

## Canonical Forms

- Canonical Means “Unique”
- All Possible Functions can be Expressed in One and Only One Way in a Canonical Form
- Canonical Forms may be Circuit Diagrams or Algebraic Equations

## Minterms and Maxterms

- Consider a function of Three variables  $x$ ,  $y$ , and  $z$
- Since each Variable may be Complemented or Uncomplemented there are  $2^3=8$  Different Combinations
- When Combinations are Combined with AND they are Called Minterms
- When Combinations are combined with OR they are Called Maxterms

# Minterms and Maxterms

**Table 2-3**  
Minterms and Maxterms for Three Binary Variables

x	y	z	Minterms		Maxterms	
			Term	Designation	Term	Designation
0	0	0	$x'y'z'$	$m_0$	$x + y + z$	$M_0$
0	0	1	$x'y'z$	$m_1$	$x + y + z'$	$M_1$
0	1	0	$x'yz'$	$m_2$	$x + y' + z$	$M_2$
0	1	1	$x'yz$	$m_3$	$x + y' + z'$	$M_3$
1	0	0	$xy'z'$	$m_4$	$x' + y + z$	$M_4$
1	0	1	$xy'z$	$m_5$	$x' + y + z'$	$M_5$
1	1	0	$xyz'$	$m_6$	$x' + y' + z$	$M_6$
1	1	1	$xyz$	$m_7$	$x' + y' + z'$	$M_7$

For  $n$  Variables there are  $2^n$  Minterms/Maxterms

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## Sum-of-Minterms Form

- Canonical Form – *Standard Products*
- Determine the Set of Minterms for Which a Function is 1-valued
  - These are called “Minterms of the Function”
- Combine all Minterms with a + Operation
- This is a 2-Level Form

## Sum-of-Minterms Example

$x$	$y$	$z$	$f_1$	
0	0	0	0	
0	0	1	1	← $\overline{\overline{x}} \overline{\overline{y}} z$
0	1	0	0	
0	1	1	0	
1	0	0	1	← $x \overline{\overline{y}} \overline{\overline{z}}$
1	0	1	0	
1	1	0	0	
1	1	1	1	← $x y z$

$$f_1 = \overline{\overline{x}} \overline{\overline{y}} z + x \overline{\overline{y}} \overline{\overline{z}} + x y z$$

$$f_1 = \sum (1, 4, 7)$$

## Product-of-Maxterms Form

- Canonical Form – *Standard Sums*
- Determine the Set of Maxterms for Which a Function is 0-valued
  - These are called “Maxterms of the Function”
- Must Complement Each Literal
- Combine all Maxterms with a • Operation
- This is a 2-Level Form

## Product-of-Maxterms Example

$x$	$y$	$z$	$f_1$	
0	0	0	0	← $x + y + z$
0	0	1	1	
0	1	0	0	← $x + \bar{y} + \bar{z}$
0	1	1	0	← $x + y + z$
1	0	0	1	
1	0	1	0	← $\bar{x} + \bar{y} + \bar{z}$
1	1	0	0	← $\bar{x} + y + \bar{z}$
1	1	1	1	

$$f_1 = (x + y + z)(x + \bar{y} + z)(x + \bar{y} + \bar{z})$$

$$(\bar{x} + y + \bar{z})(\bar{x} + \bar{y} + z)$$

$$f_1 = \prod (0, 2, 3, 5, 6)$$

## Other Notation

$$f_1 = \sum (1, 4, 7) = m_1 + m_4 + m_7$$

$$f_1 = \prod (0, 2, 3, 5, 6) = M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6$$

What is the function:

$$f = m_0 + m_1 + m_2 + m_3 + m_4 + m_5 + m_6 + m_7$$

## Conversion of Canonic Forms

$$f_1 = \sum (1, 4, 7) = m_1 + m_4 + m_7$$

$$\overline{f_1} = \sum (0, 2, 3, 5, 6) = m_0 + m_2 + m_3 + m_5 + m_6$$

$$\overline{\overline{f_1}} = \overline{m_0 + m_2 + m_3 + m_5 + m_6}$$

$$f_1 = m_0 + m_2 + m_3 + m_5 + m_6$$

DeMorgan's Theorem:

$$f_1 = \overline{m_0} \cdot \overline{m_2} \cdot \overline{m_3} \cdot \overline{m_5} \cdot \overline{m_6}$$

$$\overline{m_i} = M_i$$

$$f_1 = M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6$$

## Standard Forms

- Canonical Forms USUALLY NOT Smallest (in terms of literals)
- Each minterm/maxterm contains  $n$  literals
- Standard Forms Contain Terms with  $n$  or Fewer Literals
- Sum-Of-Products (SOP) form
- Product-Of-Sums (POS) form
- These are Also Two-level Forms

## Standard Forms Examples

$$F_1 = \bar{y} + xy + \bar{x} y \bar{z} \quad F_2 = x(\bar{y} + z)(\bar{x} + y + \bar{z})$$

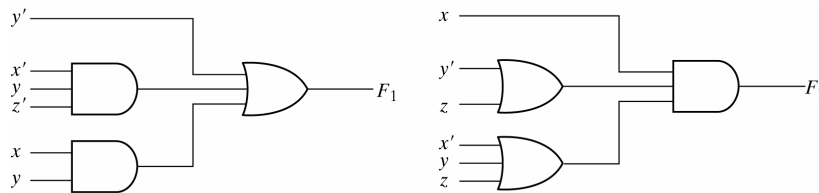


Fig. 2-3 Two-level implementation

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## Standard Forms

- Can Use Algebra to Find a Standard Form from a Canonical Form
- We Will Learn Other Methods to do this
- Commonly Known as “Simplification”
  - Seems Easy for Small Functions
  - Computationally Complex
  - Classic Problem in Switching Theory

## Multi-level Forms

- All Possible Functions can be Expressed in a Standard Two-level Form
- Multi-level Forms have more than 2 Levels

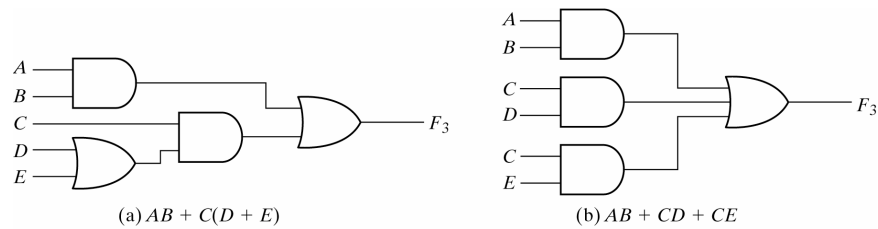


Fig. 2-4 Three- and Two-Level implementation

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## All Functions of 2-Variables

- There are  $2^{2^n}$  functions of  $n$  Variables
- AND and OR Happen to be Two of 16 Possible Functions of 2 Variables

**Table 2-7**  
*Truth Tables for the 16 Functions of Two Binary Variables*

$x$	$y$	$F_0$	$F_1$	$F_2$	$F_3$	$F_4$	$F_5$	$F_6$	$F_7$	$F_8$	$F_9$	$F_{10}$	$F_{11}$	$F_{12}$	$F_{13}$	$F_{14}$	$F_{15}$
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1

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# Names and Symbols of Functions









**Table 2-8**  
Boolean Expressions for the 16 Functions of Two Variables

Boolean functions	Operator symbol	Name	Comments
$F_0 = 0$		Null	Binary constant 0
$F_1 = xy$	$x \cdot y$	AND	$x$ and $y$
$F_2 = xy'$	$x/y$	Inhibition	$x$ , but not $y$
$F_3 = x$		Transfer	$x'$
$F_4 = x'y$	$y/x$	Inhibition	$y$ , but not $x$
$F_5 = y$		Transfer	$y$
$F_6 = xy' + x'y$	$x \oplus y$	Exclusive-OR	$x$ or $y$ , but not both
$F_7 = x + y$	$x + y$	OR	$x$ or $y$
$F_8 = (x + y)'$	$x \downarrow y$	NOR	Not-OR
$F_9 = xy + x'y'$	$(x \oplus y)'$	Equivalence	$x$ equals $y$
$F_{10} = y'$	$y'$	Complement	Not $y$
$F_{11} = x + y'$	$x \supset y$	Implication	If $y$ , then $x$
$F_{12} = x'$		Complement	Not $x$
$F_{13} = x' + y$	$x \supset y$	Implication	If $x$ , then $y$
$F_{14} = (xy)'$	$x \uparrow y$	NAND	Not-AND
$F_{15} = 1$		Identity	Binary constant 1

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## Common Circuits

- Certain Subsets Form ANY Function
  - AND, OR, NOT
  - NAND
  - NOR
  - AND, XOR

Name	Graphic symbol	Algebraic function	Truth table															
AND		$F = xy$	<table border="1"> <thead> <tr><th>x</th><th>y</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	x	y	F	0	0	0	0	1	0	1	0	0	1	1	1
x	y	F																
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0	1	0																
1	0	0																
1	1	1																
OR		$F = x + y$	<table border="1"> <thead> <tr><th>x</th><th>y</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	x	y	F	0	0	0	0	1	1	1	0	1	1	1	1
x	y	F																
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Inverter		$F = x'$	<table border="1"> <thead> <tr><th>x</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>1</td></tr> <tr><td>1</td><td>0</td></tr> </tbody> </table>	x	F	0	1	1	0									
x	F																	
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1	0																	
Buffer		$F = x$	<table border="1"> <thead> <tr><th>x</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td></tr> </tbody> </table>	x	F	0	0	1	1									
x	F																	
0	0																	
1	1																	
NAND		$F = (xy)'$	<table border="1"> <thead> <tr><th>x</th><th>y</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>0</td></tr> </tbody> </table>	x	y	F	0	0	1	0	1	1	1	0	1	1	1	0
x	y	F																
0	0	1																
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1	0	1																
1	1	0																
NOR		$F = (x + y)'$	<table border="1"> <thead> <tr><th>x</th><th>y</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>0</td></tr> </tbody> </table>	x	y	F	0	0	1	0	1	0	1	0	0	1	1	0
x	y	F																
0	0	1																
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1	0	0																
1	1	0																
Exclusive-OR (XOR)		$F = xy' + x'y$ $= x \oplus y$	<table border="1"> <thead> <tr><th>x</th><th>y</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>0</td></tr> </tbody> </table>	x	y	F	0	0	0	0	1	1	1	0	1	1	1	0
x	y	F																
0	0	0																
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1	0	1																
1	1	0																
Exclusive-NOR or equivalence		$F = xy + x'y'$ $= (x \oplus y)'$	<table border="1"> <thead> <tr><th>x</th><th>y</th><th>F</th></tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	x	y	F	0	0	1	0	1	0	1	0	0	1	1	1
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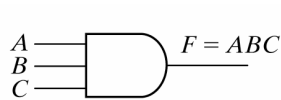
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Fig. 2-5 Digital logic gates

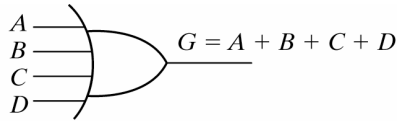


## Multi-input Logic Gates

- Many Logic Gates Can Have More than One Input
- Examples are AND and OR Gates
- Associativity Holds for These
- Can Build with Cascade of 2-input Gates



(a) Three-input AND gate



(b) Four-input OR gate

Fig. 1-6 Gates with multiple inputs

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## Multi-input Logic Gates

- Associativity Does Not Hold for NOR

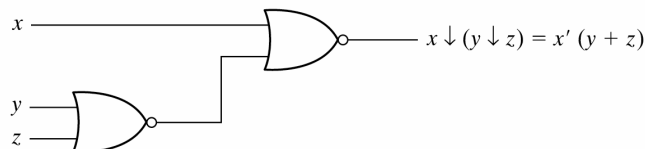
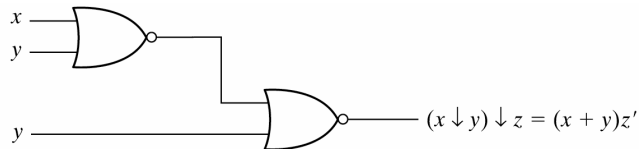
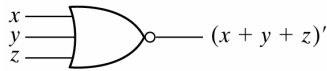


Fig. 2-6 Demonstrating the nonassociativity of the NOR operator;  $(x \downarrow y) \downarrow z \neq x \downarrow (y \downarrow z)$

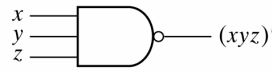
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## Multi-input Logic Gates

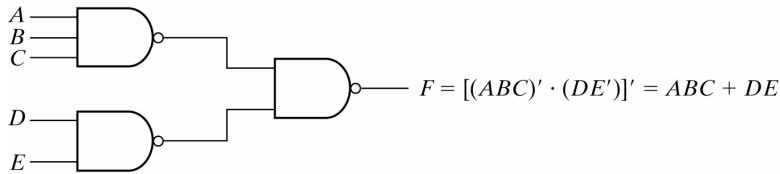
- Associativity Does Not Hold for NAND



(a) 3-input NOR gate



(b) 3-input NAND gate



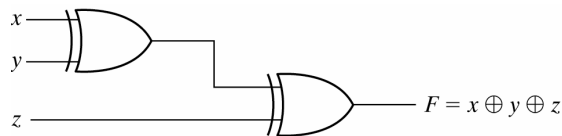
(c) Cascaded NAND gates

Fig. 2-7 Multiple-input and cascaded NOR and NAND gates

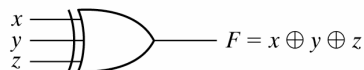
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## Exclusive-OR Gates

- Associativity Holds
- “Programmable” Inverter



(a) Using 2-input gates



(b) 3-input gate

x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

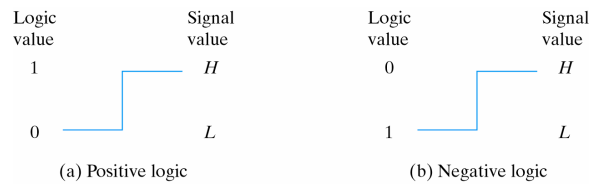
(c) Truth table

Fig. 2-8 3-input exclusive-OR gate

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## Positive and Negative Logic

- 0 and 1 are Models to use Algebra
- Circuits use High (H) and Low (L) Values
  - Voltage or Current
- Up to Designer to Interpret if  $H \rightarrow 0$  or  $H \rightarrow 1$
- Interpretations Called 'Positive'/'Negative' Logic



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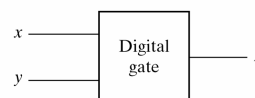
Fig. 2-9 signal assignment and logic polarity

## Positive and Negative Logic Example and Notation

*Positive-Logic  
AND becomes  
Negative-Logic  
OR*

$x$	$y$	$F$
$L$	$L$	$L$
$L$	$H$	$L$
$H$	$L$	$L$
$H$	$H$	$H$

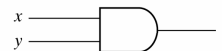
(a) Truth table with  $H$  and  $L$



(b) Gate block diagram

$x$	$y$	$z$
0	0	0
0	1	0
1	0	0
1	1	1

(c) Truth table for positive logic



(d) Positive logic AND gate

$x$	$y$	$z$
1	1	1
1	0	1
0	1	1
0	0	0

(e) Truth table for negative logic



(f) Negative logic OR gate

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Fig. 2-10 Demonstration of positive and negative logic