

State-Sensitive Points-to Analysis for the Dynamic Behavior of JavaScript Objects

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Presented by ke tian

Outlines:

- **What is the problem?**

 - track the changes of object properties

- **What is the solution/contribution?**

 - state-sensitive points-to analysis

 - + a new control-flow graph representation

- **How efficient is the solution?**

 - significant improvement (+11% precision)

Background (__proto__)

Edit This Code:

See Result »

Result:

```
<!DOCTYPE html>
<html>
<body>

<p>Test By Ke</br></p>
<p id="demo"></p>

<script>

function Foo(y){this.y = y;}

var b = new Foo(20);
b.c = 30;

document.getElementById("demo").innerHTML =
(b.__proto__ ) + '</br>' +
(b.c) + '</br>' +
(b.y) + '</br>' +
(b.__proto__ == Foo.prototype) + '</br>' +
(b.constructor)
;

</script>

</body>
</html>
```

Test By Ke

[object Object]

30

20

true

function Foo(y){this.y = y;}

b	
c	30
__proto__	

Foo.prototype

y	20
constructor	function Function() { [native code] }
__proto__	

Object.prototype

How to (formally) describe a JavaScript Object?

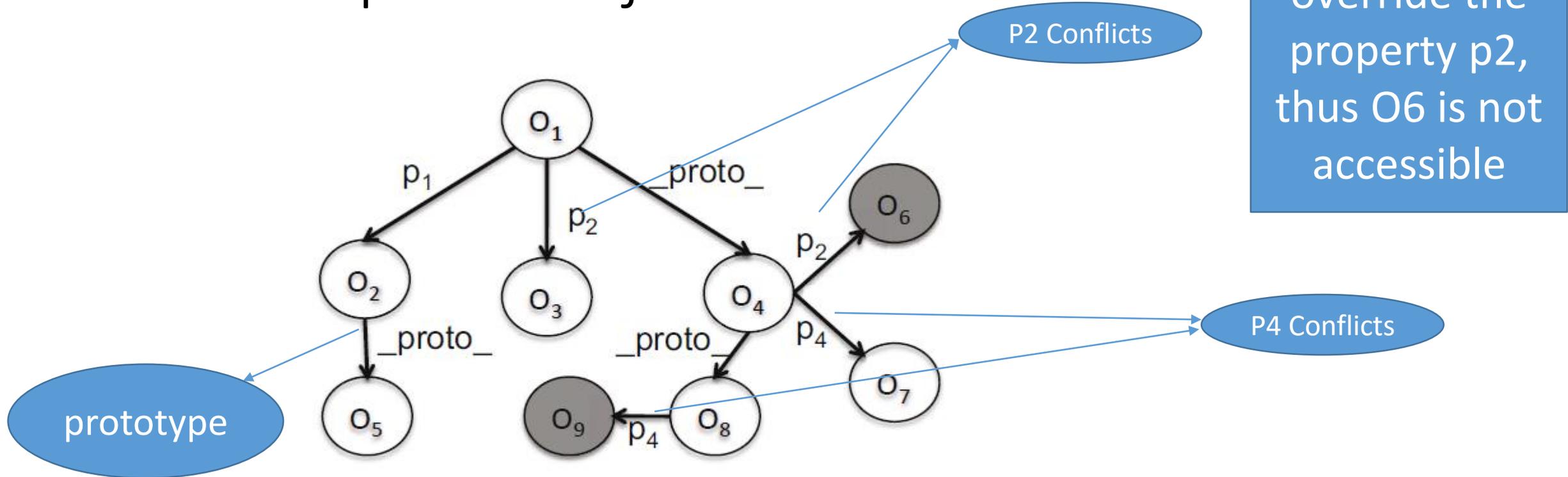
Definition 1. The *obj-ref state* at a program point denotes all of its accessible properties and their non-primitive values.

Def 1. is used to describe a type (not constant) of a JavaScript object

Definition 2. *State-update statements* are: (1) property write statement (i.e., $x.p = y$ or $x[p'] = y$), (2) property delete statement (i.e., *delete* $x.p$ or *delete* $x[p']$), and (3) an invocation that directly or indirectly results in execution of (1) and/or (2).

Def 2. Write and delete operations can result in state-update, affect obj-ref states

An example of obj-ref state



Local O3
override the
property p2,
thus O6 is not
accessible

Fig. 1. *obj-ref state* for O_1 . (Unshaded nodes only)

$\text{Obj-Ref}(O_1) = \{O_1, O_2, O_3, O_4, O_5, O_7, O_8\}$

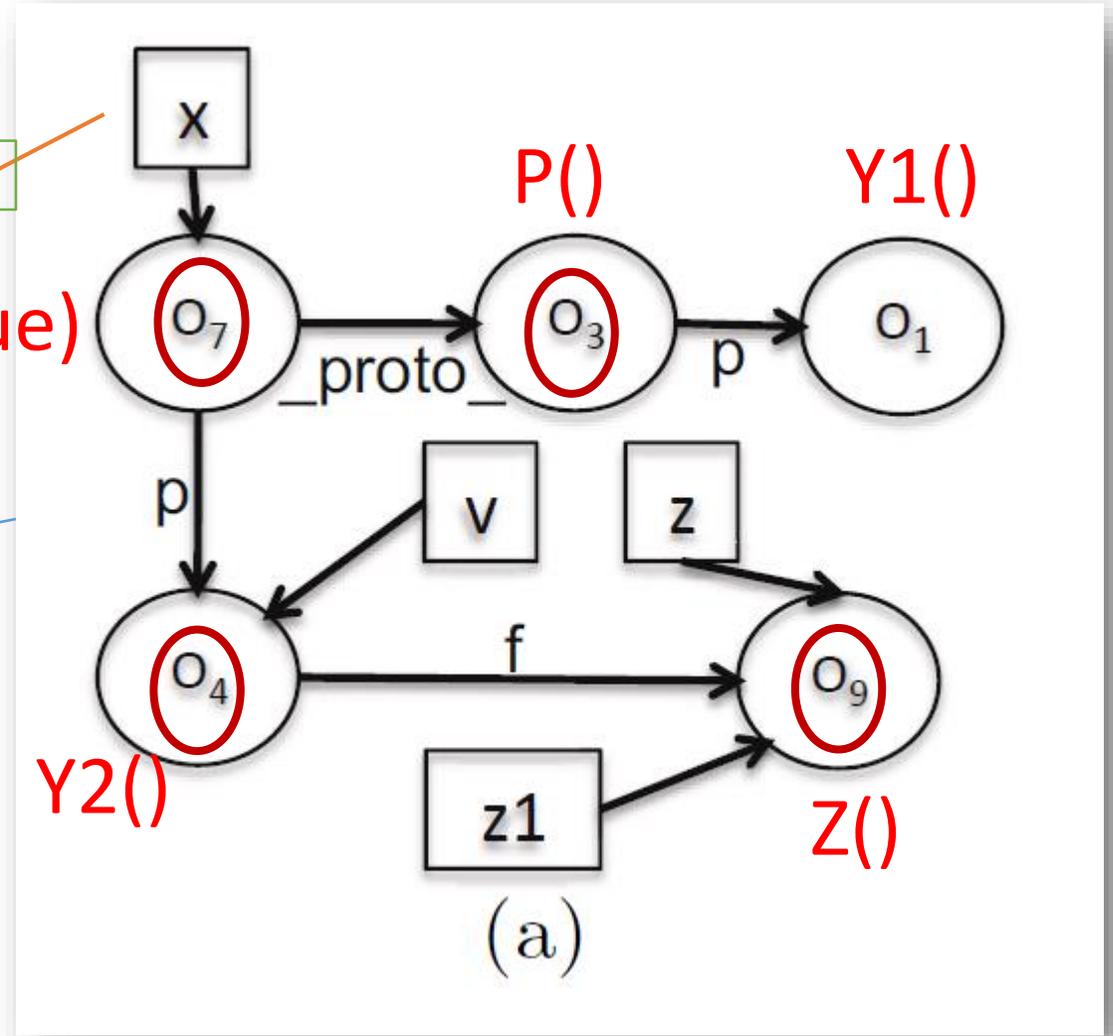
Motivating example (the problem)

```
1 function P(){ this.p = new Y1(); }
2 function X(b){
3   this.__proto__ = new P();
4   if(b) { this.p = new Y2(); }
5   else this.q = new Y3();
6 }
7 var x = new X(true);
8 x.bar = function(v, z){ v.f = z; }
9 var z1 = new Z();
10 x.bar(x.p, z1);
11 ...
12 x.p = new A();
13 ...
14 var z2 = new Z();
15 x.bar(x.p, z2);
```

create

X(true)

Line 10

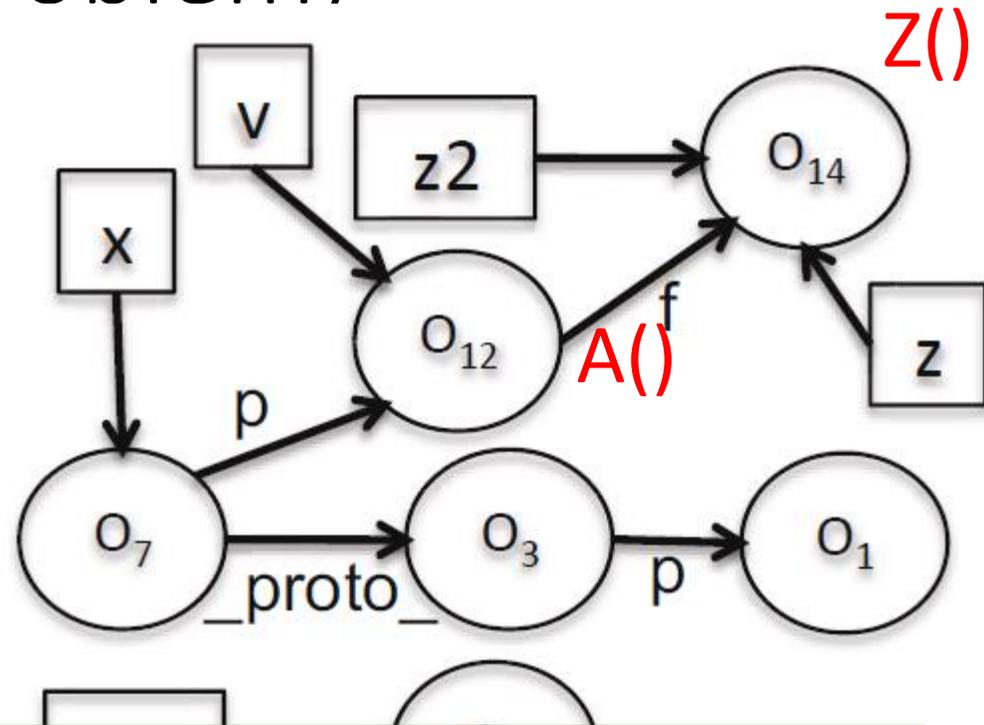


Obj-Ref(O7)={O7, O3, O4, O9}

Run-time points-to graph at line 10

Motivating example (the problem)

```
1 function P(){ this.p = new Y1(); }
2 function X(b){
3   this.__proto__ = new P();
4   if(b) { this.p = new Y2(); }
5   else this.q = new Y3();
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7 var x = new X(true);
8 x.bar = function(v, z){ v.f = z; }
9 var z1 = new Z();
10 x.bar(x.p, z1);
11 ...
12 x.p =
13 ...
14 var z2 = new Z();
15 x.bar(x.p, z2);
```

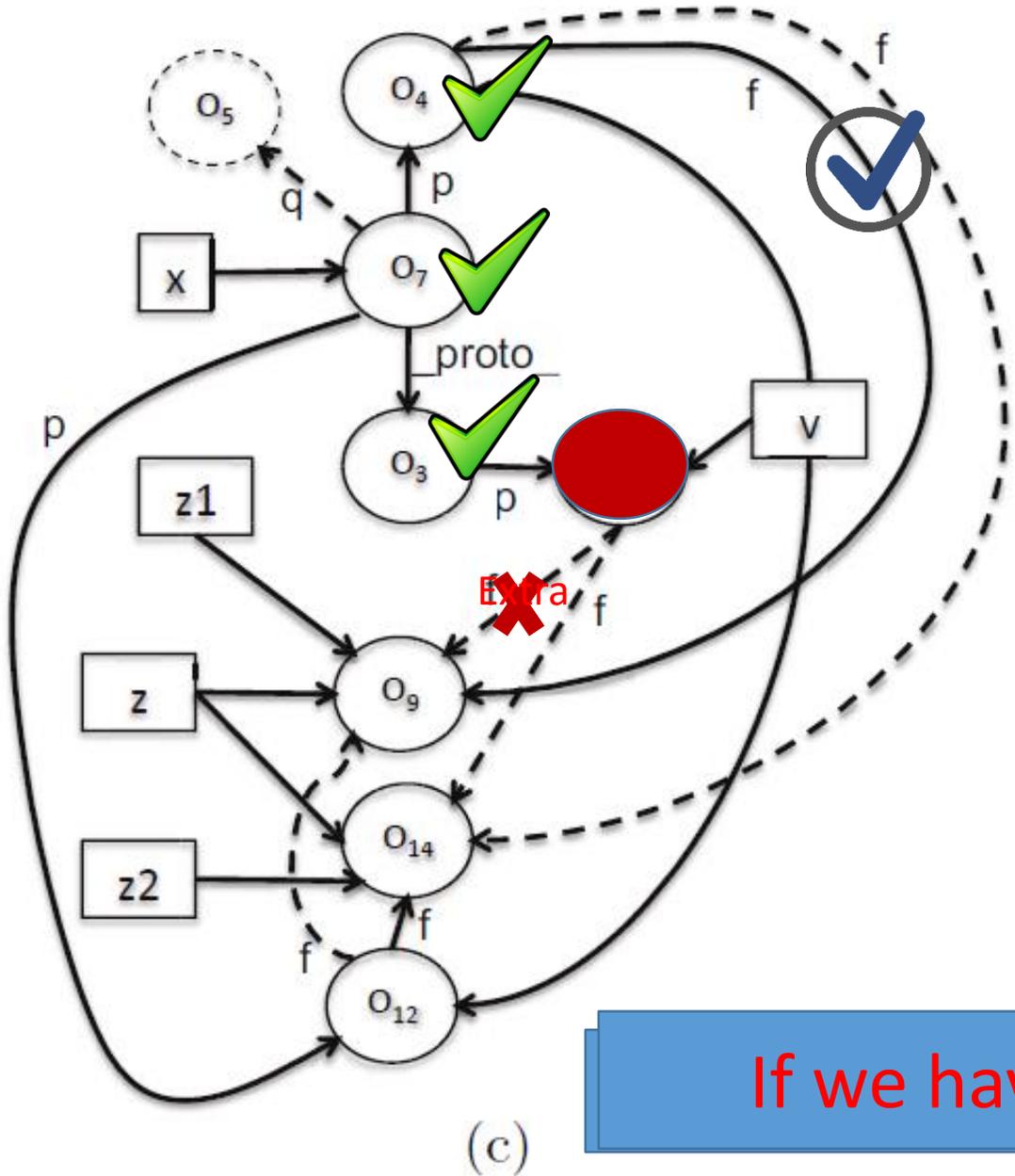


Conclusion: Obj-Ref(O7) can be changed during the runtime

Obj-Ref(O7)={O7, O12, O3, O14}

Run-time points-to graph at line 15

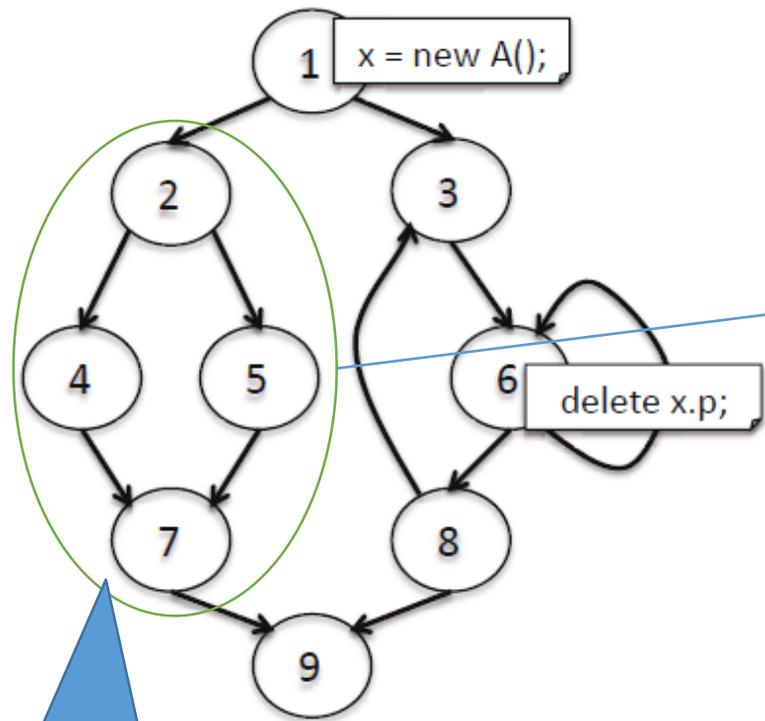
What is the problem? (imprecision)



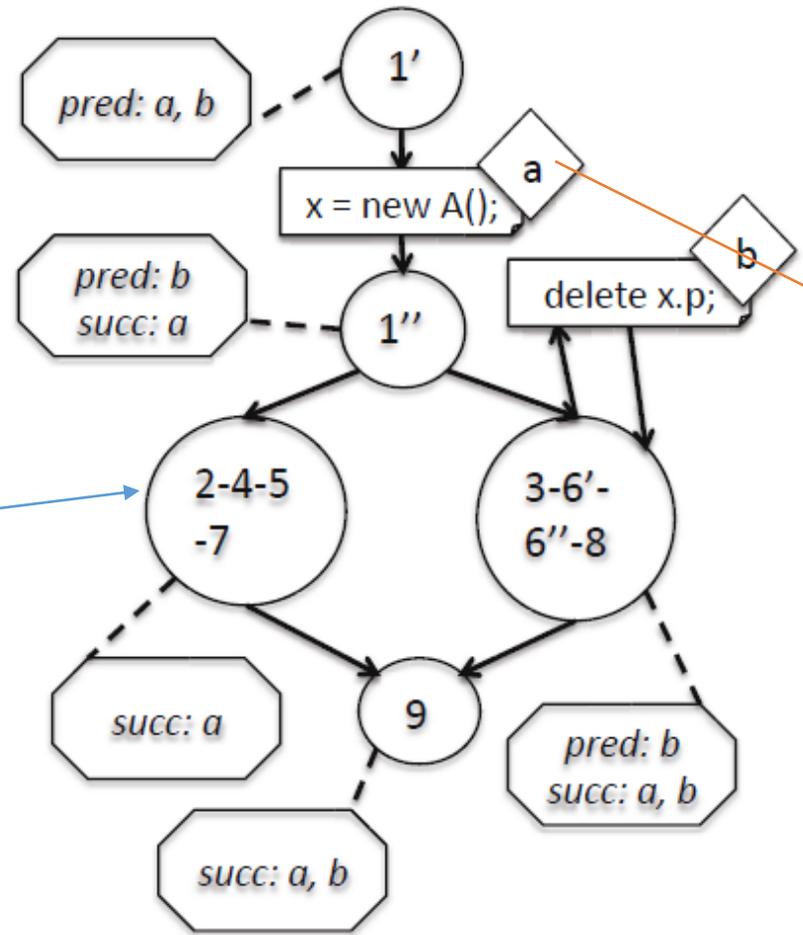
- 1 constructor polymorphism (O7.q)
- 2 Object property change
- ****3 Function invocation****
x.bar(x.p,z1) [line 10]
x.p.f = z1(O9)
E.g., not knowing O4 exists
extra : (O1 -> O9)

If we have $\text{Obj-Ref}(O7) = \{O7, O3, O4\}$ line 9

What is the Solution? (state-preserving block graph)



State-preserving Node



- Split the CFG based on state-update

Statement (new/delete)

1 -> (1', x=newA(), 1'')

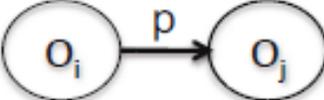
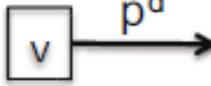
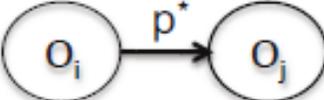
- Aggregate the state-preserving nodes in the graph

Fig. 4. SPBG generation. (a) CFG. (b) SPBG.

Partial flow sensitive

Points-to graph representation

Table 1. Expanded points-to graph with annotations

points-to graph G	node N	variable v	
		abstract object o	
		in-construction object $@o$	
	edge E	variable reference $(v, \phi o)$	
		property reference $(\langle \phi o_i, p \rangle, \phi o_j)$	
		access path $(\langle v, p \rangle, \phi o)$	
	annotation A	d annotation p^d	
		$*$ annotation p^*	

Used for object creation
(in transfer function)

$$\Phi o = \{ @o, o \}$$

Traditional points-to graph

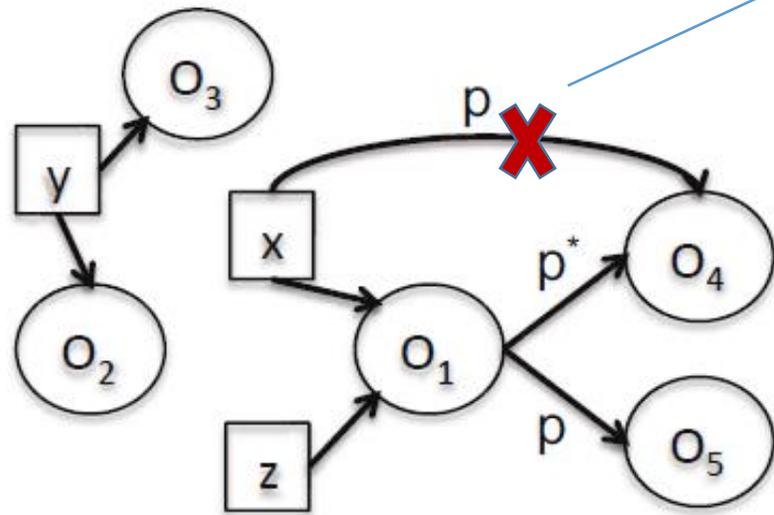
Help us to (strongly) determine
the property of variable

P^* means the relation
MAY not exist (safety)

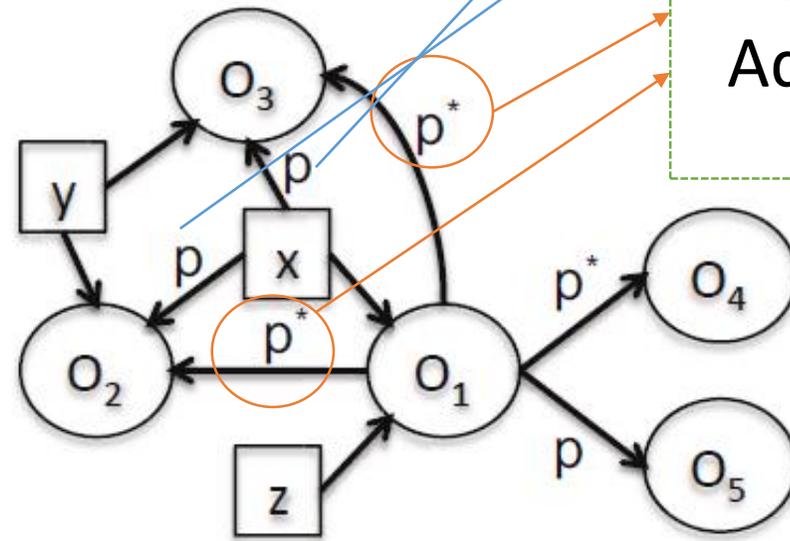
Transfer functions (update points-to graph)

- Object creation ($x = \text{new } X(a_1, a_2, \dots, a_n)$)
- Property write ($x.p = y$) **[example]**
- Property delete (delete $x.p$)
- Direct write ($x = y$)
- Property read ($x = y.p$)
- Method invocation ($x = y.m(a_1, a_2, \dots, a_n)$)

Transfer rule (example $x.p = y$)



(a)

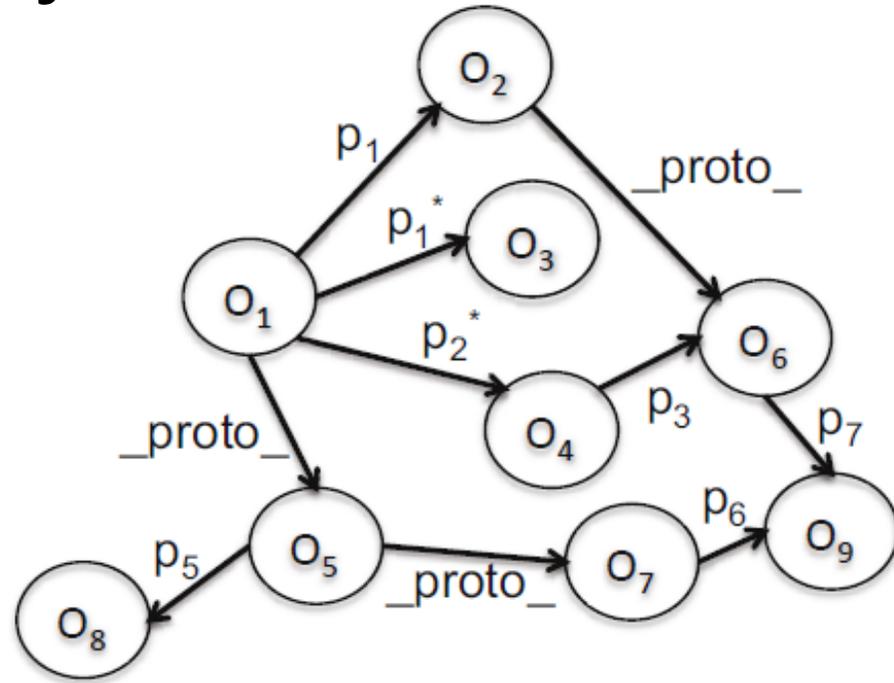


(b)

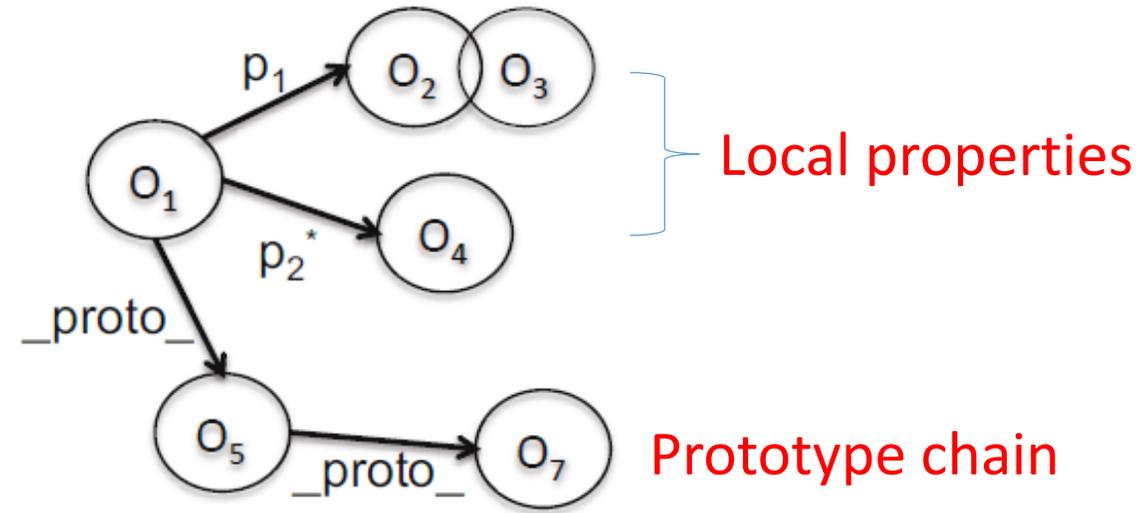
Strong update:
Delete $\langle x, p \rangle, O_4$
Add $\langle x, p \rangle, O_3$
Add $\langle x, p \rangle, O_2$
Weak update:
Add $\langle O_1, p^* \rangle, O_3$
Add $\langle O_1, p^* \rangle, O_2$

Fig. 5. Property write example. (a) Input points-to graph. (b) Updated points-to graph.

Approximation (reduce analysis overhead) use obj-ref state as context



(a)

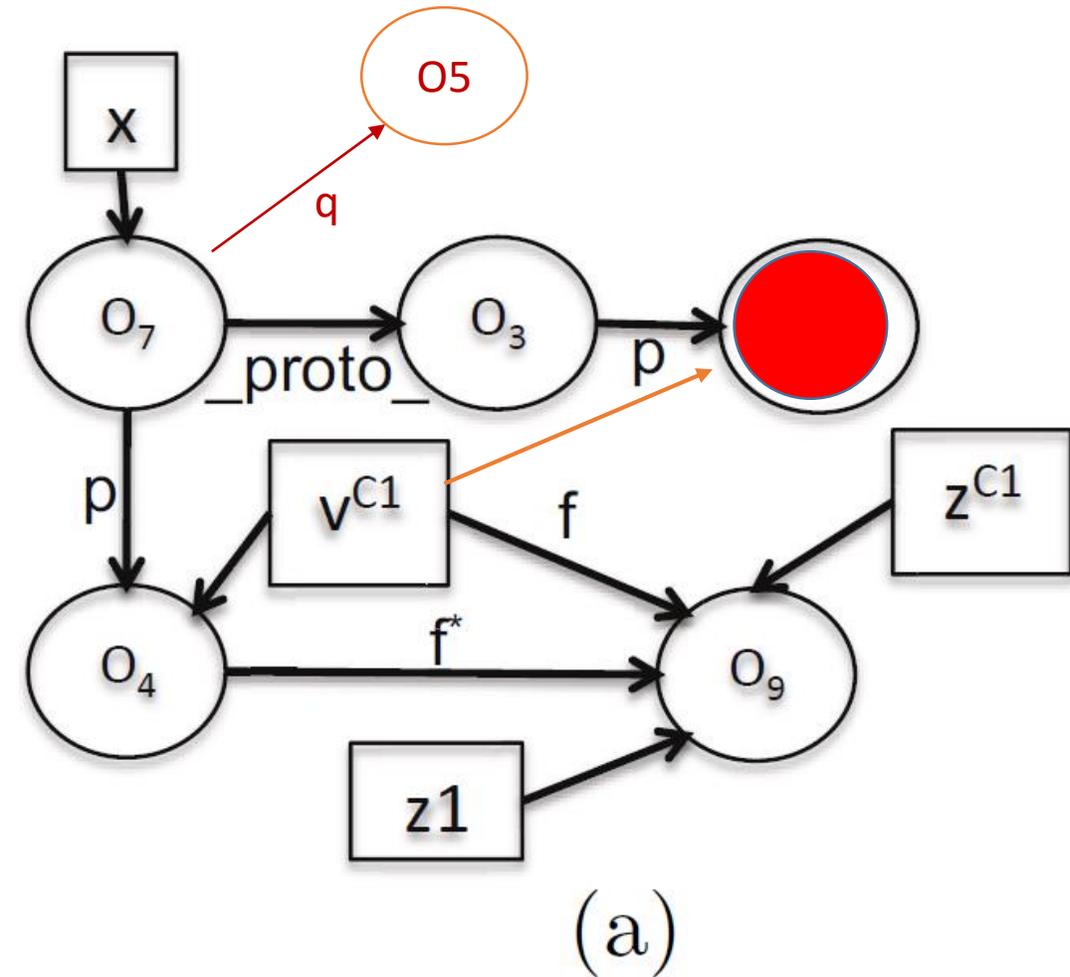


(b)

Fig. 6. Approximate *obj-ref state* as a context. (a) *obj-ref state* of O_1 . (b) Approximate *obj-ref state* of O_1 .

Trade-off :Lose the (some) precision but increase scalability

How efficient is the solution?

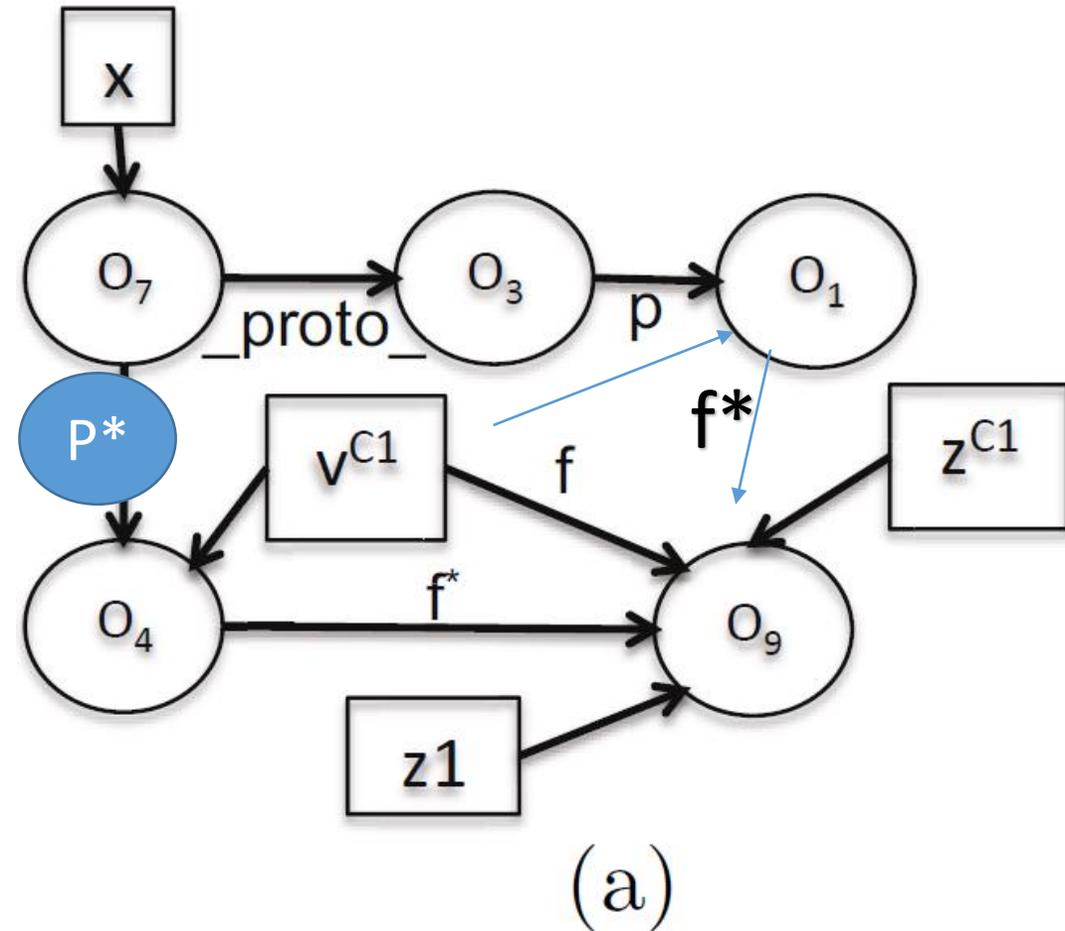


```
1  Approximate Obj-Ref(O7)      ; }
2  C1= {O7,
3      p:O4, (NO O1 here!!!)
4      __proto__: O3
5  }
6  }
7
8  x.bar = function(v, z){ v.f = z; }
9  var z1 = new Z();
10 x.bar(x.p, z1);
11 ...
12 x.p = new A();
13 ...
14 var z2 = new Z();
15 x.bar(x.p, z2);
```

insensitive Points-to graph at line 10

Fig. 2. JavaScript example

How efficient is the solution?



Points-to graph at line 10 (if p*)

What if $O7 \xrightarrow{p^*} O4$?

$C1 = \{O7,$

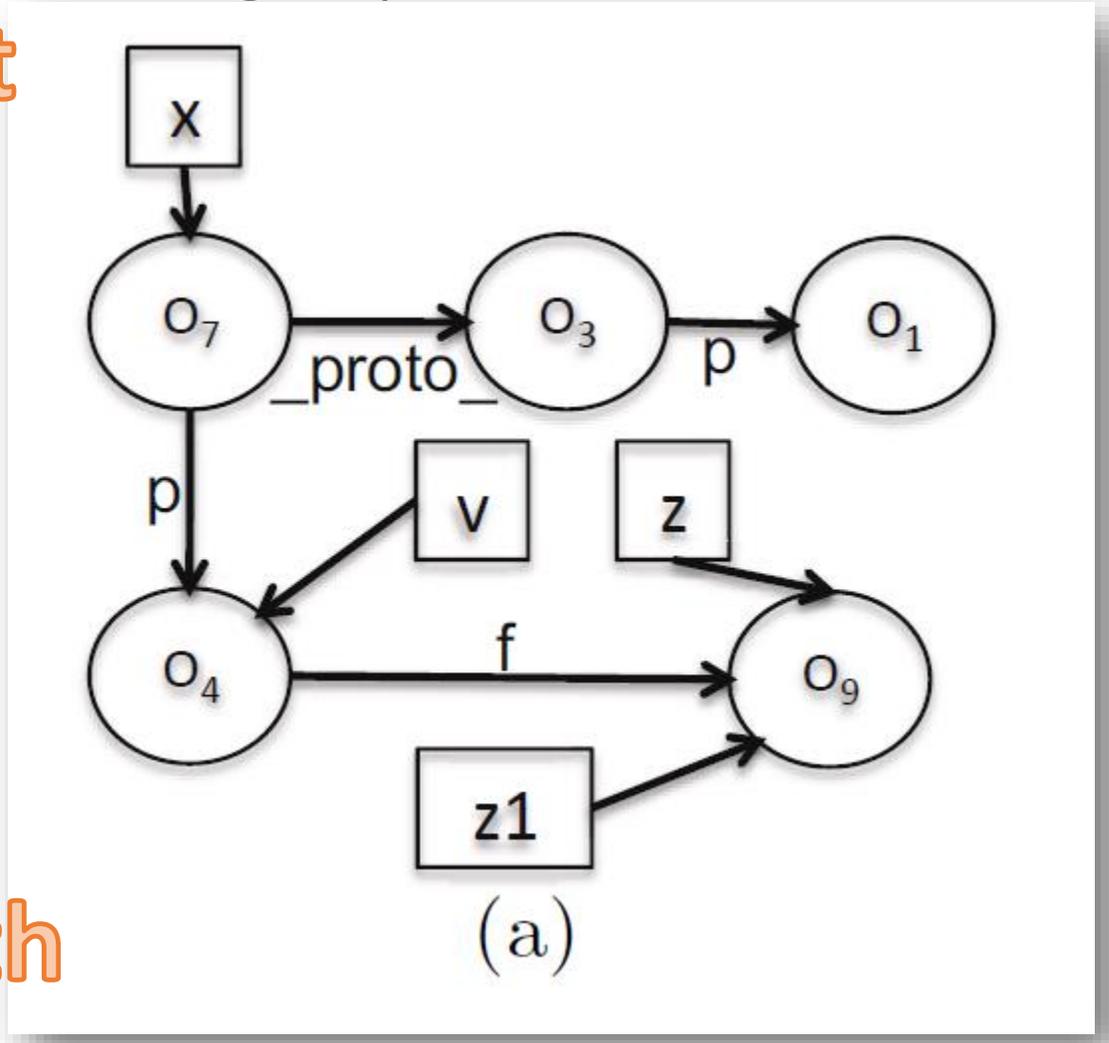
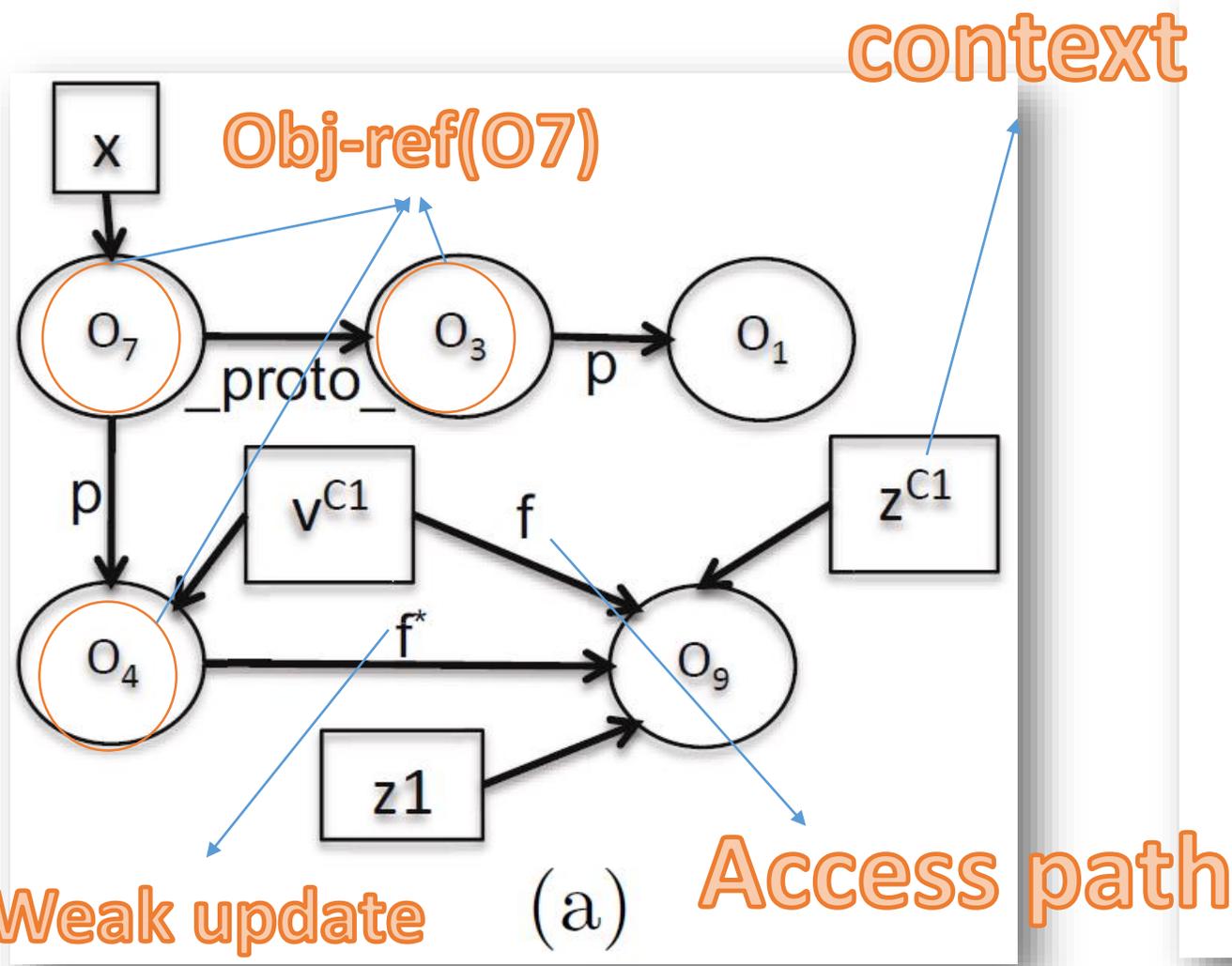
$p: O4, O1,$

$__proto__: O3\}$

```
8 x.bar = function(v, z){ v.f = z; }
9 var z1 = new Z();
10 x.bar(x.p, z1);
11 ...
12 x.p = new A();
13 ...
14 var z2 = new Z();
15 x.bar(x.p, z2);
```

Fig. 2. JavaScript example

Compare with the ground truth graph



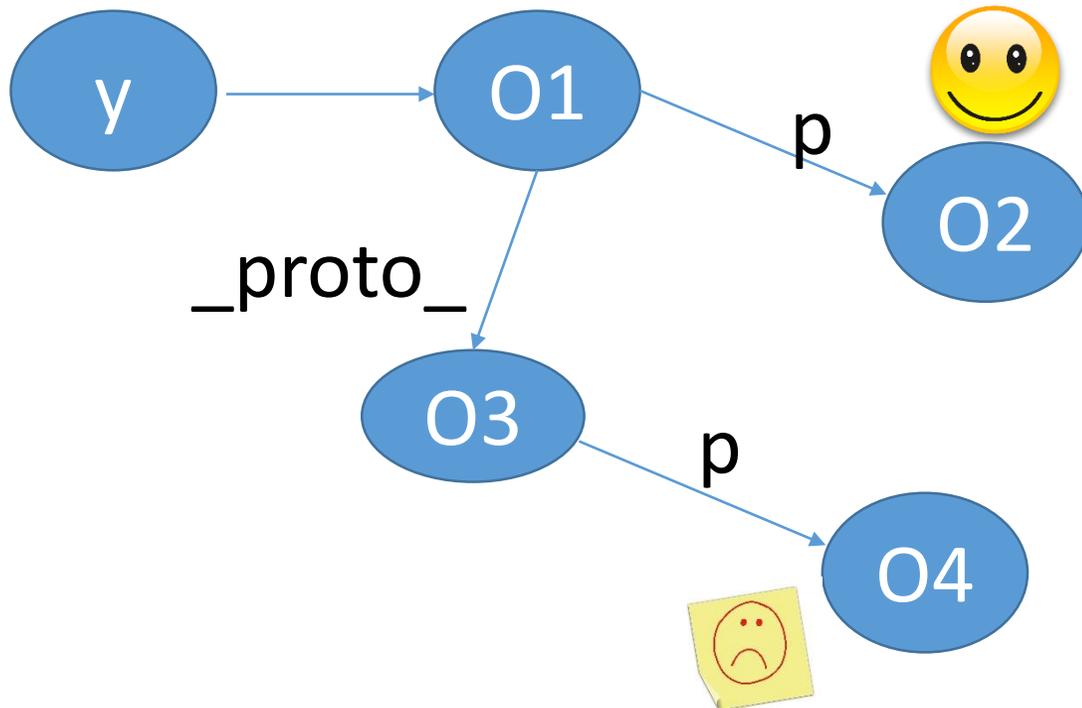
New Points-to graph at line 10 (the authors' approach) Figure 7(a)

Run-time Points-to graph at line 10 Fig 3(a)

Measurement (# of return objects)

$x = y.p$:

Return {O2} better than return {O2,O4}



$x = y.p/x=y.p(...)$
[statement s /context c]
REF(s,c) = How many
objects will be returned
Through lookuping

The smaller #,
the better precision

Table 4. *REF* analysis precision

Website	<i>Corr</i>			<i>CorrBSSS</i>		
	1	2-4	≥ 5	1	2-4	≥ 5
facebook.com	38%	52%	10%	50%	47%	3%
google.com	32%	51%	17%	53%	42%	5%
youtube.com	41%	47%	12%	54%	41%	5%
yahoo.com	48%	46%	6%	52%	45%	3%
wikipedia.org	29%	45%	26%	43%	39%	18%
amazon.com	45%	52%	3%	46%	51%	3%
twitter.com	32%	53%	15%	39%	49%	12%
blogspot.com	35%	34%	31%	53%	36%	11%
linkedin.com	34%	49%	17%	44%	50%	6%
msn.com	40%	36%	24%	48%	37%	15%
ebay.com	30%	40%	30%	46%	40%	14%
bing.com	41%	34%	25%	54%	37%	9%
<i>Geom. Mean</i>	37%	44%	15%	48%	43%	7%

Corr: correlation tracking ...
CorrBSSS: the authors' approach

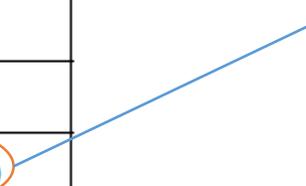
Better performance
48-37 % =11%

Overhead (acceptable)

Table 5. *REF* analysis cost (in seconds) on average per webpage

Website	<i>Corr</i>	<i>CorrBSSS</i>	overhead
facebook	17.4	45.9	163%
google	13.0	30.4	134%
youtube	31.2	75.3	141%
yahoo	28.5	54.1	90%
wiki	16.0	40.1	151%
amazon	15.1	24.2	61%
twitter	38.1	94.5	148%
blog	15.9	42.4	137%
linkedin	27.8	62.0	167%
msn	34.4	57.9	68%
ebay	8.3	27.2	227%
bing	22.1	50.4	128%
<i>Geom. Mean</i>	20.4	46.7	127%

Average overhead



Q&A

- What is “state-sensitive”? and the relation with context sensitivity.

Context sensitivity

Object

Call-site ...

State ...

- Others....