

On Demand secure routing protocol resilient to Byzantine failures

Primary Reference: B. Awerbuch, D. Holmer, C. Nita-Rotaru, and H. Rubens, “An on-demand secure routing protocol resilient to Byzantine failures,” in *ACM Workshop on Wireless Security (WiSe)*, September 2002.

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OUTLINE

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- PROTOCOL PHASES OVERVIEW
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- SECURE ROUTING PROTOCOL
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 - Byzantine fault detection
 - Link Weight management
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INTRODUCTION

- Ad Hoc Networks
 - No Fixed infrastructure
 - Highly suitable for emergency deployments
 - Search and rescue missions
 - High mobility and constrained power resources
 - Routing must converge quickly
- Traditional versus On demand routing
 - Distance vector and Link state consume power
 - Not designed to track topology changes at high rate
 - On Demand routing – Route request on demand
 - Routes can be cached till the topology changes

AD HOC ROUTING VULNERABILITES

- Common routing Vulnerabilities
 - Malicious nodes advertising false information
 - Try to redirect routes
 - Perform denial of service attack
 - More vulnerable to attack due to co-operative nature
- Byzantine Attacks
 - Attacker may have full control of the network
 - Performs *Black Hole* attacks
 - Create routing loops and high delay paths
- Protocol Assumptions
 - Attribute the fault to the link rather than the nodes
 - A path exists if there is a fault free link between the nodes
 - Provide routing survivability under the adversarial model

PROTOCOL PHASES OVERVIEW

- Route discovery with fault avoidance
 - Using flooding and a fault link weight list, find a least weight path
- Byzantine fault detection
 - Finds faulty links from source to destination
 - Adaptive probing technique is used
 - Faulty link is identified after $\log n$ faults
 - n is the length of the path
- Link weight management
 - Maintains a weight list of links
 - Faulty links are rehabilitated over time
 - Serves as input to the Route discovery phase

PROBLEM DEIFINITION AND MODEL

- Network Model
 - Requires bi-directional link between nodes
 - Do not address the attacks against lower layers
- Security model and considered attacks
 - Only source and destination are trusted
 - Intermediate nodes can be authenticated
 - Intermediate nodes may exhibit Byzantine behavior
- Byzantine Behavior
 - A node's behavior which causes disruption/degradation
 - Generally referred to as faults
- Problem Definition
 - Robust on demand routing protocol
 - Operating under the models above

SECURE ROUTING PROTOCOL

- On Demand protocol functionalities
 - Establishes reliability metric based on past history
 - High weights indicate low reliability
 - Each node maintains a weight list, updates it dynamically
 - Faulty links are identified using adaptive probing technique
 - Faulty links are avoided using the fault detection phase
- Phases block diagram

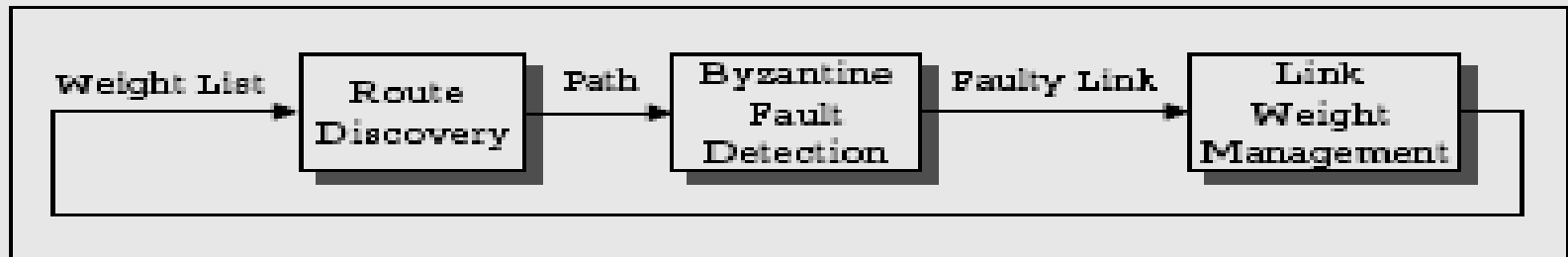


Figure 1: Secure Routing Protocol Phases

SECURE ROUTING PROTOCOL

- Route Discovery with fault avoidance
 - Flood the route request to guarantee delivery
 - Flood the response to prevent adversary from blocking the path
 - If blocked, then the route discovery will not have an input
 - No route caching is used due to security implications
 - Chooses a minimum weight path from the knowledge of LWM
 - Digital signing at intermediate nodes provides security
- Digital signatures
 - Authenticates the source sending requests
 - An unauthorized request will fail and will be dropped
 - Provides secure and legitimate route request

SECURE ROUTING PROTOCOL

- Request Initiation
 - Digitally signed request is broadcasted
- Request Propagation
 - Propagated through flooding at intermediate nodes
 - Verify matching request, source's sign and rebroadcast
- Request Receipt/Response Initiation
 - The destination verifies the request, creates response
 - Response includes sequence number and weight list
- Response propagation/receipt
 - Compute the total weight by summing the individual weights
 - If less than the previous request, verify and broadcast
 - If the path in the response is better, update the route

SECURE ROUTING PROTOCOL

- Code to perform route discovery operation

Code executed at node source when a new route to node destination is needed:

```
(1) CreateSignSend( REQUEST, destination, source, req_sequence, weight_list )
```

Code executed at node this_node when a request message req is received:

```
(2) if( Find( requests_list, req ) = NULL )
```

```
(3)     VerifySignature( req.source, req.signature )
```

```
(4)     if( this_node = req.destination )
```

```
(5)         CreateSignSend( RESPONSE, req.destination, req.source, req.req_sequence, req.high_weights_list )
```

```
(6)     else
```

```
(7)         Broadcast( req )
```

```
(8)     endif
```

```
(9)     InsertList( requests_list, req )
```

```
(10) endif
```

SECURE ROUTING PROTOCOL

- Code to perform route discovery operation

Request
initiation

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(9)     InsertList( requests_list, req )
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```
(10) endif
```

Indicates no matching requests

SECURE ROUTING PROTOCOL

Code executed at node `this_node` when a response message `res` is received:

```
(11) update = false
(12) prev_node = res.destination
(13) total_weight = 0
(14) for( i = 0; i < res.no_hops; i++ )
(15)     total_weight += LinkWeight( res.weight_list, prev_node, res.hops[i].node )
(16)     prev_node = res.hops[i].node
(17) endfor
(18) res.total_weight = total_weight + LinkWeight( res.weight_list, prev_node, this_node )
(19) prev_response = Find( responses_list, res )
(20) if( prev_response ≠ NULL )
(21)     if( res.total_weight ≥ prev_response.total_weight )
(22)         update = true
(23)     endif
(24) else
(25)     update = true
(26) endif
(27) if( update )
(28)     VerifySignature( res.destination, res.signature )
(29)     for( i = 0; i < res.no_hops; i++ )
(30)         VerifySignature( res.hops[i].node, res.hops[i].signature )
(31)     endfor
(32)     if( this_node = source )
(33)         UpdateList( path_list, res )
(34)     else
(35)         CreateSignSend( res, this_node )
(36)         UpdateList( responses_list, res )
(37)     endif
(38) endif
```

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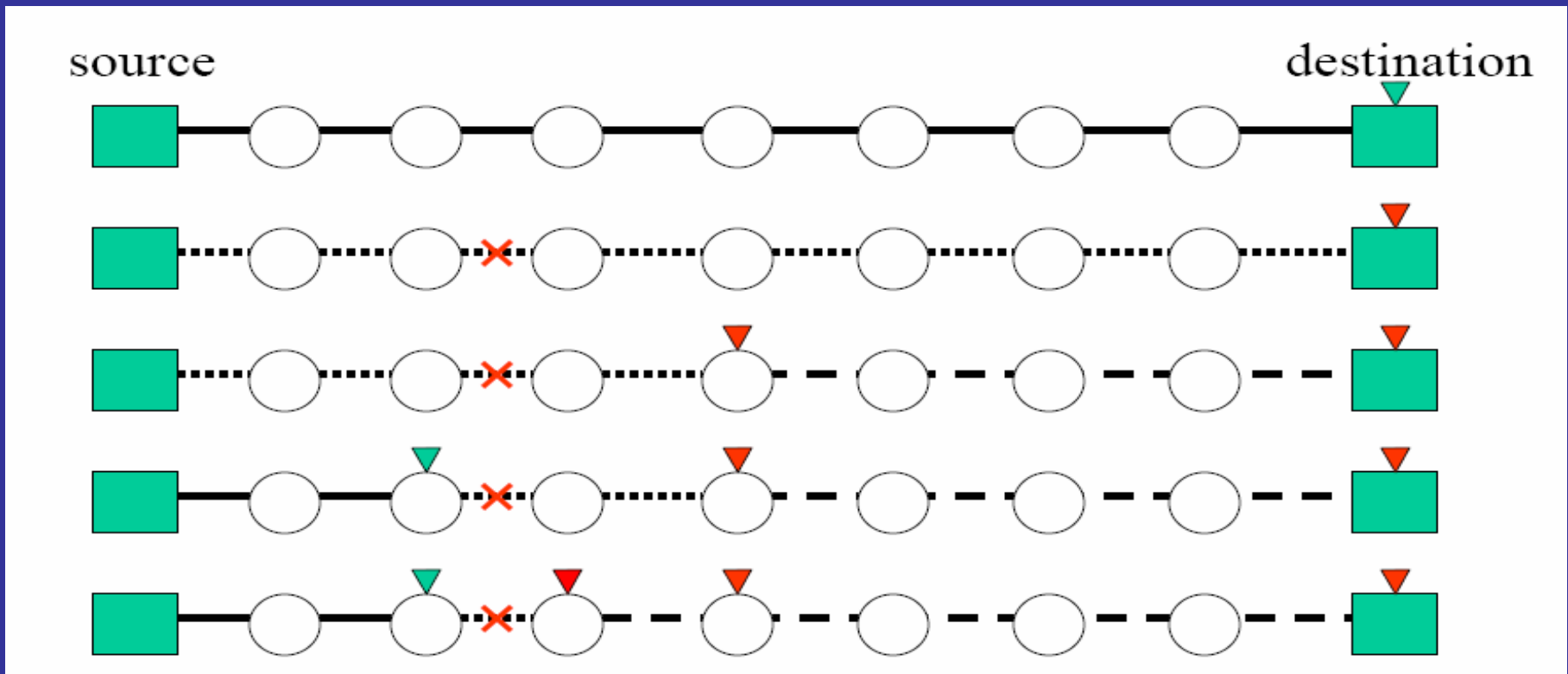
Total weight calculation

SECURE ROUTING PROTOCOL

- Byzantine Fault Detection
 - Detection algorithm uses ACKs with timeout
 - Packets not received before timeouts are considered as loss
 - Loss threshold is set which sets the limit of damage by attacker
 - Network performance is not affected due to identification of links
 - DSR Topology changes notification [2] allows optimization
 - A route error message is propagated to inform the change
- Fault Detection overview
 - Source keeps track of ACKs received
 - If no. of recent losses exceeds threshold, register a fault
 - Start a binary search on the path to identify faulty link

SECURE ROUTING PROTOCOL

- Fault detection overview
 - Probe nodes to search for faulty links



SECURE ROUTING PROTOCOL

- Probe specification
 - Specified in the list in the same order as they appear in path
 - Probe send ACK based on verifying HMAC and encryption
 - Verify the contents of the packets at all nodes
- Acknowledgement specification
 - ACKs from previous nodes are combined and sent
 - This ensures reception of ACKs from the nodes in the path
 - Thus the source discovers a loss on an interval
- Probe retirement – due to overhead
 - Counter is assigned for each joined probe
 - Successful ACK decreases the counter by one.
 - When Counter becomes zero, withdraw probe

SECURE ROUTING PROTOCOL

- Link Weight Management
 - Provides input to the route discovery process
 - Management scheme maintains the weight list
 - Weight list generated from the faulty link output of previous phase
 - When a link is faulty, its weight is doubled
 - Higher weight indicates lower reliability
 - Lower reliability routes will not comprise the path

CONCLUSION AND FUTURE WORK

- Successfully identifies faulty links
- Byzantine behavior is easily identified
- Survivability is provided even in presence of attacks
- Other Byzantine behaviors to be mitigated
- Threshold levels to be studied
- Route caching with security

REFERENCES

- [1] B. Awerbuch, D. Holmer, C. Nita-Rotaru, and H. Rubens, “An on-demand secure routing protocol resilient to Byzantine failures,” in *ACM Workshop on Wireless Security (WiSe)*, September 2002.
- [2] D.B. Johnson, D. A. Maltz, and J. Broch, DSR: The Dynamic source routing protocol for multi hop wireless ad hoc networks, in *Ad Hoc networking*, Ch. 5, pp. 139-172. Addison-Wesley, 2001.