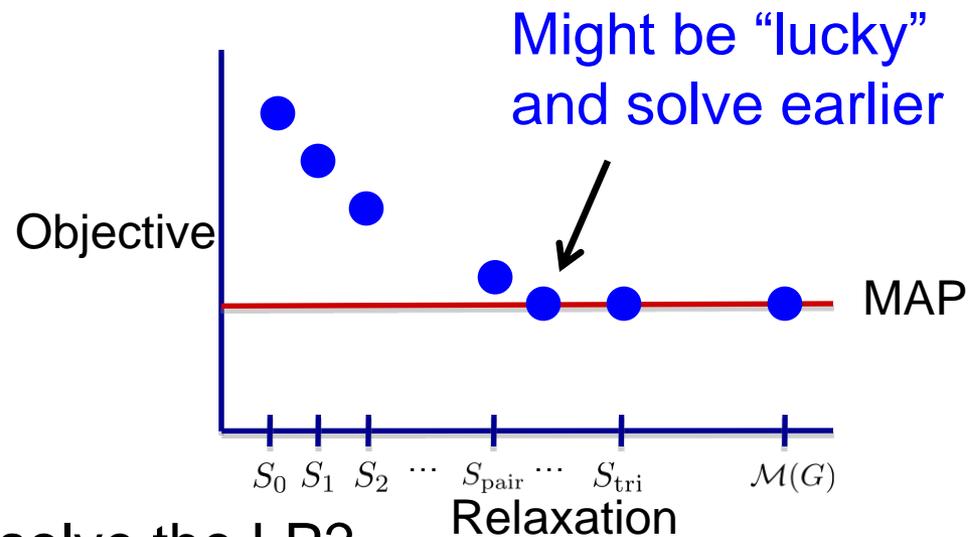
$$\sum_{x_k, x_l} \mu_{ijkl}(x_i, x_j, x_k, x_l) = \mu_{ij}(x_i, x_j)$$

“  
Quadruplet  
consistency

# Tightening the LP

$$\max_x \sum_{(i,j) \in E} \theta_{ij}(x_i, x_j) \leq \max_{\mu \in S} \sum_{(i,j) \in E} \sum_{x_i, x_j} \mu_{ij}(x_i, x_j) \theta_{ij}(x_i, x_j)$$

Such that  $\mathcal{M}(G) \subseteq S$

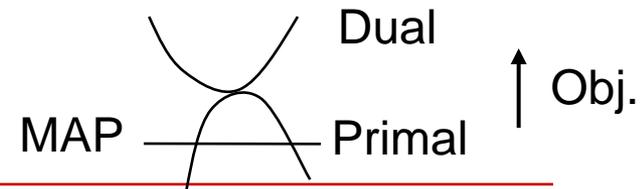


Great! But...

- ❑ Can we efficiently solve the LP?
- ❑ What clusters to add?
- ❑ How do we avoid re-solving?

# Our solution

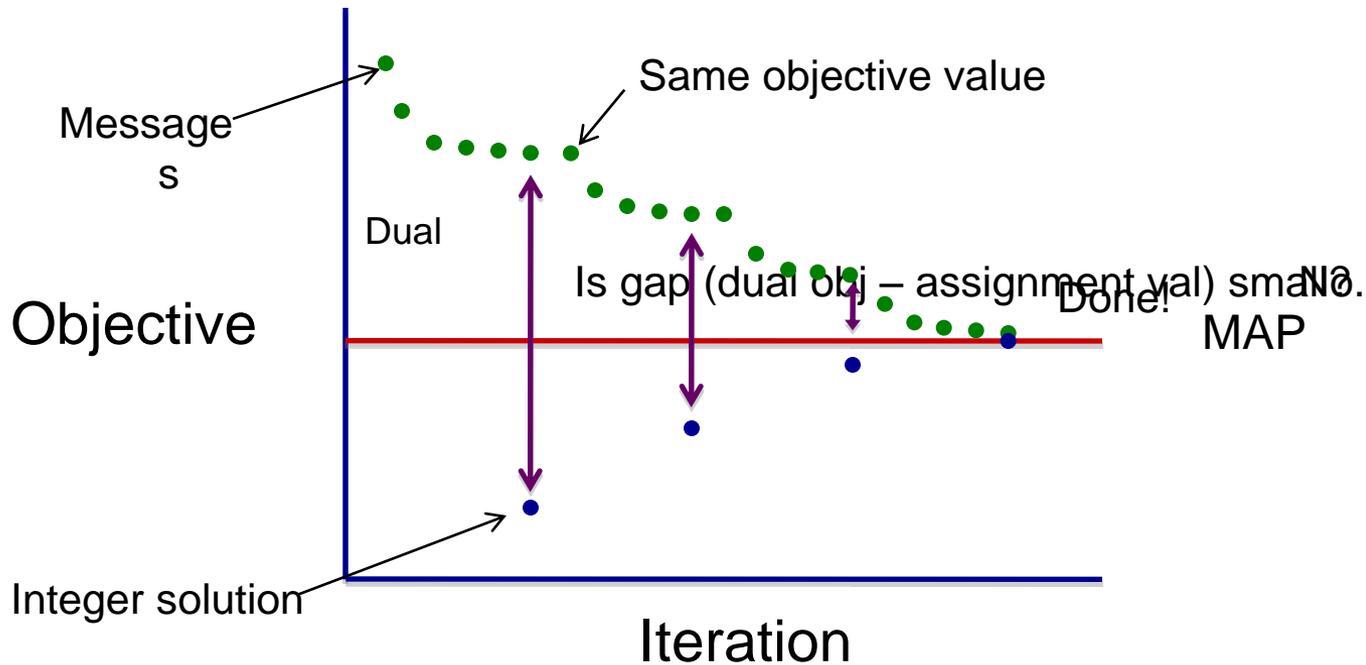
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- ❑ Can we efficiently solve the LP?
  - ❑ We work in the *dual* LP (Globerson & Jaakkola '07)
  - ❑ Dual can be solved by efficient message-passing algorithm
  - ❑ Corresponds to coordinate-descent algorithm
  
- ❑ What cluster to add next?
  - ❑ We propose a greedy *bound minimization* algorithm
  - ❑ Add clusters with guaranteed improvement – upper bound gets tighter
  
- ❑ How do we avoid re-solving?
  - ❑ “Warm start” of new messages using the old messages

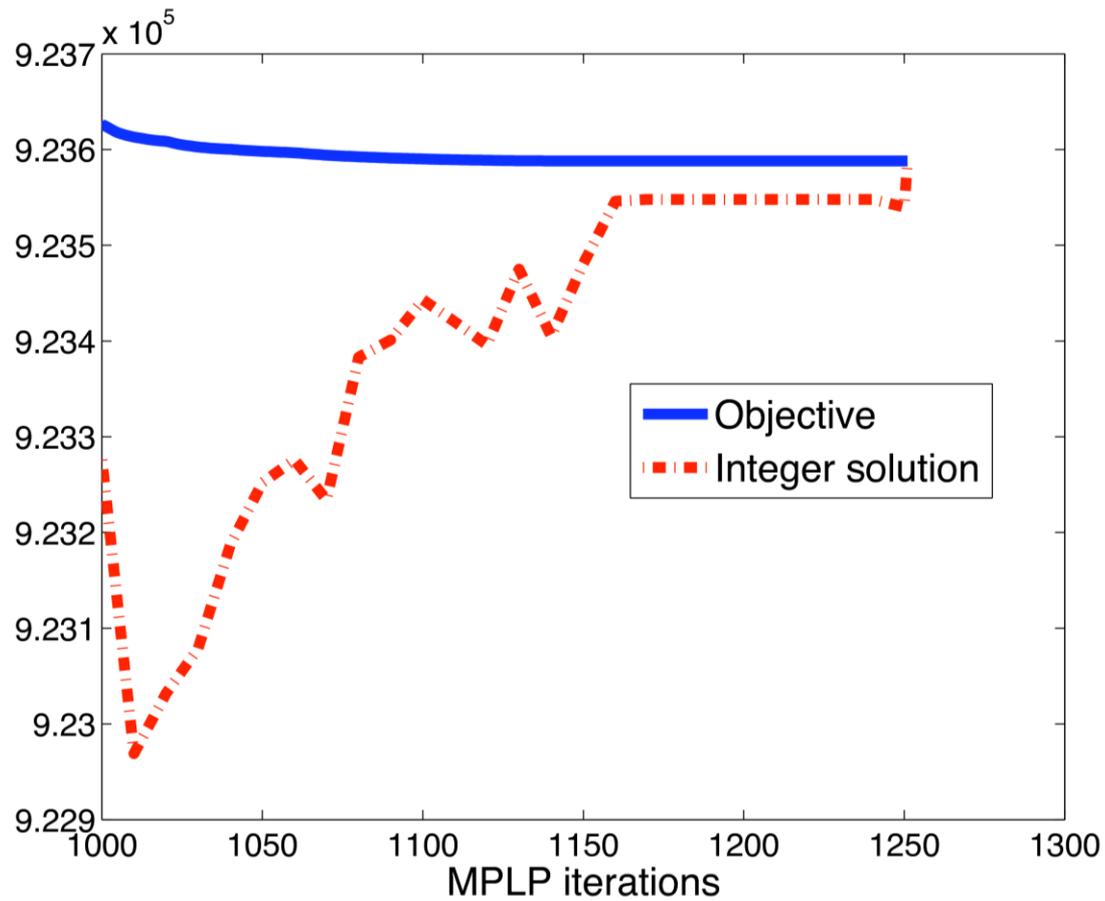
# Dual algorithm

1. Run message-passing
2. Decode assignment from messages
3. Choose a cluster to add to relaxation
4. Warm start: initialize new cluster messages

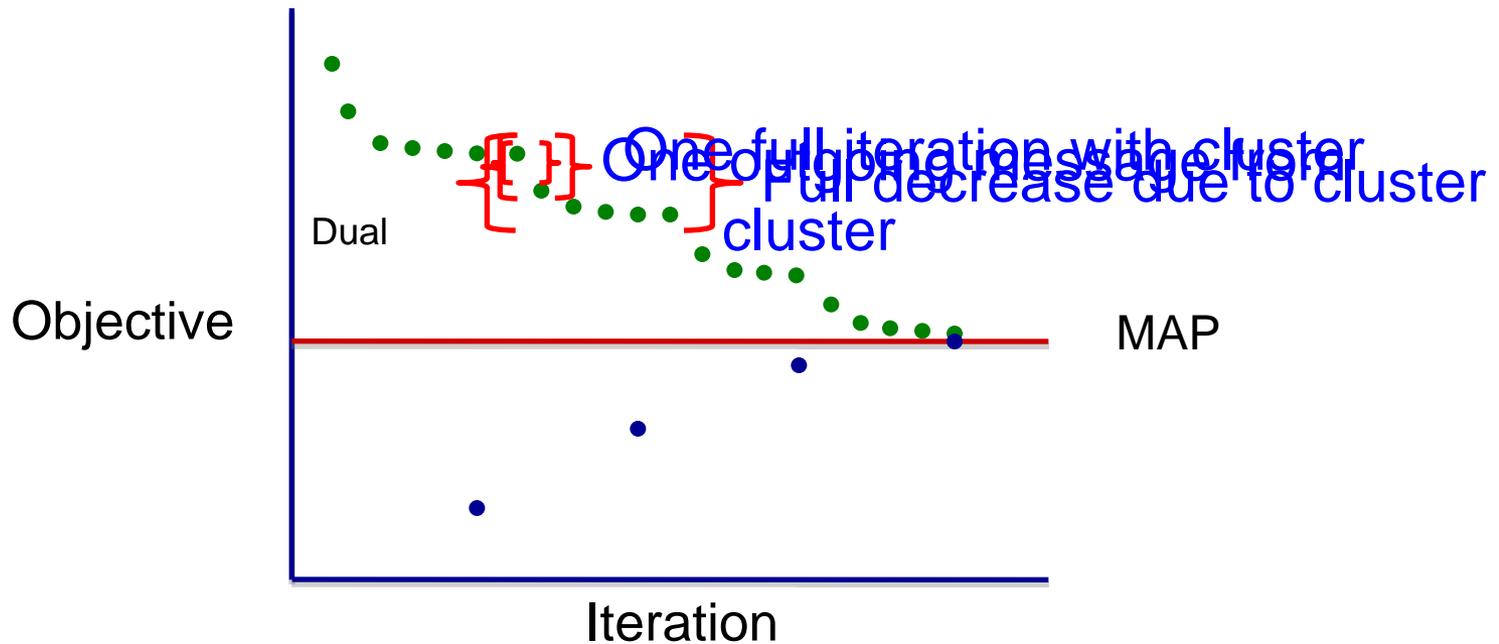


# Dual algorithm

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# What cluster to add next?



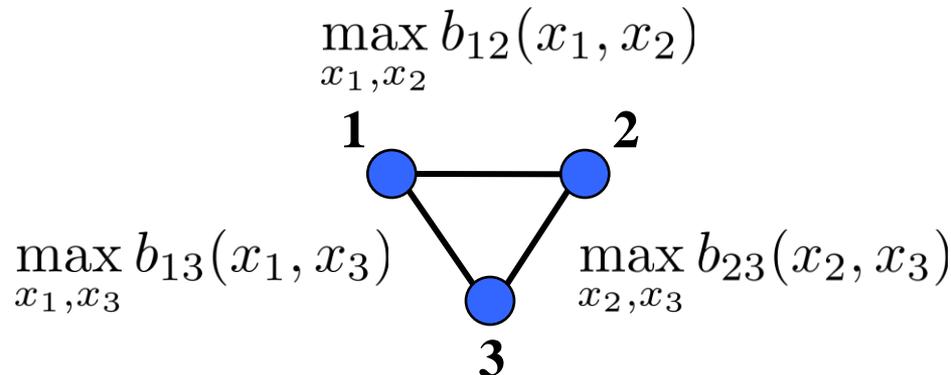
$$\sum_{e \in c} \max_{x_e} b_e(x_e) - \max_{x_c} \left[ \sum_{e \in c} b_e(x_e) \right]$$

# What cluster to add next?

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$$\sum_{e \in c} \max_{x_e} b_e(x_e) - \max_{x_c} \left[ \sum_{e \in c} b_e(x_e) \right]$$

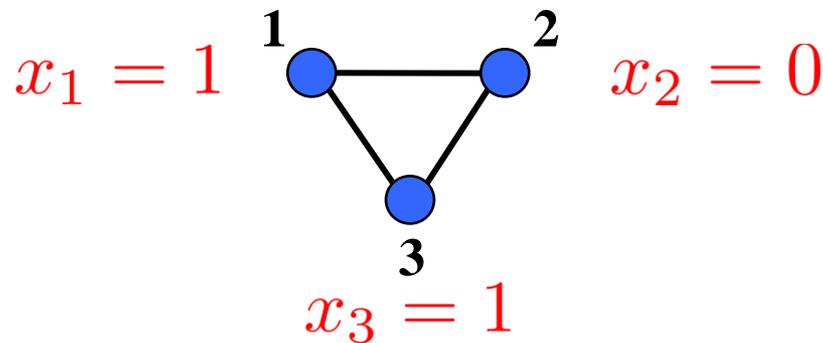
$$\max_{x_1, x_2, x_3} [b_{12}(x_1, x_2) + b_{23}(x_2, x_3) + b_{13}(x_1, x_3)]$$



# What cluster to add next?

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$$\underbrace{\sum_{e \in c} \max_{x_e} b_e(x_e)}_{3 * 99} - \underbrace{\max_{x_c} \left[ \sum_{e \in c} b_e(x_e) \right]}_{2 * 99 - 10}$$



If dual  $b_{ij}(x_i, x_j) = 99$  if  $x_i \neq x_j$  decreases, there was frustration  
 $b_{ij}(x_i, x_j) = -10$  otherwise

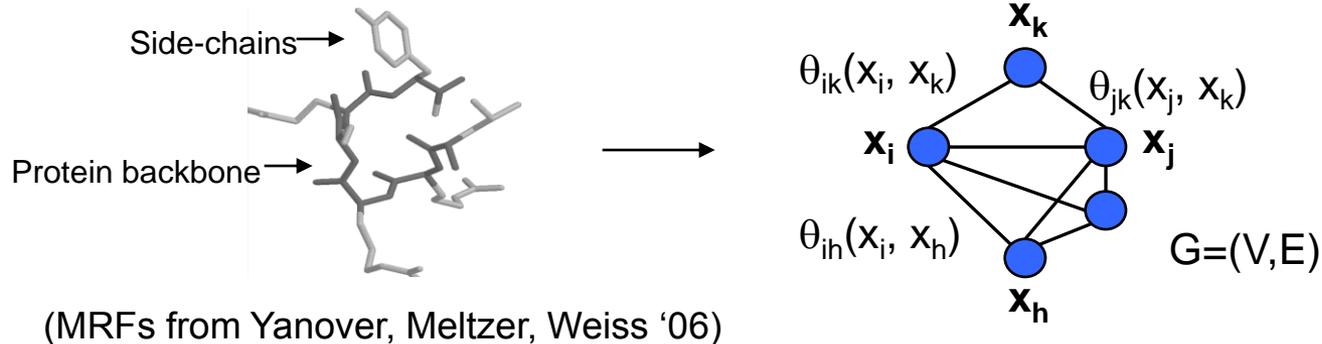
# Related Work

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- Region-pursuit algorithm for generalized BP (Welling UAI '04)
  - Iteratively adds the regions that most change region free energy (algorithmically very similar)
  - Found that sometimes adding regions gave worse results
  - Our approach circumvents this by working with the dual LP
  
- Cutting-plane algorithm using cycle inequalities (Sontag & Jaakkola '08)
  - Selection criteria of *constraint violation* instead of *bound minimization*
  - SJ can efficiently *find* violated constraints, but *re-solving* is hard in primal
  
- Other dual formulations (Werner '05, Kolmogorov & Wainwright '05, Johnson et al. '07, Komodakis et al. '07, Globerson & Jaakkola '07)
  
- Concurrently, similar approach proposed by (Werner CVPR

# Experiments: Protein design

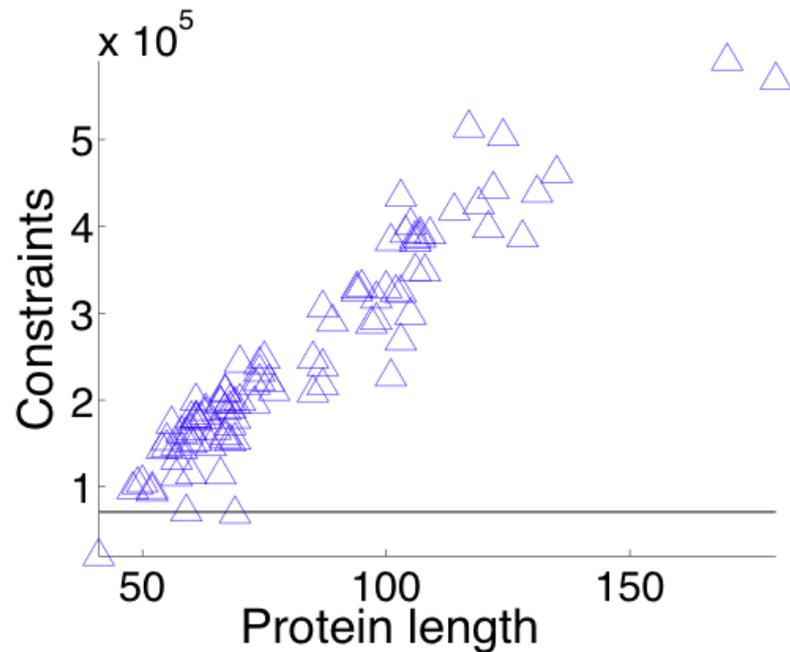
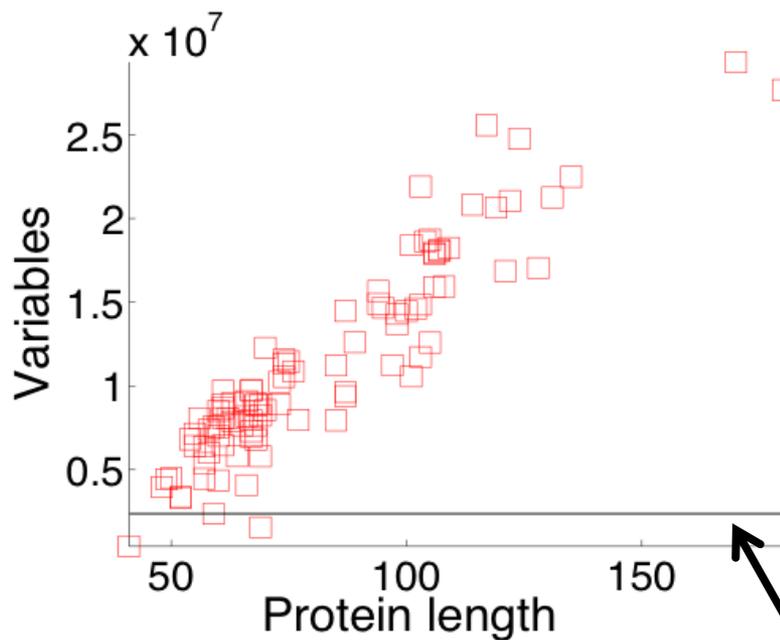
- Given protein's 3D shape, choose amino-acids giving the most stable structure



- Each state corresponds to a choice of amino-acid and side-chain angle
- MRFs have 41-180 variables, each variable with 95-158 states
- Hard to solve
  - Very large treewidth
  - Many small cycles (20,000 triangles) and frustration

# Primal LP, pairwise, is large

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CPLEX can only run on 3:  
must move to dual!

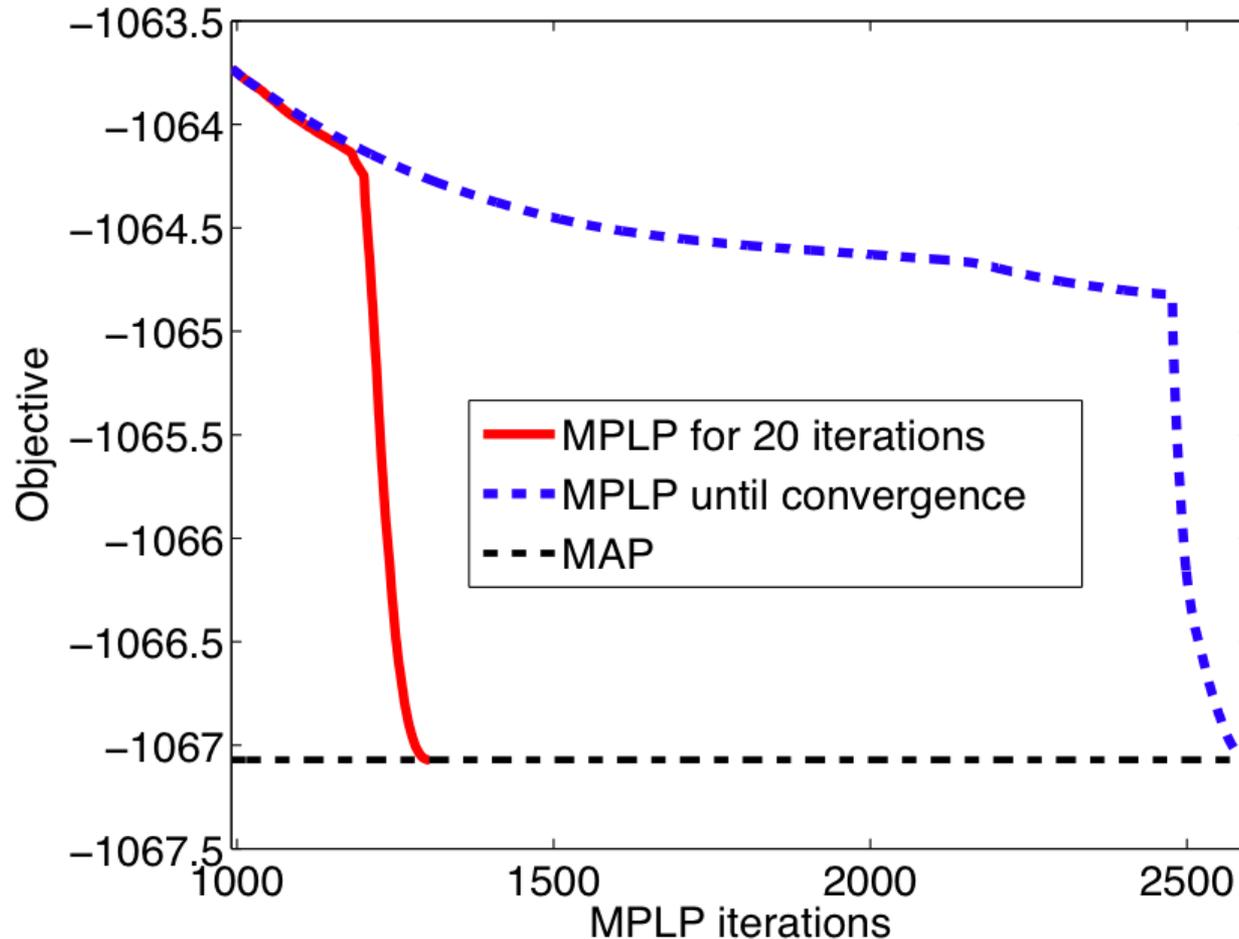
# Protein design results

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- Pairwise consistency solves only 2 of the 97 proteins  
(Yanover, Meltzer, Weiss, JMLR '06)
- With triplets, we solve 96 of 97 protein design problems  
(!!!)
- Between 5 and 735 triplets needed (median: 145)
  - Out of 20,000
  - Each triplet message needs >1 million computations
- 9.7 hours/problem (max: 11 days)

# Faster to stop before for convergence

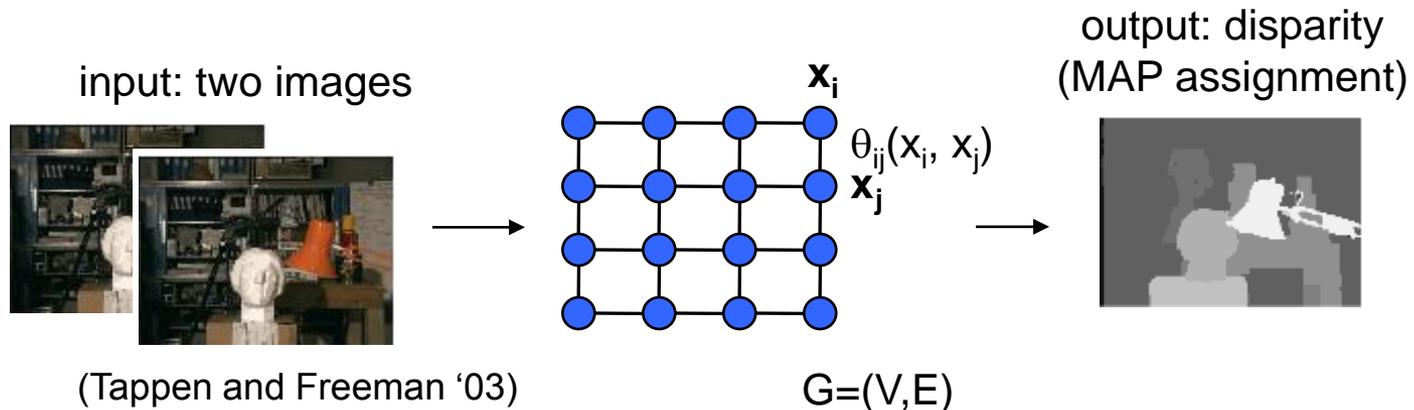
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# Experiments: Stereo vision

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- How far away are these objects?



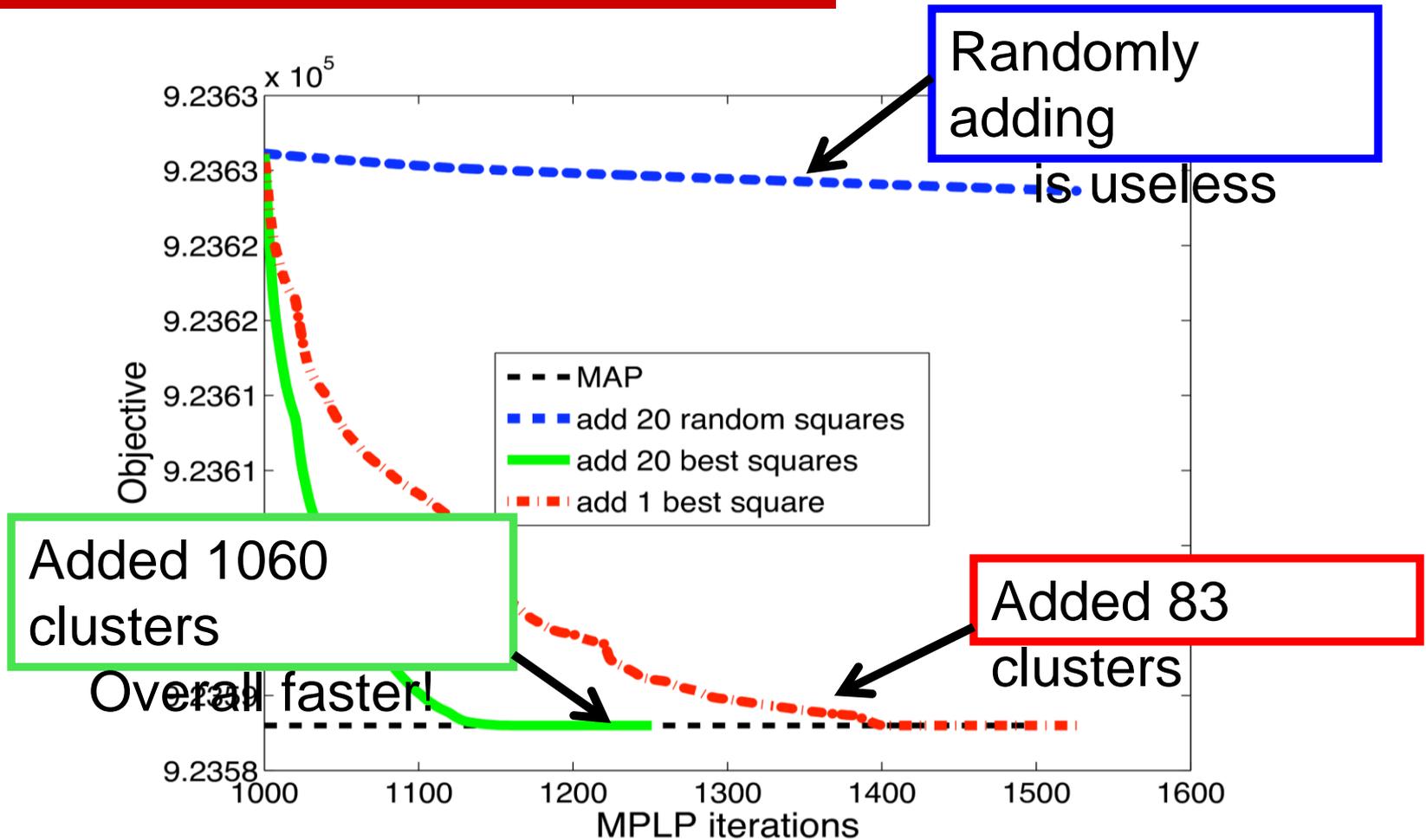
- 116 x154 pixels (13,000 variables), each with 16 states
- Hard to solve
  - Treewidth is over 230
  - Many short cycles: 13,000 squares (4-cycles)
  - Non-convex potentials

# Stereo vision results

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- 10 images: variations on 'Tsukuba' sequence
- Pairwise consistency solves 6 of the 10 images
- We solve all of them

# How aggressively to add clusters?



# Future work

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- ❑ Efficiently searching over clusters in the dual
- ❑ Structured prediction with large MRFs
- ❑ Extension to marginals and partition function
- ❑ Will soon release optimized code

# Conclusions

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- We give an algorithm to add clusters to message-passing
- Directly minimizes upper bound on MAP given by LP relaxation
- Using only a small number of clusters, solves some difficult real-world problems